Nevio Benvenuto | Giovanni Cherubini | Stefano Tomasin

2nd Edition

Algorithms for Comunications Systems and their Applications

WILEY

Algorithms for Communications Systems and their Applications

Algorithms for Communications Systems and their Applications

Second Edition

Nevio Benvenuto University of Padua Italy

Giovanni Cherubini IBM Research Zurich Switzerland

Stefano Tomasin University of Padua Italy

WILEY

This edition first published 2021 © 2021 John Wiley & Sons Ltd

All rights reserved. No part of this publication may be reproduced, stored in a retrieval system, or transmitted, in any form or by any means, electronic, mechanical, photocopying, recording or otherwise, except as permitted by law. Advice on how to obtain permission to reuse material from this title is available at http://www.wiley.com/go/permissions.

The right of Nevio Benvenuto, Giovanni Cherubini, and Stefano Tomasin to be identified as the authors of this work has been asserted in accordance with law.

Registered Offices

John Wiley & Sons, Inc., 111 River Street, Hoboken, NJ 07030, USA John Wiley & Sons Ltd, The Atrium, Southern Gate, Chichester, West Sussex, PO19 8SQ, UK

Editorial Office

The Atrium, Southern Gate, Chichester, West Sussex, PO19 8SQ, UK

For details of our global editorial offices, customer services, and more information about Wiley products visit us at www.wiley.com.

Wiley also publishes its books in a variety of electronic formats and by print-on-demand. Some content that appears in standard print versions of this book may not be available in other formats.

Limit of Liability/Disclaimer of Warranty

While the publisher and authors have used their best efforts in preparing this work, they make no representations or warranties with respect to the accuracy or completeness of the contents of this work and specifically disclaim all warranties, including without limitation any implied warranties of merchantability or fitness for a particular purpose. No warranty may be created or extended by sales representatives, written sales materials or promotional statements for this work. The fact that an organization, website, or product is referred to in this work as a citation and/or potential source of further information does not mean that the publisher and authors endorse the information or services the organization, website, or product may provide or recommendations it may make. This work is sold with the understanding that the publisher is not engaged in rendering professional services. The advice and strategies contained herein may not be suitable for your situation. You should consult with a specialist where appropriate. Further, readers should be aware that websites listed in this work may have changed or disappeared between when this work was written and when it is read. Neither the publisher nor authors shall be liable for any loss of profit or any other commercial damages, including but not limited to special, incidental, consequential, or other damages.

Library of Congress Cataloging-in-Publication Data

Names: Benvenuto, Nevio, author. | Cherubini, Giovanni, 1957- author. | Tomasin, Stefano, 1975- author.

Title: Algorithms for communications systems and their applications / Nevio Benvenuto, University of Padua, Italy, Giovanni Cherubini, IBM Research Zurich, Switzerland, Stefano Tomasin, University of Padua, Italy.

Description: Second edition. | Hoboken, NJ, USA : Wiley, 2021. | Includes bibliographical references and index.

- Identifiers: LCCN 2020004346 (print) | LCCN 2020004347 (ebook) | ISBN 9781119567967 (cloth) | ISBN 9781119567974 (adobe pdf) | ISBN 9781119567981 (epub)
- Subjects: LCSH: Signal processing–Mathematics. | Telecommunication systems–Mathematics. | Algorithms.

Classification: LCC TK5102.9 .B46 2020 (print) | LCC TK5102.9 (ebook) | DDC 621.382/2-dc23

LC record available at https://lccn.loc.gov/2020004346

LC ebook record available at https://lccn.loc.gov/2020004347

Cover Design: Wiley Cover Image: © betibup33/Shutterstock

Set in 9.5/12.5pt STIXGeneral by SPi Global, Chennai, India

10 9 8 7 6 5 4 3 2 1

To Adriana, to Antonio, Claudia, and Mariuccia, and in memory of Alberto

Contents

Preface xxv Acknowledgements xxvii

1	Elements of signal theory 1
1.1	Continuous-time linear systems 1
1.2	Discrete-time linear systems 2
	Discrete Fourier transform 7
	The DFT operator 7
	Circular and linear convolution via DFT 8
	Convolution by the overlap-save method 10
	IIR and FIR filters 11
1.3	Signal bandwidth 14
	The sampling theorem 17
	Heaviside conditions for the absence of signal distortion 17
1.4	Passband signals and systems 18
	Complex representation 18
	Relation between a signal and its complex representation 21
	Baseband equivalent of a transformation 26
	Envelope and instantaneous phase and frequency 28
1.5	Second-order analysis of random processes 29
1.5.1	Correlation 29
	Properties of the autocorrelation function 30
1.5.2	Power spectral density 30
	Spectral lines in the PSD 30
	Cross power spectral density 31
	Properties of the PSD 32
	PSD through filtering 32
1.5.3	PSD of discrete-time random processes 32
	Spectral lines in the PSD 33
	PSD through filtering 34
	Minimum-phase spectral factorization 35
1.5.4	PSD of passband processes 36
	PSD of in-phase and quadrature components 36
	Cyclostationary processes 38
1.6	The autocorrelation matrix 43

1.7	Examples of random processes 46
1.8	Matched filter 52
	White noise case 53
1.9	Ergodic random processes 55
1.9.1	Mean value estimators 57
	Rectangular window 58
	Exponential filter 59
	General window 59
1.9.2	Correlation estimators 60
	Unbiased estimate 60
	Biased estimate 60
1.9.3	Power spectral density estimators 61
	Periodogram or instantaneous spectrum 61
	Welch periodogram 62
	Blackman and Tukey correlogram 63
	Windowing and window closing 63
1.10	Parametric models of random processes 65
	ARMA 65
	MA 67
	AR 6/
	Spectral factorization of AR models 69
	Palatian between ADMA MA and AD models 70
1 10 1	Autocorrelation of AP processor 72
1.10.1	Spectral estimation of an AP process 72
1.10.2	Some useful relations 75
	AR model of sinusoidal processes 77
1.11	Guide to the bibliography 78
	Bibliography 78
Appendix	1.A Multirate systems 79
1.A.1	Fundamentals 79
1.A.2	Decimation 81
1.A.3	Interpolation 83
1.A.4	Decimator filter 84
1.A.5	Interpolator filter 86
1.A.6	Rate conversion 88
1.A.7	Time interpolation 90
	Linear interpolation 90
	Quadratic interpolation 91
1.A.8	The noble identities 91
1.A.9	The polyphase representation 92
	Efficient implementations 93
Appendix	1.B Generation of a complex Gaussian noise 98
Appendix	1.C Pseudo-noise sequences 99
	Maximal-length 99
	CAZAC 101
	Gold 102

2	The Wiener filter 105
2.1	The Wiener filter 105
	Matrix formulation 106
	Optimum filter design 107
	The principle of orthogonality 109
	Expression of the minimum mean-square error 110
	Characterization of the cost function surface 110
	The Wiener filter in the <i>z</i> -domain 111
2.2	Linear prediction 114
	Forward linear predictor 115
	Optimum predictor coefficients 115
	Forward prediction error filter 116
	Relation between linear prediction and AR
	models 117
	First- and second-order solutions 117
23	The least squares method 118
2.5	Data windowing 119
	Matrix formulation 119
	Correlation matrix 120
	Determination of the optimum filter coefficients 120
231	The principle of orthogonality 121
2.3.1	Minimum cost function 121
	The normal equation using the data matrix 122
	Geometric interpretation: the projection operator 122
232	Solutions to the LS problem 123
2.3.2	Singular value decomposition 124
	Minimum norm solution 125
24	The estimation problem 126
2.7	Estimation of a random variable 126
	MMSE estimation 127
	Extension to multiple observations 128
	Linear MMSE estimation of a random variable 120
	Linear MMSE estimation of a random variable 129
2 4 1	The Cremér Bee lower bound 131
2.4.1	Extension to voctor peremeter 132
2.5	Examples of application 134
2.5	Identification of a linear discrete time system 134
2.5.1	Identification of a continuous time system 135
2.5.2	Concellation of an interfering signal 138
2.5.5	Cancellation of a sinusoidal interferer with known
2.J.4	frequency 130
255	Eable concellation in digital subscriber loops 140
2.3.3	Concellation of a nominal in interform 141
2.3.0	Pibliography 142
A	Dibiliography 142
Appendix	2.A The Levinson–Durbin algorithm 142
	Lattice filters 144 The Delevete Covin eleveride 145
	The Delsarte–Genin algorithm 145

3	Adaptive transversal filters 147
31	The MSE design criterion 148
311	The steepest descent or gradient algorithm 148
5.1.1	Stability 149
	Conditions for convergence 150
	Adaptation gain 151
	Transient behaviour of the MSE 152
3.1.2	The least mean square algorithm 153
	Implementation 154
	Computational complexity 155
	Conditions for convergence 155
3.1.3	Convergence analysis of the LMS algorithm 156
	Convergence of the mean 157
	Convergence in the mean-square sense: real scalar case 157
	Convergence in the mean-square sense: general case 159
	Fundamental results 161
	Observations 162
	Final remarks 163
3.1.4	Other versions of the LMS algorithm 163
	Leaky LMS 164
	Sign algorithm 164
	Normalized LMS 164
	Variable adaptation gain 165
3.1.5	Example of application: the predictor 166
3.2	The recursive least squares algorithm 171
	Normal equation 172
	Derivation 173
	Initialization 174
	Recursive form of the minimum cost function 175
	Convergence 1/6
	Eventuational complexity 1/6
2.2	Example of application: the predictor 1//
3.3 2.2.1	Fast recursive algorithms 1//
3.3.1 3.4	Examples of application 178
3.4	Identification of a linear discrete-time system 178
5.7.1	Finite alphabet case 170
342	Cancellation of a sinusoidal interferer with known frequency 181
5.1.2	Bibliography 181
	biolography 101
4	Transmission channels 183
4.1	Radio channel 183
4.1.1	Propagation and used frequencies in radio transmission 183
	Basic propagation mechanisms 184
	Frequency ranges 184
4.1.2	Analog front-end architectures 185
	Radiation masks 185
	Conventional superheterodyne receiver 186

x

	Alternative architectures 187
	Direct conversion receiver 187
	Single conversion to low-IF 188
	Double conversion and wideband IF 188
4.1.3	General channel model 189
	High power amplifier 189
	Transmission medium 191
	Additive noise 191
	Phase noise 101
414	Narrowband radio channel model 103
4.1.4	Equivalent circuit at the receiver 105
	Multipath 106
	Dath loss as a function of distance 107
415	Factions and a function of distance 197
4.1.5	Fading effects in propagation models 200
	Macroscopic fading or shadowing 200
	Microscopic fading 201
4.1.6	Doppler shift 202
4.1.7	Wideband channel model 204
	Multipath channel parameters 205
	Statistical description of fading channels 206
4.1.8	Channel statistics 208
	Power delay profile 208
	Coherence bandwidth 209
	Doppler spectrum 210
	Coherence time 211
	Doppler spectrum models 211
	Power angular spectrum 211
	Coherence distance 212
	On fading 212
4.1.9	Discrete-time model for fading channels 213
	Generation of a process with a pre-assigned spectrum 214
4.1.10	Discrete-space model of shadowing 216
4.1.11	Multiantenna systems 218
	Line of sight 218
	Discrete-time model 219
	Small number of scatterers 220
	Large number of scatterers 220
	Blockage effect 222
4.2	Telephone channel 222
4.2.1	Distortion 222
4.2.2	Noise sources 222
	Quantization noise: 222
	Thermal noise: 224
4.2.3	Echo 224
	Bibliography 225
Appendix	4.A Discrete-time NB model for mmWave channels 226
4.A.1	Angular domain representation 226
	с

5	Vector quantization 229
5.1	Basic concept 229
5.2	Characterization of VO 230
	Parameters determining VO performance 231
	Comparison between VO and scalar quantization 232
5.3	Optimum quantization 233
	Generalized Llovd algorithm 233
5.4	The Linde, Buzo, and Gray algorithm 235
	Choice of the initial codebook 236
	Splitting procedure 236
	Selection of the training sequence 238
5.4.1	k-means clustering 239
5.5	Variants of VQ 239
	Tree search VQ 239
	Multistage VQ 240
	Product code VQ 240
5.6	VQ of channel state information 242
	MISO channel quantization 242
	Channel feedback with feedforward information 244
5.7	Principal component analysis 244
5.7.1	PCA and k-means clustering 246
	Bibliography 248
6	Digital transmission model and channel capacity 249
6.1	Digital transmission model 249
6.2	Detection 253
6.2.1	Optimum detection 253
	ML 254
())	MAP 254
6.2.2	Solid detection 200
	LLKs associated to bits of BMAP 230
602	Simplified expressions 238
0.2.3 6.2	Receiver strategies 200 Relevant parameters of the digital transmission model 260
0.5	Reletions among perspectors 261
64	Fror probability 262
6.5	Capacity 265
651	Discrete-time AWGN channel 266
652	SISO parrowband AWGN channel 266
0.5.2	Channel gain 267
6.5.3	SISO dispersive AGN channel 267
6.5.4	MIMO discrete-time NB AWGN channel 269
	Continuous-time model 270
	MIMO dispersive channel 270
6.6	Achievable rates of modulations in AWGN channels 270
6.6.1	Rate as a function of the SNR per dimension 271
6.6.2	Coding strategies depending on the signal-to-noise ratio 272
	Coding gain 274

xii

6.6.3	Achievable rate of an AWGN channel using PAM 275
	Bibliography 276
Appendix	6.A Gray labelling 277
Appendix	6.B The Gaussian distribution and Marcum functions 278
6.B.1	The Q function 278
6.B.2	Marcum function 279
7	Single-carrier modulation 281
7.1	Signals and systems 281
7.1.1	Baseband digital transmission (PAM) 281
	Modulator 281
	Transmission channel 283
	Receiver 283
	Power spectral density 284
712	Passband digital transmission (OAM) 285
,	Modulator 285
	Power spectral density 286
	Three equivalent representations of the modulator 287
	Coherent receiver 288
713	Baseband equivalent model of a OAM system 288
7.1.5	Signal analysis 288
714	Characterization of system elements 201
/.1.4	Transmitter 201
	Transmission channel 201
	Receiver 203
7.2	Intersymbol interference 204
1.2	Discrete-time equivalent system 204
	Nyquist pulses 205
	Eve diagram 208
73	Performance analysis 302
1.5	Signal to poice ratio 302
	Symbol array probability in the absence of ISL 202
	Matched filter receiver 202
74	Channel equalization 204
7.4	Zero foreing equalizer 204
7.4.1	Lineer equalizer 205
1.4.2	Ontimum receiver in the presence of poice and ISL 205
	Alternative derivation of the UD equalization 206
	Signal to poise ratio at detector 210
7 4 2	Signal-to-hoise ratio at detector 510
7.4.5	Le with a finite number of coefficients 570
	Adaptive LE 311
7 4 4	Practionally spaced equalizer 313
/.4.4	Decision feedback equalizer 315
	Design of a DFE with a finite number of coefficients 318
	Design of a fractionally spaced DFE 320
	Signal-to-noise ratio at the decision point 322
7 4 5	Kemarks 322
1.4.5	Frequency domain equalization 323

	DFE with data frame using a unique word 323
7.4.6	LE-ZF 326
7.4.7	DFE-ZF with IIR filters 327
	DFE-ZF as noise predictor 331
	DFE as ISI and noise predictor 331
7.4.8	Benchmark performance of LE-ZF and DFE-ZF 333
	Comparison 333
	Performance for two channel models 334
7.4.9	Passband equalizers 335
	Passband receiver structure 335
	Optimization of equalizer coefficients and carrier phase offset 337
	Adaptive method 338
7.5	Optimum methods for data detection 340
	Maximum a posteriori probability (MAP) criterion 341
7.5.1	Maximum-likelihood sequence detection 341
	Lower bound to error probability using MLSD 342
	The Viterbi algorithm 343
	Computational complexity of the VA 346
7.5.2	Maximum a posteriori probability detector 347
	Statistical description of a sequential machine 347
	The forward–backward algorithm 348
	Scaling 351
	The log likelihood function and the Max-Log-MAP criterion 352
	LLRs associated to bits of BMAP 353
	Relation between Max-Log–MAP and Log–MAP 354
7.5.3	Optimum receivers 354
7.5.4	The Ungerboeck's formulation of MLSD 356
7.5.5	Error probability achieved by MLSD 358
	Computation of the minimum distance 361
7.5.6	The reduced-state sequence detection 365
	Trellis diagram 365
	The RSSE algorithm 367
	Further simplification: DFSE 369
7.6	Numerical results obtained by simulations 370
	QPSK over a minimum-phase channel 370
	QPSK over a non-minimum phase channel 370
	8-PSK over a minimum phase channel 372
	8-PSK over a non-minimum phase channel 372
7.7	Precoding for dispersive channels 373
7.7.1	Tomlinson–Harashima precoding 374
7.7.2	Flexible precoding 376
7.8	Channel estimation 378
7.8.1	The correlation method 378
7.8.2	The LS method 379
	Formulation using the data matrix 380
7.8.3	Signal-to-estimation error ratio 380
	Computation of the signal-to-estimation error ratio 381
	On the selection of the channel length 384

7.8.4	Channel estimation for multirate systems 384
7.8.5	The LMMSE method 385
7.9	Faster-than-Nyquist Signalling 386
	Bibliography 387
Appendix	7.A Simulation of a QAM system 389
Appendix	7.B Description of a finite-state machine 393
Appendix	7.C Line codes for PAM systems 394
7.C.1	Line codes 394
	Non-return-to-zero format 395
	Return-to-zero format 396
	Biphase format 397
	Delay modulation or Miller code 398
	Block line codes 398
	Alternate mark inversion 398
7.C.2	Partial response systems 399
	The choice of the PR polynomial 401
	Symbol detection and error probability 404
	Precoding 406
	Error probability with precoding 407
	Alternative interpretation of PR systems 408
7.D	Implementation of a QAM transmitter 410
8	Multicarrier modulation 413
8.1	MC systems 413
8.2	Orthogonality conditions 414
	Time domain 415
	Frequency domain 415
	z-Transform domain 415
8.3	Efficient implementation of MC systems 416
	MC implementation employing matched filters 416
	Orthogonality conditions in terms of the polyphase components 418
	MC implementation employing a prototype filter 419
8.4	Non-critically sampled filter banks 422
8.5	Examples of MC systems 426
	OFDM or DMT 426
	Filtered multitone 427
8.6	Analog signal processing requirements in MC systems 429
8.6.1	
	Analog filter requirements 429
	Analog filter requirements 429 Interpolator filter and virtual subchannels 429
	Analog filter requirements429Interpolator filter and virtual subchannels429Modulator filter430
8.6.2	Analog filter requirements 429 Interpolator filter and virtual subchannels 429 Modulator filter 430 Power amplifier requirements 431
8.6.2 8.7	Analog filter requirements 429 Interpolator filter and virtual subchannels 429 Modulator filter 430 Power amplifier requirements 431 Equalization 432
8.6.2 8.7 8.7.1	Analog filter requirements 429 Interpolator filter and virtual subchannels 429 Modulator filter 430 Power amplifier requirements 431 Equalization 432 OFDM equalization 432
8.6.2 8.7 8.7.1 8.7.2	Analog filter requirements 429 Interpolator filter and virtual subchannels 429 Modulator filter 430 Power amplifier requirements 431 Equalization 432 OFDM equalization 432 FMT equalization 434
8.6.2 8.7 8.7.1 8.7.2	Analog filter requirements 429 Interpolator filter and virtual subchannels 429 Modulator filter 430 Power amplifier requirements 431 Equalization 432 OFDM equalization 432 FMT equalization 434 Per-subchannel fractionally spaced equalization 434
8.6.2 8.7 8.7.1 8.7.2	Analog filter requirements 429 Interpolator filter and virtual subchannels 429 Modulator filter 430 Power amplifier requirements 431 Equalization 432 OFDM equalization 432 FMT equalization 434 Per-subchannel fractionally spaced equalization 434 Per-subchannel <i>T</i> -spaced equalization 435
8.6.2 8.7 8.7.1 8.7.2	Analog filter requirements 429 Interpolator filter and virtual subchannels 429 Modulator filter 430 Power amplifier requirements 431 Equalization 432 OFDM equalization 432 FMT equalization 434 Per-subchannel fractionally spaced equalization 434 Per-subchannel T-spaced equalization 435 Alternative per-subchannel T-spaced equalization 436

447

	OTFS equalization 437
8.9	Channel estimation in OFDM 437
0.9	Instantaneous estimate or LS method 438
	LMMSE 440
	The LS estimate with truncated impulse response 440
8.9.1	Channel estimate and pilot symbols 441
8.10	Multiuser access schemes 442
8.10.1	OFDMA 442
8.10.2	SC-FDMA or DFT-spread OFDM 443
8.11	Comparison between MC and SC systems 444
8.12	Other MC waveforms 445
	Bibliography 446
9	Transmission over multiple input multiple output channels
9.1	The MIMO NB channel 44/
	Spatial multiplexing and spatial diversity 431
0.2	CSI only at the receiver 452
9.2	SIMO combiner 452
9.2.1	Fourier and diversity 455
922	MIMO combiner 455
1.2.2	Zero-forcing 456
	MMSE 456
9.2.3	MIMO non-linear detection and decoding 457
2.2.5	V-BLAST system 457
	Spatial modulation 458
9.2.4	Space-time coding 459
	The Alamouti code 459
	The Golden code 461
9.2.5	MIMO channel estimation 461
	The least squares method 462
	The LMMSE method 463
9.3	CSI only at the transmitter 463
9.3.1	MISO linear precoding 463
	MISO antenna selection 464
9.3.2	MIMO linear precoding 465
	ZF precoding 465
9.3.3	MIMO non-linear precoding 466
	Dirty paper coding 467
	TH precoding 468
9.3.4	Channel estimation for CSIT 469
9.4	CSI at both the transmitter and the receiver 469
9.5	Hybrid beamforming $4/0$
0.6	Hybrid beamforming and angular domain representation 4/2 Multiveer MIMO: broadcast channel 472
9.0	$\frac{1}{2}$
	CSI only at the transmitter = 473
961	CSI at both the transmitter and the receivers $\sqrt{73}$
9.0.1	Block diagonalization 473
	BIOW GIUGOHUHLAHOH 775

	User selection 474
	Joint spatial division and multiplexing 475
9.6.2	Broadcast channel estimation 476
9.7	Multiuser MIMO: multiple-access channel 476
	CSI only at the transmitters 477
	CSI only at the receiver 477
071	CSI at both the transmitters and the receiver 477
J./.1	Block diagonalization 477
072	Multiple access channel estimation 478
9.7.2	Manapire ACCess channel estimation 478
9.0	Channel hardening 478
9.8.1	Channel hardening 4/8
9.8.2	Multiuser channel orthogonality 4/9
	Bibliography 479
10	
10	Spread-spectrum systems 483
10.1	Spread-spectrum techniques 483
10.1.1	Direct sequence systems 483
	Classification of CDMA systems 490
	Synchronization 490
10.1.2	Frequency hopping systems 491
	Classification of FH systems 491
10.2	Applications of spread-spectrum systems 493
10.2.1	Anti-jamming 494
10.2.2	Multiple access 496
10.2.3	Interference rejection 496
10.3	Chip matched filter and rake receiver 496
	Number of resolvable rays in a multipath channel 497
	Chip matched filter 498
10.4	Interference 500
	Detection strategies for multiple-access systems 502
10.5	Single-user detection 502
	Chip equalizer 502
	Symbol equalizer 503
10.6	Multiuser detection 504
10.6 1	Block equalizer 504
10.6.2	Interference cancellation detector 506
10.0.2	Successive interference concellation 506
	Derellel interference concellation 507
1062	ML multiveen detector 508
10.0.3	ML multiuser detector 508
	Correlation matrix 508
	Whitening filter 508
10.7	Multicarrier CDMA systems 509
	Bibliography 510
Appendix	10.A Walsh Codes 511
11	Channel codes 515
11.1	System model 516

11.2.1 Theory of binary codes with group structure 518

Block codes 517

11.2

	Properties 518
	Parity check matrix 520
	Code generator matrix 522
	Decoding of binary parity check codes 523
	Cosets 523
	Two conceptually simple decoding methods 524
	Syndrome decoding 525
11.2.2	Fundamentals of algebra 527
	modulo- q arithmetic 528
	Polynomials with coefficients from a field 530
	Modular arithmetic for polynomials 531
	Devices to sum and multiply elements in a finite field 534
	Remarks on finite fields 535
	Roots of a polynomial 538
	Minimum function 541
	Methods to determine the minimum function 542
	Properties of the minimum function 544
11.2.3	Cyclic codes 545
	The algebra of cyclic codes 545
	Properties of cyclic codes 546
	Encoding by a shift register of length $r = 551$
	Encoding by a shift register of length $k = 552$
	Hard decoding of cyclic codes 552
	Hamming codes 554
	Burst error detection 556
11.2.4	Simplex cyclic codes 556
	Property 557
	Relation to PN sequences 558
11.2.5	BCH codes 558
	An alternative method to specify the code polynomials 558
	Bose-Chaudhuri–Hocquenhem codes 560
	Binary BCH codes 562
	Reed–Solomon codes 564
	Decoding of BCH codes 566
	Efficient decoding of BCH codes 568
11.2.6	Performance of block codes 575
11.3	Convolutional codes 576
11.3.1	General description of convolutional codes 579
	Parity check matrix 581
	Generator matrix 581
	Transfer function 582
	Catastrophic error propagation 585
11.3.2	Decoding of convolutional codes 586
	Interleaving 587
	Two decoding models 587
	Decoding by the Viterbi algorithm 588
	Decoding by the forward-backward algorithm 589
	Sequential decoding 590
11.3.3	Performance of convolutional codes 592

Contents

11.4	Puncturing 593		
11.5	Concatenated codes 593		
	The soft-output Viterbi algorithm 593		
11.6	Turbo codes 597		
	Encoding 597		
	The basic principle of iterative decoding 600		
	FBA revisited 601		
	Iterative decoding 608		
	Performance evaluation 610		
11.7	Iterative detection and decoding 611		
11.8	Low-density parity check codes 614		
11.8.1	Representation of LDPC codes 614		
	Matrix representation 614		
	Graphical representation 615		
1182	Encoding 616		
11.0.2	Encoding procedure 616		
1183	Decoding 617		
11.0.5	Hard decision decoder 617		
	The sum-product algorithm decoder 619		
	The LR-SPA decoder 622		
	The LLR-SPA or log-domain SPA decoder 623		
	The min-sum decoder 625		
	Other decoding algorithms 625		
11.8.4	Example of application 625		
11.0.4	Performance and coding gain 625		
1185	Comparison with turbo codes 627		
11.0.5	Polar codes 627		
11.9	Encoding 628		
11.9.1	Internal CRC 630		
	LLRs associated to code bits 631		
1192	Tanner graph 631		
11.9.2	Decoding algorithms 633		
11.9.5	Successive cancellation decoding – the principle 634		
	Successive cancellation decoding – the algorithm 635		
	Successive cancellation list decoding 638		
	Other decoding algorithms 639		
1194	Frozen set design 640		
11.7.4	Genie-aided SC decoding 640		
	Design based on density evolution 641		
	Channel polarization 643		
1195	Puncturing and shortening 644		
11.7.5	Puncturing $64A$		
	Shortoning 645		
	Frozen set design 647		
1106	Parformance 647		
11.9.0	FEHOIHIdile 04/ Milestenes in channel coding 649		
11.10	Bibliography 640		
Annondin	11 A Non binary parity about and a 652		
Appendix	Linear codes 653		

	Parity check matrix 654			
	Code generator matrix 655			
	Decoding of non-binary parity check codes 656			
	Two concentually simple decoding methods 656			
	Syndrome decoding 657			
	Syndrome accounting 0.57			
12	Trellis coded modulation 659			
12.1	Linear TCM for one- and two-dimensional signal sets 660			
12.1.1	Fundamental elements 660			
	Basic TCM scheme 661			
	Example 662			
12.1.2	Set partitioning 664			
12.1.3	Lattices 666			
12.1.4	Assignment of symbols to the transitions in the trellis 671			
12.1.5	General structure of the encoder/bit-mapper 675			
	Computation of d_{free} 677			
12.2	Multidimensional TCM 679			
	Encoding 680			
	Decoding 682			
12.3	Rotationally invariant TCM schemes 684			
	Bibliography 685			
13	lechniques to achieve capacity 687			
13 13.1	Capacity achieving solutions for multicarrier systems 687			
13 13.1 13.1.1	Capacity achieving solutions for multicarrier systems 687 Achievable bit rate of OFDM 687			
13 13.1 13.1.1 13.1.2	Iechniques to achieve capacity687Capacity achieving solutions for multicarrier systems687Achievable bit rate of OFDM687Waterfilling solution688			
13 13.1 13.1.1 13.1.2	Iechniques to achieve capacity687Capacity achieving solutions for multicarrier systems687Achievable bit rate of OFDM687Waterfilling solution688Iterative solution689			
13 13.1 13.1.1 13.1.2 13.1.3	Iechniques to achieve capacity687Capacity achieving solutions for multicarrier systems687Achievable bit rate of OFDM687Waterfilling solution688Iterative solution689Achievable rate under practical constraints689			
13 13.1 13.1.1 13.1.2 13.1.3	Iechniques to achieve capacity687Capacity achieving solutions for multicarrier systems687Achievable bit rate of OFDM687Waterfilling solution688Iterative solution689Achievable rate under practical constraints689Effective SNR and system margin in MC systems690			
13 13.1 13.1.1 13.1.2 13.1.3	Iechniques to achieve capacity687Capacity achieving solutions for multicarrier systems687Achievable bit rate of OFDM687Waterfilling solution688Iterative solution689Achievable rate under practical constraints689Effective SNR and system margin in MC systems690Uniform power allocation and minimum rate per subchannel690			
13 13.1 13.1.1 13.1.2 13.1.3 13.1.4	Iechniques to achieve capacity687Capacity achieving solutions for multicarrier systems687Achievable bit rate of OFDM687Waterfilling solution688Iterative solution689Achievable rate under practical constraints689Effective SNR and system margin in MC systems690Uniform power allocation and minimum rate per subchannel690The bit and power loading problem revisited691			
13 13.1 13.1.1 13.1.2 13.1.3 13.1.4	Iechniques to achieve capacity687Capacity achieving solutions for multicarrier systems687Achievable bit rate of OFDM687Waterfilling solution688Iterative solution689Achievable rate under practical constraints689Effective SNR and system margin in MC systems690Uniform power allocation and minimum rate per subchannel690The bit and power loading problem revisited691Problem formulation692			
13 13.1 13.1.1 13.1.2 13.1.3 13.1.4	Iechniques to achieve capacity687Capacity achieving solutions for multicarrier systems687Achievable bit rate of OFDM687Waterfilling solution688Iterative solution689Achievable rate under practical constraints689Effective SNR and system margin in MC systems690Uniform power allocation and minimum rate per subchannel690The bit and power loading problem revisited691Problem formulation692Some simplifying assumptions692			
13 13.1 13.1.1 13.1.2 13.1.3 13.1.4	Iechniques to achieve capacity 687Capacity achieving solutions for multicarrier systems687Achievable bit rate of OFDM687Waterfilling solution688Iterative solution689Achievable rate under practical constraints689Effective SNR and system margin in MC systems690Uniform power allocation and minimum rate per subchannel690The bit and power loading problem revisited691Problem formulation692Some simplifying assumptions692On loading algorithms693			
13 13.1 13.1.1 13.1.2 13.1.3 13.1.4	Iechniques to achieve capacity 687Capacity achieving solutions for multicarrier systems687Achievable bit rate of OFDM687Waterfilling solution688Iterative solution689Achievable rate under practical constraints689Effective SNR and system margin in MC systems690Uniform power allocation and minimum rate per subchannel690The bit and power loading problem revisited691Problem formulation692Some simplifying assumptions692On loading algorithms693The Hughes-Hartogs algorithm694			
13 13.1 13.1.1 13.1.2 13.1.3 13.1.4	Techniques to achieve capacity 687Capacity achieving solutions for multicarrier systems687Achievable bit rate of OFDM687Waterfilling solution688Iterative solution689Achievable rate under practical constraints689Effective SNR and system margin in MC systems690Uniform power allocation and minimum rate per subchannel690The bit and power loading problem revisited691Problem formulation692Some simplifying assumptions692On loading algorithms693The Hughes-Hartogs algorithm694The Krongold–Ramchandran–Jones algorithm694			
13 13.1 13.1.1 13.1.2 13.1.3 13.1.4	Iechniques to achieve capacity 687Capacity achieving solutions for multicarrier systems687Achievable bit rate of OFDM687Waterfilling solution688Iterative solution689Achievable rate under practical constraints689Effective SNR and system margin in MC systems690Uniform power allocation and minimum rate per subchannel690The bit and power loading problem revisited691Problem formulation692Some simplifying assumptions692On loading algorithms693The Hughes-Hartogs algorithm694The Krongold–Ramchandran–Jones algorithm694The Chow–Cioffi–Bingham algorithm696			
13 13.1 13.1.1 13.1.2 13.1.3 13.1.4	Iechniques to achieve capacity 687Capacity achieving solutions for multicarrier systems687Achievable bit rate of OFDM687Waterfilling solution688Iterative solution689Achievable rate under practical constraints689Effective SNR and system margin in MC systems690Uniform power allocation and minimum rate per subchannel690The bit and power loading problem revisited691Problem formulation692Some simplifying assumptions692On loading algorithms693The Hughes-Hartogs algorithm694The Chow–Cioffi–Bingham algorithm696Comparison698			
13 13.1 13.1.1 13.1.2 13.1.3 13.1.4	Iechniques to achieve capacity 687Capacity achieving solutions for multicarrier systems687Achievable bit rate of OFDM687Waterfilling solution688Iterative solution689Achievable rate under practical constraints689Effective SNR and system margin in MC systems690Uniform power allocation and minimum rate per subchannel690The bit and power loading problem revisited691Problem formulation692Some simplifying assumptions692On loading algorithms693The Hughes-Hartogs algorithm694The Chow–Cioffi–Bingham algorithm696Comparison698Capacity achieving solutions for single carrier systems698			
13 13.1 13.1.1 13.1.2 13.1.3 13.1.4 13.2	Techniques to achieve capacity 687Capacity achieving solutions for multicarrier systems687Achievable bit rate of OFDM687Waterfilling solution688Iterative solution689Achievable rate under practical constraints689Effective SNR and system margin in MC systems690Uniform power allocation and minimum rate per subchannel690The bit and power loading problem revisited691Problem formulation692Some simplifying assumptions692On loading algorithms693The Hughes-Hartogs algorithm694The Chow–Cioffi–Bingham algorithm696Comparison698Capacity achieving solutions for single carrier systems698Achieving capacity702			
13 13.1 13.1.1 13.1.2 13.1.3 13.1.4	Techniques to achieve capacity 687Capacity achieving solutions for multicarrier systems687Achievable bit rate of OFDM687Waterfilling solution688Iterative solution689Achievable rate under practical constraints689Effective SNR and system margin in MC systems690Uniform power allocation and minimum rate per subchannel690The bit and power loading problem revisited691Problem formulation692Some simplifying assumptions692On loading algorithms693The Hughes-Hartogs algorithm694The Chow-Cioffi-Bingham algorithm696Comparison698Capacity achieving solutions for single carrier systems698Achieving capacity702Bibliography703			
13 13.1 13.1.1 13.1.2 13.1.3 13.1.4 13.2 13.2	Iechniques to achieve capacity 687Capacity achieving solutions for multicarrier systems687Achievable bit rate of OFDM687Waterfilling solution688Iterative solution689Achievable rate under practical constraints689Effective SNR and system margin in MC systems690Uniform power allocation and minimum rate per subchannel690The bit and power loading problem revisited691Problem formulation692Some simplifying assumptions692On loading algorithms693The Hughes-Hartogs algorithm694The Chow–Cioffi–Bingham algorithm696Comparison698Capacity achieving solutions for single carrier systems698Achieving capacity702Bibliography703			
13 13.1 13.1.1 13.1.2 13.1.3 13.1.4 13.2 13.2 14 14.1	 Techniques to achieve capacity 687 Capacity achieving solutions for multicarrier systems 687 Achievable bit rate of OFDM 687 Waterfilling solution 688 Iterative solution 689 Achievable rate under practical constraints 689 Effective SNR and system margin in MC systems 690 Uniform power allocation and minimum rate per subchannel 690 The bit and power loading problem revisited 691 Problem formulation 692 Some simplifying assumptions 692 On loading algorithms 693 The Hughes-Hartogs algorithm 694 The Krongold–Ramchandran–Jones algorithm 694 The Chow–Cioffi–Bingham algorithm 696 Comparison 698 Capacity achieving solutions for single carrier systems 698 Achieving capacity 702 Bibliography 703 			

- 14.2 The phase-locked loop 707
- 14.2.1 PLL baseband model 708 Linear approximation 709

14.2.2	Analysis of the PLL in the presence of additive noise 711
	Noise analysis using the linearity assumption 711
14.2.3	Analysis of a second-order PLL 713
14.3	Costas loop 716
14.3.1	PAM signals 716
14.3.2	QAM signals 719
14.4	The optimum receiver 720
	Timing recovery 721
	Carrier phase recovery 725
14.5	Algorithms for timing and carrier phase recovery 725
14.5.1	ML criterion 726
	Assumption of slow time varying channel 726
14.5.2	Taxonomy of algorithms using the ML criterion 726
	Feedback estimators 727
	Early-late estimators 728
14.5.3	Timing estimators 729
	Non-data aided 729
	NDA synchronization via spectral estimation 732
	Data aided and data directed 733
	Data and phase directed with feedback: differentiator scheme 735
	Data and phase directed with feedback: Mueller and Muller scheme 735
	Non-data aided with feedback 738
14.5.4	Phasor estimators 738
	Data and timing directed 738
	Non-data aided for <i>M</i> -PSK signals 738
	Data and timing directed with feedback 739
14.6	Algorithms for carrier frequency recovery 740
14.6.1	Frequency offset estimators 741
	Non-data aided 741
	Non-data aided and timing independent with feedback 742
	Non-data aided and timing directed with feedback 743
14.6.2	Estimators operating at the modulation rate 743
	Data aided and data directed 744
	Non-data aided for <i>M</i> -PSK 744
14.7	Second-order digital PLL 744
14.8	Synchronization in spread-spectrum systems 745
14.8.1	The transmission system 745
	Transmitter 745
	Optimum receiver 745
14.8.2	Timing estimators with feedback 746
	Non-data aided: non-coherent DLL 747
	Non-data aided modified code tracking loop 747
	Data and phase directed: coherent DLL 747
14.9	Synchronization in OFDM 751
14.9.1	Frame synchronization 751
	Effects of STO 751
	Schmidl and Cox algorithm 752
14.9.2	Carrier frequency synchronization 754

14.10	Estimator performance 755 Other synchronization solutions 755 Synchronization in SC-FDMA 756 Bibliography 756
15	Self-training equalization 759
15.1	Problem definition and fundamentals 759
	Minimization of a special function 762
15.2	Three algorithms for PAM systems 765
	The Sato algorithm 765
	Benveniste–Goursat algorithm 766
	Stop-and-go algorithm 700
153	The contour algorithm for PAM systems 767
15.5	Simplified realization of the contour algorithm 769
15.4	Self-training equalization for partial response systems 770
	The Sato algorithm 770
	The contour algorithm 772
15.5	Self-training equalization for QAM systems 773
	The Sato algorithm 773
15.5.1	Constant-modulus algorithm 775
	The contour algorithm 776
	Joint contour algorithm and carrier phase tracking 777
15.6	Examples of applications 779
Ammondia	Bibliography $/83$
Appendix	(15.A) On the convergence of the contour algorithm 784
16	Low-complexity demodulators 787
16.1	Phase-shift keying 787
16.1.1	Differential PSK 787
	Error probability of <i>M</i> -DPSK 789
16.1.2	Differential encoding and coherent demodulation 791
	Differentially encoded BPSK 791
	Multilevel case 791
16.2	(D)PSK non-coherent receivers 793
16.2.1	Baseband differential detector 793
16.2.2	IF-band (1 bit) differential detector /94
1623	Signal at detection point 790
16.2.5	Optimum receivers for signals with random phase 798
10.5	ML criterion 799
	Implementation of a non-coherent ML receiver 800
	Error probability for a non-coherent binary FSK system 804
	Performance comparison of binary systems 806
16.4	Frequency-based modulations 807
16.4.1	Frequency shift keying 807
	Coherent demodulator 808
	Non-coherent demodulator 808

	Limiter–discriminator FM demodulator 809			
16.4.2	Minimum-shift keying 810			
	Power spectral density of CPFSK 812			
	Performance 814			
	MSK with differential precoding 815			
16.4.3	Remarks on spectral containment 816			
16.5	Gaussian MSK 816			
16.5.1	Implementation of a GMSK scheme 819			
	Configuration I 821			
	Configuration II 821			
	Configuration III 822			
1652	Linear approximation of a GMSK signal 824			
10.5.2	Performance of GMSK 824			
	Performance in the presence of multipath 820			
	Bibliography 830			
Appendix	16 A Continuous phase modulation 831			
Аррения	Alternative definition of CPM 831			
	Adventages of CPM 832			
	Advantages of CPM 852			
17	Applications of interference cancellation 833			
17 1	Eable and near and grossfally cancellation for DAM systems 834			
17.1	Crosstalk cancellation and full duplay transmission 825			
	Polyphasa structure of the appealler 836			
	Conceller at symbol rate 826			
	Adaptive separation 827			
	Adaptive canceller 857			
17.0	Eache compellation for OAM systems 842			
17.2	Echo cancellation for QEDM systems 642			
17.5	Echo cancellation for OFDM systems 844			
17.4	Multiuser detection for VDSL 840			
17.4.1	Upstream power back-on 850			
17.4.2	Comparison of PBO methods 851			
	Bibliography 855			
10				
10 1	Examples of communication systems 837			
18.1	The 5G cellular system 857			
18.1.1	Cells in a wireless system 857			
18.1.2	The release 15 of the 3GPP standard 858			
18.1.3	Radio access network 859			
	Time-trequency plan 859			
	NR data transmission chain 861			
	OFDM numerology 861			
	Channel estimation 862			
18.1.4	Downlink 862			
	Synchronization 863			
	Initial access or beam sweeping 864			
	Channel estimation 865			
	Channel state information reporting 865			
18.1.5	Uplink 865			

	Transform precoding numerology 866
	Channel estimation 866
	Synchronization 866
	Timing advance 867
18.1.6	Network slicing 867
18.2	GSM 868
10.2	Radio subsystem 870
183	Wireless local area networks 872
1010	Medium access control protocols 872
18.4	DECT 873
18.5	Bluetooth 875
18.6	Transmission over unshielded twisted pairs 875
18.61	Transmission over UTP in the customer service area 876
18.6.2	High-speed transmission over UTP in local area networks 880
18.7	Hybrid fibre/coaxial cable networks 881
1017	Ranging and power adjustment in OFDMA systems 885
	Ranging and power adjustment for unlink transmission 886
	Bibliography 889
Appendix	x 18.A Duplexing 890
	Three methods 890
Appendix	x 18.B Deterministic access methods 890
19	High-speed communications over twisted-pair cables 893
	5 1 1
19.1	Quaternary partial response class-IV system 893
19.1	Quaternary partial response class-IV system 893 Analog filter design 893
19.1	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894
19.1	Quaternary partial response class-IV system893Analog filter design893Received signal and adaptive gain control894Near-end crosstalk cancellation895
19.1	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894 Near-end crosstalk cancellation 895 Decorrelation filter 895
19.1	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894 Near-end crosstalk cancellation 895 Decorrelation filter 895 Adaptive equalizer 895
19.1	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894 Near-end crosstalk cancellation 895 Decorrelation filter 895 Adaptive equalizer 895 Compensation of the timing phase drift 896
19.1	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894 Near-end crosstalk cancellation 895 Decorrelation filter 895 Adaptive equalizer 895 Compensation of the timing phase drift 896 Adaptive equalizer coefficient adaptation 896
19.1	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894 Near-end crosstalk cancellation 895 Decorrelation filter 895 Adaptive equalizer 895 Compensation of the timing phase drift 896 Adaptive equalizer coefficient adaptation 896 Convergence behaviour of the various algorithms 897
19.1	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894 Near-end crosstalk cancellation 895 Decorrelation filter 895 Adaptive equalizer 895 Compensation of the timing phase drift 896 Adaptive equalizer coefficient adaptation 896 Convergence behaviour of the various algorithms 897 VLSI implementation 897
19.1 19.1.1	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894 Near-end crosstalk cancellation 895 Decorrelation filter 895 Adaptive equalizer 895 Compensation of the timing phase drift 896 Adaptive equalizer coefficient adaptation 896 Convergence behaviour of the various algorithms 897 VLSI implementation 897 Adaptive digital NEXT canceller 897
19.1	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894 Near-end crosstalk cancellation 895 Decorrelation filter 895 Adaptive equalizer 895 Compensation of the timing phase drift 896 Adaptive equalizer coefficient adaptation 896 Convergence behaviour of the various algorithms 897 VLSI implementation 897 Adaptive digital NEXT canceller 897 Adaptive digital equalizer 900
19.1 19.1.1	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894 Near-end crosstalk cancellation 895 Decorrelation filter 895 Adaptive equalizer 895 Compensation of the timing phase drift 896 Adaptive equalizer coefficient adaptation 896 Convergence behaviour of the various algorithms 897 VLSI implementation 897 Adaptive digital NEXT canceller 897 Adaptive digital equalizer 900 Timing control 904
19.1	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894 Near-end crosstalk cancellation 895 Decorrelation filter 895 Adaptive equalizer 895 Compensation of the timing phase drift 896 Adaptive equalizer coefficient adaptation 896 Convergence behaviour of the various algorithms 897 VLSI implementation 897 Adaptive digital NEXT canceller 897 Adaptive digital equalizer 900 Timing control 904 Viterbi detector 906
19.1 19.1.1 19.2	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894 Near-end crosstalk cancellation 895 Decorrelation filter 895 Adaptive equalizer 895 Compensation of the timing phase drift 896 Adaptive equalizer coefficient adaptation 896 Convergence behaviour of the various algorithms 897 VLSI implementation 897 Adaptive digital NEXT canceller 897 Adaptive digital equalizer 900 Timing control 904 Viterbi detector 906 Dual-duplex system 906
19.1 19.1.1 19.2	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894 Near-end crosstalk cancellation 895 Decorrelation filter 895 Adaptive equalizer 895 Compensation of the timing phase drift 896 Adaptive equalizer coefficient adaptation 896 Convergence behaviour of the various algorithms 897 VLSI implementation 897 Adaptive digital NEXT canceller 897 Adaptive digital equalizer 900 Timing control 904 Viterbi detector 906 Dual-duplex system 906 Dual-duplex transmission 906
19.1 19.1.1 19.2	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894 Near-end crosstalk cancellation 895 Decorrelation filter 895 Adaptive equalizer 895 Compensation of the timing phase drift 896 Adaptive equalizer coefficient adaptation 896 Convergence behaviour of the various algorithms 897 VLSI implementation 897 Adaptive digital NEXT canceller 897 Adaptive digital equalizer 900 Timing control 904 Viterbi detector 906 Dual-duplex system 906 Dual-duplex transmission 906 Physical layer control 908
19.1 19.1.1 19.2	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894 Near-end crosstalk cancellation 895 Decorrelation filter 895 Adaptive equalizer 895 Compensation of the timing phase drift 896 Adaptive equalizer coefficient adaptation 896 Convergence behaviour of the various algorithms 897 VLSI implementation 897 Adaptive digital NEXT canceller 897 Adaptive digital equalizer 900 Timing control 904 Viterbi detector 906 Dual-duplex system 906 Dual-duplex transmission 906 Physical layer control 908 Coding and decoding 909
19.1 19.1.1 19.2 19.2.1	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894 Near-end crosstalk cancellation 895 Decorrelation filter 895 Adaptive equalizer 895 Compensation of the timing phase drift 896 Adaptive equalizer coefficient adaptation 896 Convergence behaviour of the various algorithms 897 VLSI implementation 897 Adaptive digital NEXT canceller 897 Adaptive digital equalizer 900 Timing control 904 Viterbi detector 906 Dual-duplex system 906 Dual-duplex transmission 906 Physical layer control 908 Coding and decoding 909 Signal processing functions 912
19.1 19.1.1 19.2 19.2.1	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894 Near-end crosstalk cancellation 895 Decorrelation filter 895 Adaptive equalizer 895 Compensation of the timing phase drift 896 Adaptive equalizer coefficient adaptation 896 Convergence behaviour of the various algorithms 897 VLSI implementation 897 Adaptive digital NEXT canceller 897 Adaptive digital equalizer 900 Timing control 904 Viterbi detector 906 Dual-duplex system 906 Dual-duplex transmission 906 Physical layer control 908 Coding and decoding 909 Signal processing functions 912 The 100BASE-T2 transmitter 912
19.1 19.1.1 19.2 19.2.1	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894 Near-end crosstalk cancellation 895 Decorrelation filter 895 Adaptive equalizer 895 Compensation of the timing phase drift 896 Adaptive equalizer coefficient adaptation 896 Convergence behaviour of the various algorithms 897 VLSI implementation 897 Adaptive digital NEXT canceller 897 Adaptive digital equalizer 900 Timing control 904 Viterbi detector 906 Dual-duplex system 906 Dual-duplex transmission 906 Physical layer control 908 Coding and decoding 909 Signal processing functions 912 The 100BASE-T2 transmitter 913
19.1 19.1.1 19.2 19.2.1	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894 Near-end crosstalk cancellation 895 Decorrelation filter 895 Adaptive equalizer 895 Compensation of the timing phase drift 896 Adaptive equalizer coefficient adaptation 896 Convergence behaviour of the various algorithms 897 VLSI implementation 897 Adaptive digital NEXT canceller 897 Adaptive digital equalizer 900 Timing control 904 Viterbi detector 906 Dual-duplex system 906 Dual-duplex transmission 906 Physical layer control 908 Coding and decoding 909 Signal processing functions 912 The 100BASE-T2 transmitter 912 The 100BASE-T2 receiver 913 Computational complexity of digital receive filters 914
19.1 19.1.1 19.2 19.2.1	Quaternary partial response class-IV system 893 Analog filter design 893 Received signal and adaptive gain control 894 Near-end crosstalk cancellation 895 Decorrelation filter 895 Adaptive equalizer 895 Compensation of the timing phase drift 896 Adaptive equalizer coefficient adaptation 896 Convergence behaviour of the various algorithms 897 VLSI implementation 897 Adaptive digital NEXT canceller 897 Adaptive digital equalizer 900 Timing control 904 Viterbi detector 906 Dual-duplex system 906 Dual-duplex transmission 906 Physical layer control 908 Coding and decoding 909 Signal processing functions 912 The 100BASE-T2 transmitter 912 The 100BASE-T2 receiver 913 Computational complexity of digital receive filters 914 Bibliography 915

Preface

The motivation for writing this book is twofold. On the one hand, we provide a teaching tool for advanced courses in communications systems. On the other hand, we present a collection of fundamental algorithms and structures useful as an in-depth reference for researchers and engineers. The contents reflect our experience in teaching university courses on algorithms for telecommunications, as well as our professional experience acquired in industrial research laboratories.

The text illustrates the steps required for solving problems posed by the design of systems for reliable communications over wired or wireless channels. In particular, we have focused on fundamental developments in the field in order to provide the reader with the necessary insight to design practical systems.

The second edition of this book has been enriched by new solutions in fields of application and standards that have emerged since the first edition of 2002. To name one, the adoption of multiple antennas in wireless communication systems has received a tremendous impulse in recent years, and an entire chapter is now dedicated to this topic. About error correction, *polar* codes have been invented and are considered for future standards. Therefore, they also have been included in this new book edition. On the standards side, cellular networks have evolved significantly, thus we decided to dedicate a large part of a chapter to the new fifth-generation (5G) of cellular networks, which is being finalized at the time of writing. Moreover, a number of transmission techniques that have been designed and studied for application to 5G systems, with special regard to multi-carrier transmission, have been treated in this book. Lastly, many parts have been extensively integrated with new material, rewritten, and improved, with the purpose of illustrating to the reader their connection with current research trends, such as advances in machine learning.

Acknowledgements

We gratefully acknowledge all who have made the realization of this book possible. In particular, the editing of the various chapters would never have been completed without the contributions of numerous students in our courses on Algorithms for Telecommunications. Although space limitations preclude mentioning them all by name, we nevertheless express our sincere gratitude. We also thank Christian Bolis and Chiara Paci for their support in developing the software for the book, Charlotte Bolliger and Lilli M. Pavka for their assistance in administering the project, and Urs Bitterli and Darja Kropaci for their help with the graphics editing. For text processing, also for the Italian version, the contribution of Barbara Sicoli and Edoardo Casarin was indispensable; our thanks also go to Jane Frankenfield Zanin for her help in translating the text into English. We are pleased to thank the following colleagues for their invaluable assistance throughout the revision of the book: Antonio Assalini, Leonardo Bazzaco, Paola Bisaglia, Matthieu Bloch, Alberto Bononi, Alessandro Brighente, Giancarlo Calvagno, Giulio Colavolpe, Roberto Corvaja, Elena Costa, Daniele Forner, Andrea Galtarossa, Antonio Mian, Carlo Monti, Ezio Obetti, Riccardo Rahely, Roberto Rinaldo, Antonio Salloum, Fortunato Santucci, Andrea Scaggiante, Giovanna Sostrato, and Luciano Tomba. We gratefully acknowledge our colleague and mentor Jack Wolf for letting us include his lecture notes in the chapter on channel codes. We also acknowledge the important contribution of Ingmar Land on writing the section on polar codes. An acknowledgement goes also to our colleagues Werner Bux and Evangelos Eleftheriou of the IBM Zurich Research Laboratory, and Silvano Pupolin of the University of Padova, for their continuing support. Finally, special thanks go to Hideki Ochiai of Yokohama National University and Jinhong Yuan of University of New South Wales for hosting Nevio Benvenuto in the Fall 2018 and Spring 2019, respectively: both colleagues provided an ideal setting for developing the new book edition.

The Greek alphabet					
α	А	alpha	ν	Ν	nu
β	В	beta	ξ	Ξ	xi
γ	Γ	gamma	0	0	omicron
δ	Δ	delta	π	П	pi
ε, ε	Е	epsilon	ρ, ρ	Р	rho
ζ	Ζ	zeta	σ, ς	Σ	sigma
η	Н	eta	τ	Т	tau
θ, ϑ	Θ	theta	υ	Y	upsilon
ı	Ι	iota	ϕ, φ	Φ	phi
к	K	kappa	χ	Х	chi
λ	Λ	lambda	Ψ	Ψ	psi
μ	М	mu	ω	Ω	omega

To make the reading of the adopted symbols easier, the Greek alphabet is reported below.

Chapter 1

Elements of signal theory

In this chapter, we recall some concepts on signal theory and random processes. For an in-depth study, we recommend the companion book [1]. First, we introduce various forms of the Fourier transform. Next, we provide the complex representation of passband signals and their baseband equivalent. We will conclude with the study of random processes, with emphasis on the statistical estimation of first- and second-order ergodic processes, i.e. periodogram, correlogram, auto-regressive (AR), moving-average (MA), and auto-regressive moving average (ARMA) models.

1.1 Continuous-time linear systems

A time-invariant continuous-time continuous-amplitude linear system, also called analog filter, is represented in Figure 1.1, where x and y are the input and output signals, respectively, and h denotes the filter impulse response.



Figure 1.1 Analog filter as a time-invariant linear system with continuous domain.

The output at a certain instant $t \in \mathbb{R}$, where \mathbb{R} denotes the set of real numbers, is given by the convolution integral

$$y(t) = \int_{-\infty}^{\infty} h(t-\tau) x(\tau) \, d\tau = \int_{-\infty}^{\infty} h(\tau) x(t-\tau) \, d\tau$$
(1.1)

denoted in short

$$y(t) = x * h(t) \tag{1.2}$$

We also introduce the Fourier transform of the signal $x(t), t \in \mathbb{R}$,

$$\mathcal{X}(f) = \mathcal{F}[x(t)] = \int_{-\infty}^{+\infty} x(t) \ e^{-j2\pi f t} \ dt \qquad f \in \mathbb{R}$$
(1.3)

where $j = \sqrt{-1}$. The inverse Fourier transform is given by

$$x(t) = \int_{-\infty}^{\infty} \mathcal{X}(f) \ e^{j2\pi ft} \ df$$
(1.4)

Algorithms for Communications Systems and their Applications, Second Edition. Nevio Benvenuto, Giovanni Cherubini, and Stefano Tomasin.

© 2021 John Wiley & Sons Ltd. Published 2021 by John Wiley & Sons Ltd.

In the frequency domain, (1.2) becomes

$$\mathcal{Y}(f) = \mathcal{X}(f) \ \mathcal{H}(f), \quad f \in \mathbb{R}$$
 (1.5)

where \mathcal{H} is the filter frequency response. The magnitude of the frequency response, $|\mathcal{H}(f)|$, is usually called *magnitude response* or *amplitude response*.

General properties of the Fourier transform are given in Table 1.1,¹ where we use two important functions

step function:
$$1(t) = \begin{cases} 1 & t > 0 \\ 0 & t < 0 \end{cases}$$
(1.6)

sign function:
$$\operatorname{sgn}(t) = \begin{cases} 1 & t > 0 \\ -1 & t < 0 \end{cases}$$
 (1.7)

Moreover, we denote by $\delta(t)$ the *Dirac impulse* or delta function,

$$\delta(t) = \frac{d1(t)}{dt} \tag{1.8}$$

where the derivative is taken in the generalized sense.

Definition 1.1

We introduce two functions that will be extensively used:

$$\operatorname{rect}(f) = \begin{cases} 1 & |f| < \frac{1}{2} \\ 0 & \operatorname{elsewhere} \end{cases}$$
(1.9)

$$\operatorname{sinc}(t) = \frac{\sin(\pi t)}{\pi t}$$
(1.10)

The following relation holds

$$\mathcal{F}[\operatorname{sinc}(Ft)] = \frac{1}{F}\operatorname{rect}\left(\frac{f}{F}\right)$$
(1.11)

as illustrated in Figure 1.2.

Further examples of signals and relative Fourier transforms are given in Table 1.2.

We reserve the notation H(s) to indicate the Laplace transform of $h(t), t \in \mathbb{R}$:

$$H(s) = \int_{-\infty}^{+\infty} h(t)e^{-st}dt$$
(1.12)

with s complex variable; H(s) is also called the *transfer function* of the filter. A class of functions H(s) often used in practice is characterized by the ratio of two polynomials in s, each with a finite number of coefficients.

It is easy to observe that if the curve $s = j2\pi f$ in the *s*-plane belongs to the convergence region of the integral in (1.12), then $\mathcal{H}(f)$ is related to H(s) by

$$\mathcal{H}(f) = H(s)|_{s=i2\pi f} \tag{1.13}$$

1.2 Discrete-time linear systems

A discrete-time time-invariant linear system, with sampling period T_c , is shown in Figure 1.3, where x(k) and y(k) are, respectively, the input and output signals at the time instant kT_c , $k \in \mathbb{Z}$, where \mathbb{Z} denotes

¹ x^* denotes the complex conjugate of x, while Re(x) and Im(x) denote, respectively, the real and imaginary part of x.

property	signal	Fourier transform
	x(t)	$\mathcal{X}(f)$
linearity	a x(t) + b y(t)	$a \mathcal{X}(f) + b \mathcal{Y}(f)$
duality	$\mathcal{X}(t)$	x(-f)
time inverse	x(-t)	$\mathcal{X}(-f)$
complex conjugate	$x^*(t)$	$\mathcal{X}^*(-f)$
real part	$Re[x(t)] = \frac{x(t) + x^*(t)}{2}$	$\frac{1}{2} \left[\mathcal{X}(f) + \mathcal{X}^*(-f) \right]$
imaginary part	$Im[x(t)] = \frac{x(t) - x^*(t)}{2j}$	$\frac{1}{2j} \left[\mathcal{X}(f) - \mathcal{X}^*(-f) \right]$
time scaling	$x(at), a \neq 0$	$\frac{1}{ a } \mathcal{X}\left(\frac{f}{a}\right)$
time shift	$x(t-t_0)$	$e^{-j2\pi ft_0} \mathcal{X}(f)$
frequency shift	$x(t) e^{j2\pi f_0 t}$	$\mathcal{X}(f-f_0)$
modulation	$x(t)\cos(2\pi f_0t+\varphi)$	$\frac{1}{2} \left[e^{j\varphi} \mathcal{X}(f - f_0) + e^{-j\varphi} \mathcal{X}(f + f_0) \right]$
	$x(t)\sin(2\pi f_0t+\varphi)$	$\frac{1}{2j} \left[e^{j\varphi} \mathcal{X}(f - f_0) - e^{-j\varphi} \mathcal{X}(f + f_0) \right]$
	$Re[x(t) e^{j(2\pi f_0 t + \varphi)}]$	$\left \frac{1}{2}[e^{j\varphi}\mathcal{X}(f-f_0)+e^{-j\varphi}\mathcal{X}^*(-f-f_0)]\right $
differentiation	$\frac{d}{dt} x(t)$	$j2\pi f \ \mathcal{X}(f)$
integration	$\int_{-\pi}^{t} x(\tau) d\tau = 1 * x(t)$	$\frac{1}{i2\pi f} \mathcal{X}(f) + \frac{\mathcal{X}(0)}{2} \delta(f)$
convolution	x * y(t)	$\mathcal{X}(f) \mathcal{Y}(f)$
correlation	$[x(\tau)*y^*(-\tau)](t)$	$\mathcal{X}(f) \; \mathcal{Y}^*(f)$
product	x(t) y(t)	$\mathcal{X} * \mathcal{Y}(f)$
real signal	$x(t) = x^*(t)$	$\mathcal{X}(f) = \mathcal{X}^*(-f), \mathcal{X}$ Hermitian,
		$Re[\mathcal{X}(f)]$ even, $Im[\mathcal{X}(f)]$ odd,
		$ \mathcal{X}(f) ^2$ even
imaginary signal	$x(t) = -x^*(t)$	$\mathcal{X}(f) = -\mathcal{X}^*(-f)$
real and even signal	$x(t) = x^*(t) = x(-t)$	$\mathcal{X}(f) = \mathcal{X}^*(f) = \mathcal{X}(-f),$
		${\mathcal X}$ real and even
real and odd signal	$x(t) = x^*(t) = -x(-t)$	$\mathcal{X}(f) = -\mathcal{X}^*(f) = -\mathcal{X}(-f),$
		${\mathcal X}$ imaginary and odd
Parseval theorem	$E_x = \int_{-\infty}^{+\infty} x(t) ^2 dt$	$dt = \int_{-\infty}^{+\infty} \mathcal{X}(f) ^2 df = E_{\mathcal{X}}$
Poisson sum formula	$\sum_{k=-\infty}^{+\infty} x(kT_c) = \frac{1}{2}$	$\frac{1}{T_c} \sum_{\ell'=-\infty}^{+\infty} \mathcal{X}\left(\frac{\ell'}{T_c}\right)$
	~~ W	

Table 1.1: Some general properties of the Fourier transform.



Figure 1.2 Example of signal and Fourier transform pair.



Figure 1.3 Discrete-time linear system (filter).

the set of integers. We denote by $\{x(k)\}$ or $\{x_k\}$ the entire discrete-time signal, also called *sequence*. The impulse response of the system is denoted by $\{h(k)\}, k \in \mathbb{Z}$, or more simply by h.

The relation between the input sequence $\{x(k)\}$ and the output sequence $\{y(k)\}$ is given by the convolution operation:

$$y(k) = \sum_{n=-\infty}^{+\infty} h(k-n)x(n)$$
 (1.14)

denoted as y(k) = x * h(k). In the discrete time, the delta function is simply the *Kronecker impulse*

$$\delta_n = \delta(n) = \begin{cases} 1 & n = 0\\ 0 & n \neq 0 \end{cases}$$
(1.15)

Here are some definitions holding for time-invariant linear systems.

The system is *causal* (anticausal) if h(k) = 0, k < 0 (if h(k) = 0, k > 0).

sional	Fourier transform
x(t)	$\mathcal{V}(f)$
x(t)	A(f)
$\delta(t)$	1
1 (constant)	$\delta(f)$
$e^{j2\pi f_0 t}$	$\delta(f-f_0)$
$\cos(2\pi f_0 t)$	$\left \frac{1}{2}\left[\delta(f-f_0)+\delta(f+f_0)\right]\right $
$\sin(2\pi f_0 t)$	$\left \frac{1}{2j}\left[\delta(f-f_0)-\delta(f+f_0)\right]\right $
1(<i>t</i>)	$\frac{1}{2} \delta(f) + \frac{1}{j2\pi f}$
sgn(t)	$\frac{1}{j\pi f}$
$\operatorname{rect}\left(\frac{t}{T}\right)$	$T\operatorname{sinc}(fT)$
sinc $\left(\frac{t}{T}\right)$	$T \operatorname{rect}(fT)$
$\left \left(1 - \frac{ t }{T} \right) \operatorname{rect} \left(\frac{t}{2T} \right) \right $	$T\operatorname{sinc}^2(fT)$
$e^{-at} 1(t), a > 0$	$\frac{1}{a+j2\pi f}$
$t \ e^{-at} \ 1(t), a > 0$	$\frac{1}{(a+j2\pi f)^2}$
$e^{-a t }, a > 0$	$\frac{2a}{a^2 + (2\pi f)^2}$
$e^{-at^2}, a > 0$	$\sqrt{\frac{\pi}{a}}e^{-\pi\frac{\pi}{a}f^2}$

Table 1.2: Examples of Fourier transform signal pairs.

The *transfer function* of the filter is defined as the z-transform² of the impulse response h, given by

$$H(z) = \sum_{k=-\infty}^{+\infty} h(k) z^{-k}$$
(1.16)

Let the *frequency response* of the filter be defined as

$$\mathcal{H}(f) = \mathcal{F}[h(k)] = \sum_{k=-\infty}^{+\infty} h(k) e^{-j2\pi f k T_c} = H(z)_{z=e^{j2\pi f T_c}}$$
(1.17)

The inverse Fourier transform of the frequency response yields

$$h(k) = T_c \int_{-\frac{1}{2T_c}}^{+\frac{1}{2T_c}} \mathcal{H}(f) e^{j2\pi f k T_c} df$$
(1.18)

We note the property that, for $x(k) = b^k$, where *b* is a complex constant, the output is given by $y(k) = H(b) b^k$. In Table 1.3, some further properties of the z-transform are summarized.

For discrete-time linear systems, in the frequency domain (1.14) becomes

$$\mathcal{Y}(f) = \mathcal{X}(f)\mathcal{H}(f) \tag{1.19}$$

where all functions are periodic of period $1/T_c$.

² Sometimes the *D* transform is used instead of the z-transform, where $D = z^{-1}$, and H(z) is replaced by $h(D) = \sum_{k=-\infty}^{+\infty} h(k)D^k$.

property	sequence	z transform
	<i>x</i> (<i>k</i>)	X(z)
linearity	ax(k) + by(k)	aX(z) + bY(z)
time shift	x(k-m)	$z^{-m}X(z)$
complex conjugate	$x^*(k)$	$X^{*}(z^{*})$
time inverse	x(-k)	$X\left(\frac{1}{z}\right)$
	$x^{*}(-k)$	$X^*\left(\frac{1}{z^*}\right)$
z-domain scaling	$a^{-k}x(k)$	X(az)
convolution	x * y(k)	X(z)Y(z)
correlation	$x * (y^*(-m))(k)$	$X(z)Y^*\left(\frac{1}{z^*}\right)$
real sequence	$x(k) = x^*(k)$	$X(z) = X^*(z^*)$

Table 1.3: Properties of the z-transform.

Example 1.2.1 A fundamental example of z-transform is that of the sequence:

$$h(k) = \begin{cases} a^k \ k \ge 0\\ 0 \ k < 0 \end{cases}, \qquad |a| < 1$$
(1.20)

Applying the transform (1.16), we find

$$H(z) = \frac{1}{1 - az^{-1}} \tag{1.21}$$

defined for $|az^{-1}| < 1$ or |z| > |a|.

Example 1.2.2

Let $q(t), t \in \mathbb{R}$, be a continuous-time signal with Fourier transform $Q(f), f \in \mathbb{R}$. We now consider the sequence obtained by sampling q, that is

$$h_k = q(kT_c), \quad k \in \mathbb{Z}$$
(1.22)

Using the Poisson formula of Table 1.1, we have that the Fourier transform of the sequence $\{h_k\}$ is related to Q(f) by

$$\mathcal{H}(f) = \mathcal{F}[h_k] = H\left(e^{j2\pi fT_c}\right) = \frac{1}{T_c} \sum_{\ell=-\infty}^{\infty} \mathcal{Q}\left(f - \ell \frac{1}{T_c}\right)$$
(1.23)

Definition 1.2

Let us introduce the useful pulse with parameter N, a positive integer number,

$$\operatorname{sinc}_{N}(a) = \frac{1}{N} \frac{\sin(\pi a)}{\sin\left(\pi \frac{a}{N}\right)}$$
(1.24)

and $\operatorname{sinc}_N(0) = 1$. The pulse is periodic with period N(2N) if N is odd (even). For N, very large $\operatorname{sinc}_N(a)$ approximates $\operatorname{sinc}(a)$ in the range $|a| \ll N/2$.
Example 1.2.3 For the signal

$$h_k = \begin{cases} 1 & k = 0, 1, \dots, N-1 \\ 0 & \text{otherwise} \end{cases}$$
(1.25)

with sampling period T_c , it is

$$\mathcal{H}(f) = e^{-j2\pi f \frac{N-1}{2}T_c} N \operatorname{sinc}_N(f N T_c)$$
(1.26)

Discrete Fourier transform

For a sequence with a finite number of samples, $\{g_k\}, k = 0, 1, ..., N - 1$, the Fourier transform becomes

$$\mathcal{G}(f) = \sum_{k=0}^{N-1} g_k e^{-j2\pi f k T_c}$$
(1.27)

Evaluating G(f) at the points $f = m/(NT_c)$, m = 0, 1, ..., N - 1, and setting $G_m = G(m/(NT_c))$, we obtain:

$$\mathcal{G}_{m} = \sum_{k=0}^{N-1} g_{k} W_{N}^{km}, \qquad W_{N} = e^{-j\frac{2\pi}{N}}$$
(1.28)

The sequence $\{G_m\}$, m = 0, 1, ..., N - 1, is called the discrete Fourier transform (DFT) of $\{g_k\}$, k = 0, 1, ..., N - 1. The inverse of (1.28) is given by

$$g_k = \frac{1}{N} \sum_{m=0}^{N-1} \mathcal{G}_m W_N^{-km}, \qquad k = 0, 1, \dots, N-1$$
(1.29)

We note that, besides the factor 1/N, the expression of the inverse DFT (IDFT) coincides with that of the DFT, provided W_N^{-1} is substituted with W_N .

We also observe that the direct computation of (1.28) requires N(N - 1) complex additions and N^2 complex multiplications; however, the algorithm known as fast Fourier transform (FFT) computes the DFT by $N \log_2 N$ complex additions and $\left(\frac{N}{2}\log_2 N - N\right)$ complex multiplications.³

A simple implementation is also available when the DFT size is an integer power of some numbers (e.g. 2, 3, and 5). The efficient implementation of a DFT with length power of n (2, 3, and 5) is denoted as radix-n FFT. Moreover, if the DFT size is the product of integer powers of these numbers, the DFT can be implemented as a cascade of FFTs. In particular, by letting $M = 2^{\alpha_2}$, $L = 3^{\alpha_3} \cdot 5^{\alpha_5}$, the DFT of size N = LM can be implemented as the cascade of L M-size DFTs, the multiplication by twiddle factors (operating only on the phase of the signal) and an L-size DFT. Applying again the same approach to the inner M-size DFT, we obtain that the N-size DFT is the cascade of 2^{α_2} FFTs of size $3^{\alpha_3}5^{\alpha_5}$, each implemented by 3^{α_3} FFTs of size 5^{α_5} .

The DFT operator

The DFT operator can be expressed in matrix form as

$$\boldsymbol{F} = \begin{bmatrix} 1 & 1 & 1 & \dots & 1 \\ 1 & W_N & W_N^2 & \dots & W_N^{(N-1)} \\ 1 & W_N^2 & W_N^4 & \dots & W_N^{2(N-1)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & W_N^{(N-1)} & W_N^{(N-1)2} & \dots & W_N^{(N-1)(N-1)} \end{bmatrix}$$
(1.30)

³ The computational complexity of the FFT is often expressed as $N \log_2 N$.

with elements $[F]_{i,n} = W_N^{in}$, i, n = 0, 1, ..., N - 1. The inverse operator (IDFT) is given by⁴

$$F^{-1} = \frac{1}{N}F^* \tag{1.31}$$

We note that $F = F^T$ and $(1/\sqrt{N})F$ is a unitary matrix.⁵

The following property holds: if C is a right circulant square matrix, i.e. its rows are obtained by successive shifts to the right of the first row, then FCF^{-1} is a diagonal matrix whose elements are given by the DFT of the first row of C. This property is exploited in the most common modulation scheme (see Chapter 8).

Introducing the vector formed by the samples of the sequence $\{g_k\}, k = 0, 1, \dots, N-1$,

$$\boldsymbol{g}^{T} = [g_0, g_1, \dots, g_{N-1}] \tag{1.32}$$

and the vector of its transform coefficients

$$\boldsymbol{\mathcal{G}}^{T} = [\mathcal{G}_{0}, \mathcal{G}_{1}, \dots, \mathcal{G}_{N-1}] = \text{DFT}[\boldsymbol{g}]$$
 (1.33)

from (1.28) we have

$$\boldsymbol{\mathcal{G}} = \boldsymbol{F}\boldsymbol{g} \tag{1.34}$$

Moreover, based on (1.31), we obtain

$$g = \frac{1}{N} F^* \mathcal{G} \tag{1.35}$$

Circular and linear convolution via DFT

Let the two sequences x and h have a finite support of L_x and N samples, respectively, (see Figure 1.4) with $L_x > N$:

$$x(k) = 0 k < 0 k > L_x - 1 (1.36)$$

and

$$h(k) = 0 \qquad k < 0 \qquad k > N - 1 \tag{1.37}$$

We define the periodic signals of period L,



Figure 1.4 Time-limited signals: $\{x(k)\}, k = 0, 1, ..., L_x - 1$, and $\{h(k)\}, k = 0, 1, ..., N - 1$.

⁴ T stands for transpose and ^{*H*} for transpose complex conjugate or Hermitian.

⁵ A square matrix A is unitary if $A^{H}A = I$, where I is the identity matrix, i.e. a matrix for which all elements are zero except the elements on the main diagonal that are all equal to one.

where in order to avoid time aliasing, it must be

$$L \ge \max\{L_x, N\} \tag{1.39}$$

Definition 1.3

The circular convolution between x and h is a periodic sequence of period L defined as

$$y^{(circ)}(k) = h \bigotimes^{L} x(k) = \sum_{i=0}^{L-1} h_{rep_{L}}(i) \ x_{rep_{L}}(k-i)$$
(1.40)

with *main period* corresponding to k = 0, 1, ..., L - 1.

Then, if we indicate with $\{\mathcal{X}_m\}$, $\{\mathcal{H}_m\}$, and $\{\mathcal{Y}_m^{(circ)}\}$, m = 0, 1, ..., L - 1, the *L*-point DFT of sequences *x*, *h*, and $y^{(circ)}$, respectively, we obtain

$$\mathcal{Y}_m^{(circ)} = \mathcal{X}_m \mathcal{H}_m, \qquad m = 0, 1, \dots, L - 1 \tag{1.41}$$

In vector notation (1.33), (1.41) becomes⁶

$$\boldsymbol{\mathcal{Y}}^{(circ)} = \left[\boldsymbol{\mathcal{Y}}_{0}^{(circ)}, \boldsymbol{\mathcal{Y}}_{1}^{(circ)}, \dots, \boldsymbol{\mathcal{Y}}_{L-1}^{(circ)}\right]^{T} = \text{diag}\left\{\text{DFT}[\boldsymbol{x}]\right\}\boldsymbol{\mathcal{H}}$$
(1.42)

where \mathcal{H} is the column vector given by the *L*-point DFT of the sequence *h*, completed with L - N zeros. We are often interested in the *linear convolution* between *x* and *h* given by (1.14):

$$y(k) = x * h(k) = \sum_{i=0}^{N-1} h(i)x(k-i)$$
(1.43)

whose support is $k = 0, 1, \dots, L_x + N - 2$.

We give below two relations between the circular convolution $y^{(circ)}$ and the linear convolution y.

Relation 1. For

$$L \ge L_r + N - 1 \tag{1.44}$$

by comparing (1.43) with (1.40), the two convolutions $y^{(circ)}$ and y coincide only for the instants k = 0, 1, ..., L - 1, i.e.

$$y(k) = y^{(circ)}(k), \quad k = 0, 1, \dots, L-1$$
 (1.45)

To compute the convolution between the two finite-length sequences x and h, (1.44) and (1.45) require that both sequences be completed with zeros (zero padding) to get a length of $L = L_x + N - 1$ samples. Then, taking the *L*-point DFT of the two sequences, performing the product (1.41), and taking the inverse transform of the result, one obtains the desired linear convolution.

Relation 2. For $L = L_x > N$, the two convolutions $y^{(circ)}$ and y coincide only for the instants k = N - 1, $N, \dots, L - 1$, i.e.

$$y^{(circ)}(k) = y(k)$$
 only for $k = N - 1, N, \dots, L - 1$ (1.46)

An example of circular convolution is provided in Figure 1.5. Indeed, the result of circular convolution coincides with $\{y(k)\}$, output of the linear convolution, only for a delay k such that it is avoided the product between non-zero samples of the two periodic sequences h_{rep_L} and x_{rep_L} , indicated by • and o, respectively. This is achieved only for $k \ge N - 1$ and $k \le L - 1$.

⁶ The notation diag $\{v\}$ denotes a diagonal matrix whose elements on the diagonal are equal to the elements of the vector v.



Figure 1.5 Illustration of the *circular convolution* operation between $\{x(k)\}$, k = 0, 1, ..., L - 1, and $\{h(k)\}$, k = 0, 1, ..., N - 1.

Relation 3. A relevant case wherein the cyclic convolution is equivalent to the linear convolution requires a special structure of the sequence x. Consider $x^{(cp)}$, the extended sequence of x, obtained by partially repeating x with a *cyclic prefix* of N_{cp} samples:

$$x^{(cp)}(k) = \begin{cases} x(k) & k = 0, 1, \dots, L_x - 1\\ x(L_x + k) & k = -N_{cp}, \dots, -2, -1 \end{cases}$$
(1.47)

Let $y^{(cp)}$ be the *linear convolution* between $x^{(cp)}$ and h, with support $\{-N_{cp}, \dots, L_x + N - 2\}$. If $N_{cp} \ge N - 1$, we have

$$y^{(cp)}(k) = y^{(circ)}(k), \quad k = 0, 1, \dots, L_x - 1$$
 (1.48)

Let us define

$$z(k) = \begin{cases} y^{(cp)}(k) & k = 0, 1, \dots, L_x - 1\\ 0 & \text{elsewhere} \end{cases}$$
(1.49)

then from (1.48) and (1.41) the following relation between the corresponding L_x -point DFTs is obtained:

$$\mathcal{Z}_m = \mathcal{X}_m \mathcal{H}_m, \qquad m = 0, 1, \dots, L_x - 1 \tag{1.50}$$

Convolution by the overlap-save method

For a very long sequence x, the application of (1.46) leads to the *overlap-save* method to determine the linear convolution between x and h (with $L = L_x > N$). It is not restrictive to assume that the first (N - 1) samples of the sequence $\{x(k)\}$ are zero. If this were not True, it would be sufficient to shift the input by (N - 1) samples. A fast procedure to compute the linear convolution $\{y(k)\}$ for instants $k = N - 1, N, \dots, L - 1$, operates iteratively and processes blocks of L samples, where adjacent blocks are overlapping by (N - 1) samples. The procedure operates the following first iteration:⁷

1. Loading

$$\boldsymbol{h}^{\prime T} = [h(0), h(1), \dots, h(N-1), 0, \dots, 0]$$
(1.51)

$$\mathbf{x}^{\prime T} = [x(0), x(1), \dots, x(N-1), x(N), \dots, x(L-1)]$$
(1.52)

in which we have assumed $x(k) = 0, k = 0, 1, \dots, N - 2$.

⁷ In this section, the superscript ' indicates a vector of L elements.

2. Transform

 $\mathcal{H}' = \mathrm{DFT}[\mathbf{h}']$ vector (1.53)

$$\mathcal{X}' = \text{diag}\{\text{DFT}[\mathbf{x}']\} \quad \text{matrix} \tag{1.54}$$

3. Matrix product

$$\mathcal{Y}' = \mathcal{X}' \mathcal{H}'$$
 vector (1.55)

4. Inverse transform

$$\mathbf{y}^{\prime T} = \mathrm{DFT}^{-1}[\mathbf{\mathcal{Y}}^{\prime T}] = [\ddagger, \dots, \ddagger, y(N-1), y(N), \dots, y(L-1)]$$
(1.56)

where the symbol \sharp denotes a component that is neglected.

The second iteration operates on load

$$\mathbf{x}^{\prime T} = [x((L-1) - (N-2)), \dots, x(2(L-1) - (N-2))]$$
(1.57)

and the desired output samples will be

$$y(k)$$
 $k = L, ..., 2(L-1) - (N-2)$ (1.58)

The third iteration operates on load

$$\mathbf{x}^{\prime T} = [x(2(L-1) - 2(N-2)), \dots, x(3(L-1) - 2(N-2))]$$
(1.59)

and will yield the desired output samples

$$y(k)$$
 $k = 2(L-1) - (N-2) + 1, \dots, 3(L-1) - 2(N-2)$ (1.60)

The algorithm proceeds iteratively until the entire input sequence is processed.

IIR and FIR filters

An important class of linear systems is identified by the input-output relation

$$\sum_{n=0}^{p} a_n y(k-n) = \sum_{n=0}^{q} b_n x(k-n)$$
(1.61)

where we will set $a_0 = 1$ without loss of generality.

If the system is causal, (1.61) becomes

$$y(k) = -\sum_{n=1}^{p} a_n y(k-n) + \sum_{n=0}^{q} b_n x(k-n) \qquad k \ge 0$$
(1.62)

and the transfer function for such system is

$$H(z) = \frac{Y(z)}{X(z)} = \frac{\sum_{n=0}^{q} b_n z^{-n}}{1 + \sum_{n=1}^{p} a_n z^{-n}} = \frac{b_0 \prod_{n=1}^{q} (1 - z_n z^{-1})}{\prod_{n=1}^{p} (1 - p_n z^{-1})}$$
(1.63)

where $\{z_n\}$ and $\{p_n\}$ are, respectively, the zeros and poles of H(z). Equation (1.63) generally defines an *infinite impulse response* (IIR) filter. In the case in which $a_n = 0, n = 1, 2, ..., p$, (1.63) reduces to

$$H(z) = \sum_{n=0}^{q} \mathbf{b}_{n} z^{-n}$$
(1.64)

	h(0)	h(1)	h(2)	h(3)	<i>h</i> (4)
h_1 (minimum phase)	$0.9e^{-j1.57}$	0	0	$0.4e^{-j0.31}$	$0.3e^{-j0.63}$
h_2 (maximum phase)	$0.3e^{j0.63}$	$0.4e^{j0.31}$	0	0	$0.9e^{j1.57}$
h_3 (general case)	$0.7e^{-j1.57}$	$0.24e^{j2.34}$	$0.15e^{-j1.66}$	$0.58e^{-j0.51}$	$0.4e^{-j0.63}$

Table 1.4: Impulse responses of systems having the same magnitude of the frequency response.

and we obtain a *finite impulse response* (FIR) filter, with $h(n) = b_n$, n = 0, 1, ..., q. To get the impulse response coefficients, assuming that the z-transform H(z) is known, we can expand H(z) in partial fractions and apply the linear property of the z-transform (see Table 1.3, page 6). If q < p, and assuming that all poles are distinct, we obtain

$$H(z) = \sum_{n=1}^{p} \frac{r_n}{1 - p_n z^{-1}} \implies h(k) = \begin{cases} \sum_{n=1}^{p} r_n p_n^k & k \ge 0\\ 0 & k < 0 \end{cases}$$
(1.65)

where

$$r_n = H(z) \left[1 - p_n z^{-1} \right] \Big|_{z=p_n}$$
(1.66)

We give now two definitions.

Definition 1.4

A causal system is *stable* (bounded input-bounded output stability) if $|\mathbf{p}_n| < 1, \forall n$.

Definition 1.5 The system is minimum phase (maximum phase) if $|p_n| < 1$ and $|z_n| \le 1$ ($|p_n| > 1$ and $|z_n| > 1$), $\forall n$.

Among all systems having the same magnitude response $|\mathcal{H}(e^{i2\pi fT_c})|$, the minimum (maximum) phase system presents a phase⁸ response, arg $\mathcal{H}(e^{i2\pi fT_c})$, which is below (above) the phase response of all other systems.

Example 1.2.4

It is interesting to determine the phase of a system for a given impulse response. Let us consider the system with transfer function $H_1(z)$ and impulse response $h_1(k)$ shown in Figure 1.6a. After determining the zeros of the transfer function, we factorize $H_1(z)$ as follows:

$$H_1(z) = \mathbf{b}_0 \prod_{n=1}^4 (1 - \mathbf{z}_n z^{-1})$$
(1.67)

As shown in Figure 1.6a, $H_1(z)$ is minimum phase. We now observe that the magnitude of the frequency response does not change if $1/z_n^*$ is replaced with z_n in (1.67). If we move all the zeros outside the unit circle, we get a maximum-phase system $H_2(z)$ whose impulse response is shown in Figure 1.6b. A general case, that is a transfer function with some zeros inside and others outside the unit circle, is given in Figure 1.6c. The coefficients of the impulse responses h_1 , h_2 , and h_3 are given in Table 1.4. The coefficients are normalized so that the three impulse responses have equal energy.

⁸ For a complex number c, arg c denotes the phase of c.



Figure 1.6 Impulse response magnitudes and zero locations for three systems having the same frequency response magnitude. (a) Minimum-phase system, (b) maximum-phase system, and (c) general system.

We define the *partial energy* of a causal impulse response as

$$E(k) = \sum_{i=0}^{k} |h(i)|^2$$
(1.68)

Comparing the partial-energy sequences for the three impulse responses of Figure 1.6, one finds that the minimum (maximum) phase system yields the largest (smallest) $\{E(k)\}$. In other words, the magnitude



Figure 1.7 Classification of real valued analog filters on the basis of the support of $|\mathcal{H}(f)|$. (a) $f_1 = 0, f_2 < \infty$: lowpass filter (LPF). (b) $f_1 > 0, f_2 = \infty$: highpass filter (HPF). (c) $f_1 > 0, f_2 < \infty$: passband filter (PBF). (d) $B = f_2 - f_1 \ll (f_2 + f_1)/2$: narrowband filter (NBF). (e) $f_1 = 0, f_2 = \infty$: allpass filter (APF).

of the frequency responses being equal, a minimum (maximum) phase system concentrates all its energy on the first (last) samples of the impulse response.

Extending our previous considerations also to IIR filters, if h_1 is a causal minimum-phase filter, i.e. $H_1(z) = H_{min}(z)$ is a ratio of polynomials in z^{-1} with poles and zeros inside the unit circle, then $H_{max}(z) = K H_{min}^* \left(\frac{1}{z^*}\right)$, where K is a constant, is an anticausal maximum-phase filter, i.e. $H_{max}(z)$ is a ratio of polynomials in z with poles and zeros outside the unit circle.

In the case of a minimum-phase FIR filter with impulse response $h_{min}(n)$, n = 0, 1, ..., q, $H_2(z) = z^{-q} H^*_{min}\left(\frac{1}{z^*}\right)$ is a causal maximum-phase filter. Moreover, the relation $\{h_2(n)\} = \{h_1^*(q-n)\}$, n = 0, 1, ..., q, is satisfied. In this text, we use the notation $\{h_2(n)\} = \{h_1^{B*}(n)\}$, where *B* is the *backward* operator that orders the elements of a sequence from the last to the first.

In Appendix 1.A multirate transformations for systems are described, in which the time domain of the input is different from that of the output. In particular, decimator and interpolator filters are introduced, together with their efficient implementations.

1.3 Signal bandwidth

Definition 1.6

The support of a signal $x(\xi), \xi \in \mathbb{R}$, is the set of values $\xi \in \mathbb{R}$ for which $|x(\xi)| \neq 0$.

Let us consider a filter with impulse response h and frequency response H. If h assumes real values, then H is Hermitian, $H(-f) = H^*(f)$, and |H(f)| is an even function. Depending on the support



Figure 1.8 Classification of complex-valued analog filters on the basis of support of $|\mathcal{H}(f)|$. (a) $-\infty < f_1 \le 0, \ 0 < f_2 < \infty$: lowpass filter. (b) $f_1 > 0, \ f_2 < \infty$: passband filter. (c) $f_1 > -\infty, \ f_2 < 0, \ f_3 > 0, \ f_4 < \infty$: passband filter.

of $|\mathcal{H}(f)|$, the classification of Figure 1.7 is usually done. If *h* assumes complex values, the terminology is less standard. We adopt the classification of Figure 1.8, in which the filter is a lowpass filter (LPF) if the support $|\mathcal{H}(f)|$ includes the origin; otherwise, it is a passband filter (PBF).

Analogously, for a signal x, we will use the same denomination and we will say that x is a baseband (BB) or passband (PB) signal depending on whether the support of $|\mathcal{X}(f)|, f \in \mathbb{R}$, includes or not the origin.

Definition 1.7 In general, for a *real-valued signal x*, the set of positive frequencies such that $|\mathcal{X}(f)| \neq 0$ is called *pass-band* or simply *band B*:

$$\mathcal{B} = \{ f \ge 0 : |\mathcal{X}(f)| \ne 0 \}$$
(1.69)

As $|\mathcal{X}(f)|$ is an even function, we have $|\mathcal{X}(-f)| \neq 0, f \in \mathcal{B}$. We note that \mathcal{B} is equivalent to the support of \mathcal{X} limited to positive frequencies. The bandwidth of x is given by the measure of \mathcal{B} :

$$B = \int_{\mathcal{B}} df \tag{1.70}$$

In the case of a *complex-valued signal x*, \mathcal{B} is equivalent to the support of \mathcal{X} , and \mathcal{B} is thus given by the measure of the entire support.

Observation 1.1

The signal *bandwidth* may also be given different practical definitions. Let us consider an LPF having frequency response $\mathcal{H}(f)$. The filter gain \mathcal{H}_0 is usually defined as $\mathcal{H}_0 = |\mathcal{H}(0)|$; other definitions of gain refer to the average gain of the filter in the passband \mathcal{B} , or as $\max_f |\mathcal{H}(f)|$. We give the following four definitions for the bandwidth B of h:

(a) First zero:

$$B = \min\{f > 0 : \mathcal{H}(f) = 0\}$$
(1.71)

(b) *Based on amplitude*, bandwidth at *A* dB:

$$B = \max\left\{f > 0 : \frac{|\mathcal{H}(f)|}{\mathcal{H}_0} = 10^{-\frac{A}{20}}\right\}$$
(1.72)

Typically, A = 3, 40, or 60.

(c) *Based on energy*, bandwidth at *p*%:

$$\frac{\int_{0}^{B} |\mathcal{H}(f)|^{2} df}{\int_{0}^{\infty} |\mathcal{H}(f)|^{2} df} = \frac{p}{100}$$
(1.73)

Typically, p = 90 or 99.(d) *Equivalent noise bandwidth*:

$$B = \frac{\int_0^\infty |\mathcal{H}(f)|^2 df}{\mathcal{H}_0^2} \tag{1.74}$$

Figure 1.9 illustrates the various definitions for a particular $|\mathcal{H}(f)|$. For example, with regard to the signals of Figure 1.7, we have that for an LPF $B = f_2$, whereas for a PBF $B = f_2 - f_1$.

For discrete-time filters, for which \mathcal{H} is periodic of period $1/T_c$, the same definitions hold, with the caution of considering the support of $|\mathcal{H}(f)|$ within a period, let us say between $-1/(2T_c)$ and $1/(2T_c)$. In the case of discrete-time highpass filters (HPFs), the passband will extend from a certain frequency f_1 to $1/(2T_c)$.



Figure 1.9 The real signal bandwidth following the definitions of (1) bandwidth at first zero: $B_z = 0.652$ Hz; (2) amplitude-based bandwidth: $B_{3\ dB} = 0.5$ Hz, $B_{40\ dB} = 0.87$ Hz, $B_{50\ dB} = 1.62$ Hz; (3) energy-based bandwidth: $B_{E(p=90)} = 1.362$ Hz, $B_{E(p=99)} = 1.723$ Hz; (4) equivalent noise bandwidth: $B_{req} = 0.5$ Hz.

The sampling theorem

As discrete-time signals are often obtained by sampling continuous-time signals, we will state the following fundamental theorem.

Theorem 1.1 (Sampling theorem)

Let $q(t), t \in \mathbb{R}$ be a continuous-time signal, in general complex-valued, whose Fourier transform Q(f) has support within an interval \mathcal{B} of finite measure B_0 . The samples of the signal q, taken with period T_c as represented in Figure 1.10a,

$$h_k = q(kT_c) \tag{1.75}$$

univocally represent the signal $q(t), t \in \mathbb{R}$, under the condition that the sampling frequency $1/T_c$ satisfies the relation

$$\frac{1}{T_c} \ge B_0 \tag{1.76}$$



Figure 1.10 Operation of (a) sampling and (b) interpolation.

For the proof, which is based on the relation (1.23) between a signal and its samples, we refer the reader to [2].

 B_0 is often referred to as the minimum sampling frequency. If $1/T_c < B_0$ the signal cannot be perfectly reconstructed from its samples, originating the so-called *aliasing phenomenon* in the frequency-domain signal representation.

In turn, the signal q(t), $t \in \mathbb{R}$, can be reconstructed from its samples $\{h_k\}$ according to the scheme of Figure 1.10b, where it is employed an interpolation filter having an ideal frequency response given by

$$\mathcal{G}_{I}(f) = \begin{cases} 1 & f \in \mathcal{B} \\ 0 & \text{elsewhere} \end{cases}$$
(1.77)

We note that for *real-valued baseband signals* $B_0 = 2B$. For passband signals, care must be taken in the choice of $B_0 \ge 2B$ to avoid *aliasing* between the positive and negative frequency components of Q(f).

Heaviside conditions for the absence of signal distortion

Let us consider a filter having frequency response $\mathcal{H}(f)$ (see Figures 1.1 or 1.3) given by

$$\mathcal{H}(f) = \mathcal{H}_0 e^{-j2\pi f t_0}, \quad f \in \mathcal{B}$$
(1.78)

where \mathcal{H}_0 and t_0 are two non-negative constants, and \mathcal{B} is the *passband* of the filter input signal x. Then the output is given by

$$\mathcal{Y}(f) = \mathcal{H}(f)\mathcal{X}(f) = \mathcal{H}_0\mathcal{X}(f) \ e^{-j2\pi f t_0} \tag{1.79}$$

or, in the time domain,

$$y(t) = \mathcal{H}_0 x(t - t_0)$$
 (1.80)



Figure 1.11 Characteristics of a filter satisfying the conditions for the absence of signal distortion in the frequency interval (f_1, f_2) . (a) Magnitude and (b) phase.

In other words, for a filter of the type (1.78), the signal at the input is reproduced at the output with a gain factor \mathcal{H}_0 and a delay t_0 .

A filter of the type (1.78) satisfies the Heaviside conditions for the absence of signal distortion and is characterized by

1. Constant magnitude

$$|\mathcal{H}(f)| = \mathcal{H}_0, \quad f \in \mathcal{B} \tag{1.81}$$

2. Linear phase

$$\arg \mathcal{H}(f) = -2\pi f t_0, \quad f \in \mathcal{B}$$
(1.82)

3. Constant group delay, also called envelope delay

$$\tau(f) = -\frac{1}{2\pi} \frac{d}{df} \arg \mathcal{H}(f) = t_0, \quad f \in \mathcal{B}$$
(1.83)

We underline that it is sufficient that the Heaviside conditions are verified within the support of \mathcal{X} ; as $|\mathcal{X}(f)| = 0$ outside the support, the filter frequency response may be arbitrary.

We show in Figure 1.11 the frequency response of a PBF, with bandwidth $B = f_2 - f_1$, that satisfies the conditions stated by Heaviside.

1.4 Passband signals and systems

We now provide a compact representation of passband signals and describe their transformation by linear systems.

Complex representation

For a passband signal x, it is convenient to introduce an equivalent representation in terms of a baseband signal $x^{(bb)}$.

Let x be a PB real-valued signal with Fourier transform as illustrated in Figure 1.12. The following two procedures can be adopted to obtain $x^{(bb)}$.



Figure 1.12 Transformations to obtain the baseband equivalent signal $x^{(bb)}$ around the carrier frequency f_0 using a phase splitter.



Figure 1.13 Transformations to obtain the baseband equivalent signal $x^{(bb)}$ around the carrier frequency f_0 using a phase splitter.

PB filter. Referring to Figure 1.12 and to the transformations illustrated in Figure 1.13, given x we extract its positive frequency components using an *analytic filter* or *phase splitter*, $h^{(a)}$, having the following ideal frequency response

$$\mathcal{H}^{(a)}(f) = 2 \cdot 1(f) = \begin{cases} 2 & f > 0\\ 0 & f < 0 \end{cases}$$
(1.84)

In practice, it is sufficient that $h^{(a)}$ is a complex PB filter, with $\mathcal{H}^{(a)}(f) \simeq 2$ in the passband that extends from f_1 to f_2 , as $\mathcal{X}(f)$, and stopband, in which $|\mathcal{H}^{(a)}(f)| \simeq 0$, that extends from $-f_2$ to $-f_1$. The signal $x^{(a)}$ is called the analytic signal or pre-envelope of x.

It is now convenient to introduce a suitable frequency f_0 , called *reference carrier frequency*, which belongs to the passband (f_1, f_2) of x. The filter output, $x^{(a)}$, is frequency shifted by f_0 to obtain a BB signal, $x^{(bb)}$. The signal $x^{(bb)}$ is the *baseband equivalent* of x, also named *complex envelope of x around the carrier frequency* f_0 .

Analytically, we have

$$x^{(a)}(t) = x * h^{(a)}(t) \quad \stackrel{\mathcal{F}}{\longleftrightarrow} \quad \mathcal{X}^{(a)}(f) = \mathcal{X}(f)\mathcal{H}^{(a)}(f) \tag{1.85}$$

$$x^{(bb)}(t) = x^{(a)}(t) \ e^{-j2\pi f_0 t} \quad \stackrel{r}{\longleftrightarrow} \quad \mathcal{X}^{(bb)}(f) = \mathcal{X}^{(a)}(f+f_0) \tag{1.86}$$

and in the frequency domain

$$\mathcal{X}^{(bb)}(f) = \begin{cases} 2\mathcal{X}(f+f_0) & \text{ for } f > -f_0 \\ 0 & \text{ for } f < -f_0 \end{cases}$$
(1.87)

In other words, $x^{(bb)}$ is given by the components of x at positive frequencies, scaled by 2 and frequency shifted by f_0 .

BB filter. We obtain the same result using a frequency shift of x followed by a lowpass filter (see Figures 1.14 and 1.15). It is immediate to determine the relation between the frequency responses of the filters of Figures 1.12 and 1.14:

$$\mathcal{H}(f) = \mathcal{H}^{(a)}(f + f_0) \tag{1.88}$$

From (1.88) one can derive the relation between the impulse response of the analytic filter and the impulse response of the lowpass filter:

$$h^{(a)}(t) = h(t) \ e^{j2\pi f_0 t} \tag{1.89}$$



Figure 1.14 Illustration of transformations to obtain the baseband equivalent signal $x^{(bb)}$ around the carrier frequency f_0 using a lowpass filter.



Figure 1.15 Transformations to obtain the baseband equivalent signal $x^{(bb)}$ around the carrier frequency f_0 using a lowpass filter.



Figure 1.16 Relation between a signal, its complex envelope and the analytic signal.

Relation between a signal and its complex representation

A simple analytical relation exists between a *real signal x* and its complex envelope. In fact, making use of the property $\mathcal{X}(-f) = \mathcal{X}^*(f)$, it follows

$$\mathcal{X}(f) = \mathcal{X}(f)\mathbf{1}(f) + \mathcal{X}(f)\mathbf{1}(-f) = \mathcal{X}(f)\mathbf{1}(f) + \mathcal{X}^*(-f)\mathbf{1}(-f)$$
(1.90)

or equivalently,

$$x(t) = \frac{x^{(a)}(t) + x^{(a)*}(t)}{2} = Re\left[x^{(a)}(t)\right]$$
(1.91)

Using (1.86) it also follows

$$x(t) = Re\left[x^{(bb)}(t)e^{j2\pi f_0 t}\right]$$
(1.92)

as illustrated in Figure 1.16.

Baseband components of a PB signal. We introduce the notation

$$x^{(bb)}(t) = x_I^{(bb)}(t) + jx_Q^{(bb)}(t)$$
(1.93)

where

$$x_{I}^{(bb)}(t) = Re\left[x^{(bb)}(t)\right]$$
(1.94)

and

$$x_{Q}^{(bb)}(t) = Im[x^{(bb)}(t)]$$
(1.95)

are real-valued baseband signals, named *in-phase* and *quadrature components* of x, respectively. Substituting (1.93) in (1.92), we obtain

$$x(t) = x_I^{(bb)}(t)\cos(2\pi f_0 t) - x_Q^{(bb)}(t)\sin(2\pi f_0 t)$$
(1.96)

as illustrated in Figure 1.17.

Conversely, given x, one can use the scheme of Figure 1.15 and the relations (1.94) and (1.95) to get the baseband components. If the frequency response $\mathcal{H}(f)$ has Hermitian-symmetric characteristics with respect to the origin, h is real and the scheme of Figure 1.18 holds. The scheme of Figure 1.18 employs instead an ideal Hilbert filter with frequency response given by

$$\mathcal{H}^{(h)}(f) = -j \operatorname{sgn}(f) = e^{-j\frac{\pi}{2}\operatorname{sgn}(f)}$$
(1.97)

Magnitude and phase of $\mathcal{H}^{(h)}(f)$ are shown in Figure 1.19. We note that $h^{(h)}$ phase-shifts by $-\pi/2$ the positive-frequency components of the input and by $\pi/2$ the negative-frequency components. In practice, these filter specifications are imposed only on the passband of the input signal.⁹ To simplify the notation, in block diagrams a Hilbert filter is indicated as $-\pi/2$.

$$h^{(h)}(t) = \frac{1}{\pi t}$$
(1.98)

⁹ We note that the ideal Hilbert filter in Figure 1.19 has an impulse response given by (see Table 1.2 on page 5):



Figure 1.17 Relation between a signal and its baseband components.



(a)



(b)

Figure 1.18 Relations to derive the baseband signal components. (a) Implementation using LPF and (b) Implementation using Hilbert filter.



Figure 1.19 Magnitude and phase responses of the ideal Hilbert filter.

Comparing the frequency responses of the analytic filter (1.84) and of the Hilbert filter (1.97), we obtain the relation

$$\mathcal{H}^{(a)}(f) = 1 + j\mathcal{H}^{(h)}(f) \tag{1.101}$$

Then, letting

$$x^{(h)}(t) = x * h^{(h)}(t)$$
(1.102)

the analytic signal can be expressed as

$$x^{(a)}(t) = x(t) + jx^{(h)}(t)$$
(1.103)

Consequently, from (1.86), (1.94), and (1.95), we have

$$x_{I}^{(bb)}(t) = x(t)\cos(2\pi f_{0}t) + x^{(h)}(t)\sin(2\pi f_{0}t)$$
(1.104)

$$x_Q^{(bb)}(t) = x^{(h)}(t)\cos(2\pi f_0 t) - x(t)\sin(2\pi f_0 t)$$
(1.105)

as illustrated in Figure 1.18.10

Consequently, if x is the input signal, the output of the Hilbert filter (also denoted as Hilbert transform of x) is

$$x^{(h)}(t) = \frac{1}{\pi} \int_{-\infty}^{+\infty} \frac{x(\tau)}{t - \tau} d\tau$$
(1.99)

Moreover, noting that from $(1.97) (-j \operatorname{sgn} f)(-j \operatorname{sgn} f) = -1$, taking the Hilbert transform of the Hilbert transform of a signal, we get the initial signal with the sign changed. Then it results as

$$x(t) = -\frac{1}{\pi} \int_{-\infty}^{+\infty} \frac{x^{(h)}(\tau)}{t - \tau} d\tau$$
(1.100)

¹⁰ We recall that the design of a filter, and in particular of a Hilbert filter, requires the introduction of a suitable delay. In other words, we are only able to produce an output with a delay t_D , $x^{(h)}(t - t_D)$. Consequently, in the block diagram of Figure 1.18, also x and the various sinusoidal waveforms must be delayed.

We note that in practical *systems*, transformations to obtain, e.g. the analytic signal, the complex envelope, or the Hilbert transform of a given signal, are implemented by filters. However, it is usually more convenient to perform signal analysis in the frequency domain by the Fourier transform. In the following two examples, we use frequency-domain techniques to obtain the complex envelope of a PB signal.

Example 1.4.1 Consider the sinusoidal signal

$$x(t) = A\cos(2\pi f_0 t + \varphi_0)$$
(1.106)

with

$$\mathcal{X}(f) = \frac{A}{2}e^{j\varphi_0}\delta(f - f_0) + \frac{A}{2}e^{-j\varphi_0}\delta(f + f_0)$$
(1.107)

The analytic signal is given by

$$\mathcal{X}^{(a)}(f) = A e^{j\varphi_0} \delta(f - f_0) \quad \stackrel{\mathcal{F}^{-1}}{\longleftrightarrow} \quad x^{(a)}(t) = A e^{j\varphi_0} e^{j2\pi f_0 t} \tag{1.108}$$

and

$$\mathcal{X}^{(bb)}(f) = Ae^{i\varphi_0}\delta(f) \quad \stackrel{\mathcal{F}^{-1}}{\longleftrightarrow} \quad x^{(bb)}(t) = Ae^{i\varphi_0} \tag{1.109}$$

We note that we have chosen as reference carrier frequency of the complex envelope the same carrier frequency as in (1.106).

Example 1.4.2 Let

$$x(t) = A\sin(Bt)\cos(2\pi f_0 t) \tag{1.110}$$

with the Fourier transform given by

$$\mathcal{X}(f) = \frac{A}{2B} \left[\operatorname{rect}\left(\frac{f - f_0}{B}\right) + \operatorname{rect}\left(\frac{f + f_0}{B}\right) \right]$$
(1.111)

as illustrated in Figure 1.20. Then, using f_0 as reference carrier frequency,

$$\mathcal{X}^{(bb)}(f) = \frac{A}{B} \operatorname{rect}\left(\frac{f}{B}\right)$$
(1.112)

and

$$x^{(bb)}(t) = A\operatorname{sinc}(Bt) \tag{1.113}$$

Another *analytical technique* to get the expression of the signal after the various transformations is obtained by applying the following theorem.

Theorem 1.2

Let the product of two real signals be

$$x(t) = a(t) c(t)$$
 (1.114)

where *a* is a BB signal with $\mathcal{B}_a = [0, B)$ and *c* is a PB signal with $\mathcal{B}_c = [f_0, +\infty)$. If $f_0 > B$, then the analytic signal of *x* is related to that of *c* by

$$x^{(a)}(t) = a(t) c^{(a)}(t)$$
(1.115)

Proof. We consider the general relation (1.91), valid for every real signal

$$c(t) = \frac{1}{2} c^{(a)}(t) + \frac{1}{2} c^{(a)*}(t)$$
(1.116)

Substituting (1.116) in (1.114) yields

$$x(t) = a(t)\frac{1}{2}c^{(a)}(t) + a(t)\frac{1}{2}c^{(a)*}(t)$$
(1.117)

In the frequency domain, the support of the first term in (1.117) is given by the interval $[f_0 - B, +\infty)$, while that of the second is equal to $(-\infty, -f_0 + B]$. Under the hypothesis that $f_0 \ge B$, the two terms in (1.117) have disjoint supports in the frequency domain and (1.115) is immediately obtained.

Corollary 1.1 From (1.115), we obtain

$$x^{(h)}(t) = a(t)c^{(h)}(t)$$
(1.118)

and

$$x^{(bb)}(t) = a(t)c^{(bb)}(t)$$
(1.119)

In fact, from (1.103) we get

$$x^{(h)}(t) = Im[x^{(a)}(t)]$$
(1.120)

which substituted in (1.115) yields (1.118). Finally, (1.119) is obtained by substituting (1.86),

$$x^{(bb)}(t) = x^{(a)}(t)e^{-j2\pi f_0 t}$$
(1.121)

in (1.115).

An interesting application of (1.120) is in the design of a Hilbert filter $h^{(h)}$ starting from a lowpass filter *h*. In fact, from (1.89) and (1.120), we get

$$h^{(h)}(t) = h(t)\sin(2\pi f_0 t) \tag{1.122}$$

Example 1.4.3

Let a modulated double sideband (DSB) signal be expressed as

$$x(t) = a(t)\cos(2\pi f_0 t + \varphi_0)$$
(1.123)



Figure 1.20 Frequency response of a PB signal and corresponding complex envelope.

property	(real) signal	(real) Hilbert transform	
	x(t)	$x^{(h)}(t)$	
duality	$x^{(h)}(t)$	-x(t)	
time inverse	x(-t)	$-x^{(h)}(-t)$	
even signal	x(t) = x(-t)	$x^{(h)}(t) = -x^{(h)}(-t)$, odd	
odd signal	x(t) = -x(-t)	$x^{(h)}(t) = x^{(h)}(-t)$, even	
product (see Theorem 1.2)	a(t) c(t)	$a(t) c^{(h)}(t)$	
cosinusoidal signal	$\cos(2\pi f_0 t + \varphi_0)$	$\sin(2\pi f_0 t + \varphi_0)$	
energy	$E_{x} = \int_{-\infty}^{+\infty} x(t) ^{2} dt = \int_{-\infty}^{+\infty} x^{(h)}(t) ^{2} dt = E_{x^{(h)}}$		
orthogonality	$\int_{-\infty}^{+\infty} x(t) x^{(h)}(t) dt = 0$		

Table 1.5: Some properties of the Hilbert transform.

where *a* is a BB signal with bandwidth *B*. Then, if $f_0 > B$, from the above theorem we have the following relations:

$$x^{(a)}(t) = a(t)e^{j(2\pi f_0 t + \varphi_0)}$$
(1.124)

$$x^{(h)}(t) = a(t)\sin(2\pi f_0 t + \varphi_0)$$
(1.125)

$$x^{(bb)}(t) = a(t)e^{j\varphi_0}$$
(1.126)

We list in Table 1.5 some properties of the Hilbert transformation (1.102) that are easily obtained by using the Fourier transform and the properties of Table 1.1.

Baseband equivalent of a transformation

Given a transformation involving also passband signals, it is often useful to determine an equivalent relation between baseband complex representations of input and output signals. Three transformations are given in Figure 1.21, together with their baseband equivalent. Note that schemes in Figure 1.21a,b produce very different output signals, although both use a mixer with the same carrier.

We will prove the relation illustrated in Figure 1.21b. Assuming that h is the real-valued impulse response of an LPF and using (1.92),

$$y(t) = \left\{ h * Re \left[x^{(bb)}(\tau) e^{j2\pi f_0 \tau} \left(\cos(2\pi f_0 \tau + \varphi_1) \right) \right] \right\} (t)$$

= $Re \left[\left(h * x^{(bb)} \frac{e^{-j\varphi_1}}{2} + h * x^{(bb)} \frac{e^{+j(2\pi 2 f_0 \tau + \varphi_1)}}{2} \right) (t) \right]$ (1.127)
= $Re \left[\left(h * x^{(bb)} \frac{e^{-j\varphi_1}}{2} \right) (t) \right]$

where the last equality follows because the term with frequency components around $2f_0$ is filtered by the LPF.





Figure 1.21 Passband transformations and their baseband equivalent. (a) Modulator, (b) demodulator, and (c) passband filtering.

We note, moreover, that the filter $h^{(bb)}$ in Figure 1.21 has in-phase component $h_I^{(bb)}$ and quadrature component $h_Q^{(bb)}$ that are related to $\mathcal{H}^{(a)}$ by (see (1.94) and (1.95))

$$\mathcal{H}_{I}^{(bb)}(f) = \frac{1}{2} [\mathcal{H}^{(bb)}(f) + \mathcal{H}^{(bb)*}(-f)]$$

= $\frac{1}{2} [\mathcal{H}^{(a)}(f + f_{0}) + \mathcal{H}^{(a)*}(-f + f_{0})]$ (1.128)

and

$$\mathcal{H}_{Q}^{(bb)}(f) = \frac{1}{2j} [\mathcal{H}^{(bb)}(f) - \mathcal{H}^{(bb)*}(-f)]$$

= $\frac{1}{2j} [\mathcal{H}^{(a)}(f+f_0) - \mathcal{H}^{(a)*}(-f+f_0)]$ (1.129)

Consequently, if $\mathcal{H}^{(a)}$ has Hermitian symmetry around f_0 , then

$$\mathcal{H}_I^{(bb)}(f) = \mathcal{H}_a^{(a)}(f + f_0)$$

and

$$\mathcal{H}_Q^{(bb)}(f) = 0$$

In other words, $h^{(bb)}(t) = h_I^{(bb)}(t)$ is real and the realization of the filter $\frac{1}{2}h^{(bb)}$ is simplified. In practice, this condition is verified by imposing that the filter $h^{(a)}$ has symmetrical frequency specifications around f_0 .

Envelope and instantaneous phase and frequency

We will conclude this section with a few definitions. Given a PB signal *x*, with reference to the analytic signal we define

1. Envelope

$$M_{\rm r}(t) = |x^{(a)}(t)| \tag{1.130}$$

2. Instantaneous phase

$$\varphi_x(t) = \arg x^{(a)}(t) \tag{1.131}$$

3. Instantaneous frequency

$$f_x(t) = \frac{1}{2\pi} \frac{d}{dt} \varphi_x(t) \tag{1.132}$$

In terms of the complex envelope signal $x^{(bb)}$, from (1.86) the equivalent relations follow:

$$M_x(t) = |x^{(bb)}(t)| \tag{1.133}$$

$$\varphi_x(t) = \arg x^{(bb)}(t) + 2\pi f_0 t \tag{1.134}$$

$$f_x(t) = \frac{1}{2\pi} \frac{d}{dt} [\arg x^{(bb)}(t)] + f_0$$
(1.135)

Then, from the polar representation, $x^{(a)}(t) = M_x(t) e^{j\varphi_x(t)}$ and from (1.91), a PB signal x can be written as

$$x(t) = Re[x^{(a)}(t)] = M_x(t)\cos(\varphi_x(t))$$
(1.136)

For example if $x(t) = A \cos(2\pi f_0 t + \varphi_0)$, it follows that

$$M_x(t) = A \tag{1.137}$$

$$\varphi_x(t) = 2\pi f_0 t + \varphi_0 \tag{1.138}$$

$$f_x(t) = f_0 \tag{1.139}$$

With reference to the above relations, three other definitions follow.

1. Envelope deviation

$$\Delta M_x(t) = |x^{(a)}(t)| - A = |x^{(bb)}(t)| - A \tag{1.140}$$

2. Phase deviation

$$\Delta \varphi_x(t) = \varphi_x(t) - (2\pi f_0 t + \varphi_0) = \arg x^{(bb)}(t) - \varphi_0$$
(1.141)

3. Frequency deviation

$$\Delta f_x(t) = f_x(t) - f_0 = \frac{1}{2\pi} \frac{d}{dt} \Delta \varphi_x(t)$$
(1.142)

Then (1.136) becomes

$$x(t) = [A + \Delta M_x(t)] \cos(2\pi f_0 t + \varphi_0 + \Delta \varphi_x(t))$$
(1.143)

1.5 Second-order analysis of random processes

We recall the functions related to the statistical description of random processes, especially those functions concerning second-order analysis.

1.5.1 Correlation

Let x(t) and y(t), $t \in \mathbb{R}$, be two continuous-time complex-valued random processes. We indicate the expectation operator with *E*.

1. Mean value

$$m_x(t) = E[x(t)]$$
 (1.144)

2. Statistical power

$$M_{x}(t) = E[|x(t)|^{2}]$$
(1.145)

3. Autocorrelation

$$\mathbf{r}_{x}(t, t-\tau) = E[x(t)x^{*}(t-\tau)]$$
(1.146)

4. Crosscorrelation

$$\mathbf{r}_{xy}(t, t-\tau) = E[x(t)y^*(t-\tau)]$$
(1.147)

5. Autocovariance

$$c_{x}(t, t - \tau) = E[(x(t) - m_{x}(t))(x(t - \tau) - m_{x}(t - \tau))^{*}]$$

= $r_{y}(t, t - \tau) - m_{y}(t)m_{y}^{*}(t - \tau)$ (1.148)

6. Crosscovariance

$$c_{xy}(t, t - \tau) = E[(x(t) - m_x(t))(y(t - \tau) - m_y(t - \tau))^*]$$

= $r_{xy}(t, t - \tau) - m_x(t)m_y^*(t - \tau)$ (1.149)

Observation 1.2

- x and y are orthogonal if $r_{xy}(t, t \tau) = 0$, $\forall t, \tau$. In this case, we write $x \perp y$.¹¹
- x and y are uncorrelated if $c_{xy}(t, t \tau) = 0, \forall t, \tau$.
- if at least one of the two random processes has zero mean, *orthogonality* is equivalent to *uncorrelation*.
- x is wide-sense stationary (WSS) if
 - 1. $\mathbf{m}_{x}(t) = \mathbf{m}_{x}, \forall t,$
 - 2. $\mathbf{r}_x(t, t \tau) = \mathbf{r}_x(\tau), \forall t$.

In this case, $r_x(0) = E[|x(t)|^2] = M_x$ is the *statistical power*, whereas $c_x(0) = \sigma_x^2 = M_x - |m_x|^2$ is the *variance* of *x*.

¹¹ We observe that the notion of *orthogonality between two random processes* is quite different from that of *orthogonality between two deterministic signals*. In fact, while in the deterministic case, it is sufficient that $\int_{-\infty}^{\infty} x(t)y^*(t)dt = 0$, in the random case, the crosscorrelation must be zero for all the delays and not only for the zero delay. In particular, we note that the *two random variables* v_1 and v_2 are orthogonal if $E[v_1v_2^*] = 0$.

- x and y are *jointly wide-sense stationary* if 1. $m_x(t) = m_x, m_v(t) = m_v, \forall t$,
 - 2. $\mathbf{r}_{xy}(t, t \tau) = \mathbf{r}_{xy}(\tau), \forall t.$

Properties of the autocorrelation function

- 1. $r_x(-\tau) = r_x^*(\tau), r_x(\tau)$ is a function with Hermitian symmetry. 2. $r_x(0) \ge |r_x(\tau)|$. 3. $r_x(0)r_y(0) \ge |r_{xy}(\tau)|^2$. 4. $r_{xy}(-\tau) = r_{yx}^*(\tau)$.
- 5. $r_{x^*}(\tau) = r_x^*(\tau)$.

1.5.2 Power spectral density

Given the WSS random process x(t), $t \in \mathbb{R}$, its *power spectral density* (PSD) is defined as the Fourier transform of its autocorrelation function

$$\mathcal{P}_{x}(f) = \mathcal{F}[\mathbf{r}_{x}(\tau)] = \int_{-\infty}^{+\infty} \mathbf{r}_{x}(\tau) e^{-j2\pi f\tau} d\tau \qquad (1.150)$$

The inverse transformation is given by the following formula:

$$\mathbf{r}_{x}(\tau) = \int_{-\infty}^{+\infty} \mathcal{P}_{x}(f) e^{j2\pi f\tau} df \qquad (1.151)$$

In particular from (1.151), we obtain the statistical power

$$\mathbb{M}_{x} = \mathbf{r}_{x}(0) = \int_{-\infty}^{+\infty} \mathcal{P}_{x}(f) df \qquad (1.152)$$

Hence, the name PSD for the function $\mathcal{P}_x(f)$: it represents the distribution of the statistical power in the frequency domain.

The pair of equations (1.150) and (1.151) are obtained from the Wiener-Khintchine theorem [3].

Definition 1.8

The passband \mathcal{B} of a random process x is defined with reference to its PSD function.

Spectral lines in the PSD

In many applications, it is important to detect the presence of sinusoidal components in a random process. With this intent we give the following theorem.

Theorem 1.3

The PSD of a WSS process, $\mathcal{P}_{x}(f)$, can be uniquely decomposed into a component $\mathcal{P}_{x}^{(c)}(f)$ without delta functions and a discrete component consisting of delta functions (spectral lines) $\mathcal{P}_{x}^{(d)}(f)$, so that

$$\mathcal{P}_{x}(f) = \mathcal{P}_{x}^{(c)}(f) + \mathcal{P}_{x}^{(d)}(f)$$
(1.153)

where $\mathcal{P}_{x}^{(c)}(f)$ is an ordinary (piecewise linear) function and

$$\mathcal{P}_x^{(d)}(f) = \sum_{i \in \mathcal{I}} \mathbb{M}_i \delta(f - f_i)$$
(1.154)

where \mathcal{I} identifies a discrete set of frequencies $\{f_i\}, i \in \mathcal{I}$.

The inverse Fourier transform of (1.153) yields the relation

$$\mathbf{r}_{x}(\tau) = \mathbf{r}_{x}^{(c)}(\tau) + \mathbf{r}_{x}^{(d)}(\tau)$$
(1.155)

with

$$\mathbf{r}_{x}^{(d)}(\tau) = \sum_{i \in \mathcal{I}} \mathbb{M}_{i} e^{j2\pi f_{i}\tau}$$
(1.156)

The most interesting consideration is that the following random process decomposition corresponds to the decomposition (1.153) of the PSD:

$$x(t) = x^{(c)}(t) + x^{(d)}(t)$$
(1.157)

where $x^{(c)}$ and $x^{(d)}$ are *orthogonal* processes having PSD functions

$$\mathcal{P}_{\chi^{(c)}}(f) = \mathcal{P}_{\chi}^{(c)}(f) \quad \text{and} \quad \mathcal{P}_{\chi^{(d)}}(f) = \mathcal{P}_{\chi}^{(d)}(f)$$
(1.158)

Moreover, $x^{(d)}$ is given by

$$x^{(d)}(t) = \sum_{i \in \mathcal{I}} x_i e^{j2\pi f_i t}$$
(1.159)

where $\{x_i\}$ are orthogonal random variables (r.v.s.) having statistical power

$$E[|x_i|^2] = \mathbb{M}_i, \quad i \in \mathcal{I}$$
(1.160)

where M_i is defined in (1.154).

Observation 1.3 The spectral lines of the PSD identify the periodic components in the process.

Definition 1.9

A WSS random process is said to be asymptotically uncorrelated if the following two properties hold:

(1)
$$\lim_{\tau \to \infty} r_x(\tau) = |\mathbf{m}_x|^2$$
 (1.161)

(2)
$$c_x(\tau) = r_x(\tau) - |m_x|^2$$
 is absolutely integrable (1.162)

The property (1) shows that x(t) and $x(t - \tau)$ become uncorrelated for $\tau \to \infty$.

For such processes, one can prove that

$$\mathbf{r}_{x}^{(c)}(\tau) = \mathbf{c}_{x}(\tau) \quad \text{and} \quad \mathbf{r}_{x}^{(d)}(\tau) = |\mathbf{m}_{x}|^{2}$$
(1.163)

Hence, $\mathcal{P}_{x}^{(d)}(f) = |\mathbf{m}_{x}|^{2} \delta(f)$, and the process exhibits at most a spectral line at the origin.

Cross power spectral density

One can extend the definition of PSD to two jointly WSS random processes:

$$\mathcal{P}_{xv}(f) = \mathcal{F}[\mathbf{r}_{xv}(\tau)] \tag{1.164}$$

Since $r_{xy}(-\tau) \neq r_{xy}^*(\tau)$, $\mathcal{P}_{xy}(f)$ is in general a complex function.

Properties of the PSD

- 1. $\mathcal{P}_{x}(f)$ is a *real-valued function*. This follows from property 1 of the autocorrelation.
- 2. $\mathcal{P}_x(f)$ is generally *not an even function*. However, if the process x is real valued, then both $r_x(\tau)$ and $\mathcal{P}_x(f)$ are even functions.
- 3. $\mathcal{P}_{x}(f)$ is a non-negative function.
- 4. $\mathcal{P}_{yx}(f) = \mathcal{P}^*_{xy}(f).$

5.
$$\mathcal{P}_{x^*}(f) = \mathcal{P}_x(-f)$$

Moreover, the following inequality holds:

$$0 \le |\mathcal{P}_{xy}(f)|^2 \le \mathcal{P}_{x}(f)\mathcal{P}_{y}(f) \tag{1.165}$$

Definition 1.10 (White random process)

The zero-mean random process x(t), $t \in \mathbb{R}$, is called *white* if

$$\mathbf{r}_{\mathbf{y}}(\tau) = K\delta(\tau) \tag{1.166}$$

with K a positive real number. In this case, $\mathcal{P}_{x}(f)$ is a constant, i.e.

$$\mathcal{P}_x(f) = K \tag{1.167}$$

PSD through filtering

With reference to Figure 1.22, by taking the Fourier transform of the various crosscorrelations, the following relations are easily obtained:

$$\mathcal{P}_{yx}(f) = \mathcal{P}_{x}(f)\mathcal{H}(f) \tag{1.168}$$

$$\mathcal{P}_{v}(f) = \mathcal{P}_{r}(f)|\mathcal{H}(f)|^{2} \tag{1.169}$$

$$\mathcal{P}_{v_{\mathcal{I}}}(f) = \mathcal{P}_{x}(f)\mathcal{H}(f)\mathcal{G}^{*}(f)$$
(1.170)

The relation (1.169) is of particular interest since it relates the PSDs of the output process of a filter to the PSD of the input process, through the frequency response of the filter. In the particular case in which y and z have disjoint passbands, i.e. $\mathcal{P}_{y}(f)\mathcal{P}_{z}(f) = 0$, then from (1.165) $r_{yz}(\tau) = 0$, and $y \perp z$.

1.5.3 PSD of discrete-time random processes

Let $\{x(k)\}$ and $\{y(k)\}$ be two discrete-time random processes. Definitions and properties of Section 1.5.1 remain valid also for discrete-time processes: the only difference is that the correlation is now defined



Figure 1.22 Reference scheme of PSD computations.

on discrete time and is called autocorrelation sequence (ACS). It is however interesting to review the properties of PSDs. Given a discrete-time WSS random process *x*, the PSD is obtained as

$$\mathcal{P}_{x}(f) = T_{c}\mathcal{F}[\mathbf{r}_{x}(n)] = T_{c}\sum_{n=-\infty}^{+\infty}\mathbf{r}_{x}(n)e^{-j2\pi f n T_{c}}$$
(1.171)

We note a further property: $\mathcal{P}_{x}(f)$ is a periodic function of period $1/T_{c}$. The inverse transformation yields:

$$\mathbf{r}_{x}(n) = \int_{-\frac{1}{2T_{c}}}^{\frac{1}{2T_{c}}} \mathcal{P}_{x}(f) e^{j2\pi f n T_{c}} df$$
(1.172)

In particular, the statistical power is given by

$$M_{x} = \mathbf{r}_{x}(0) = \int_{-\frac{1}{2T_{c}}}^{\frac{1}{2T_{c}}} \mathcal{P}_{x}(f) df$$
(1.173)

Definition 1.11 (White random process) A discrete-time random process $\{x(k)\}$ is white if

$$\mathbf{r}_x(n) = \sigma_x^2 \delta_n \tag{1.174}$$

In this case, the PSD is a constant:

$$\mathcal{P}_x(f) = \sigma_x^2 T_c \tag{1.175}$$

Definition 1.12

If the samples of the random process $\{x(k)\}$ are statistically independent and identically Distributed, we say that $\{x(k)\}$ has i.i.d. samples.

Generating an i.i.d. sequence is not simple; however, it is easily provided by many random number generators [4]. However, generating, storing, and processing a finite length, i.i.d. sequence requires a complex processor and a lot of memory. Furthermore, the deterministic correlation properties of such a subsequence may not be very good. Hence, in Appendix 1.C we introduce a class of pseudonoise (PN) sequences, which are deterministic and periodic, with very good correlation properties. Moreover, the symbol alphabet can be just binary.

Spectral lines in the PSD

Also the PSD of a discrete time random process can be decomposed into ordinary components and spectral lines on a period of the PSD. In particular for a discrete-time WSS *asymptotically uncorrelated* random process, the relation (1.163) and the following are true

$$\mathcal{P}_{x}^{(c)}(f) = T_{c} \sum_{n=-\infty}^{+\infty} c_{x}(n) \ e^{-j2\pi f n T_{c}}$$
(1.176)

$$\mathcal{P}_{x}^{(d)}(f) = |\mathbf{m}_{x}|^{2} \sum_{\ell=-\infty}^{+\infty} \delta\left(f - \frac{\ell}{T_{c}}\right)$$
(1.177)

We note that, if the process has non-zero mean value, the PSD exhibits lines at multiples of $1/T_c$.

Example 1.5.1

We calculate the PSD of an i.i.d. sequence $\{x(k)\}$. From

$$\mathbf{r}_{x}(n) = \begin{cases} M_{x} & n = 0\\ |\mathbf{m}_{x}|^{2} & n \neq 0 \end{cases}$$
(1.178)

it follows that

$$c_x(n) = \begin{cases} \sigma_x^2 & n = 0\\ 0 & n \neq 0 \end{cases}$$
(1.179)

Then

$$\mathbf{r}_{x}^{(c)}(n) = \sigma_{x}^{2} \delta_{n}, \quad \mathbf{r}_{x}^{(d)}(n) = |\mathbf{m}_{x}|^{2}$$
(1.180)

$$\mathcal{P}_x^{(c)}(f) = \sigma_x^2 T_c, \quad \mathcal{P}_x^{(d)}(f) = |\mathbf{m}_x|^2 \sum_{\ell=-\infty}^{+\infty} \delta\left(f - \frac{\ell}{T_c}\right)$$
(1.181)

PSD through filtering

Given the system illustrated in Figure 1.3, we want to find the relation between the PSDs of the input and output signals, assuming these processes are individually as well as jointly WSS. We introduce the z-transform of the correlation sequence:

$$P_{x}(z) = \sum_{n=-\infty}^{+\infty} \mathbf{r}_{x}(n) z^{-n}$$
(1.182)

From the comparison of (1.182) with (1.171), the PSD of x is related to $P_x(z)$ by

$$\mathcal{P}_x(f) = T_c P_x(e^{j2\pi f T_c}) \tag{1.183}$$

Using Table 1.3 in page 6, we obtain the relations between ACS and PSD listed in Table 1.6. Let the deterministic autocorrelation of h be defined as¹²

$$\mathbf{r}_{h}(n) = \sum_{k=-\infty}^{+\infty} h(k)h^{*}(k-n) = [h(m) * h^{*}(-m)](n)$$
(1.184)

whose z-transform is given by

$$P_h(z) = \sum_{n=-\infty}^{+\infty} \mathbf{r}_h(n) \ z^{-n} = H(z) \ H^*\left(\frac{1}{z^*}\right)$$
(1.185)

Table 1.6: Relations between ACS and PSD for discrete-time processes through a linear filter.

ACS	PSD
$\mathbf{r}_{yx}(n) = \mathbf{r}_x * h(n)$	$P_{yx}(z) = P_x(z)H(z)$
$r_{xy}(n) = [r_x(m) * h^*(-m)](n)$	$P_{xy}(z) = P_x(z)H^*(1/z^*)$
$\mathbf{r}_{y}(n) = \mathbf{r}_{xy} * h(n)$	$P_{y}(z) = P_{xy}(z)H(z)$
$= \mathbf{r}_x * \mathbf{r}_h(n)$	$= P_x(z)H(z)H^*(1/z^*)$

¹² We use the same symbol to indicate the correlation between random processes and the correlation between deterministic signals.

In case $P_h(z)$ is a rational function, from (1.185) one deduces that, if $P_h(z)$ has a pole (zero) of the type $e^{j\varphi}|a|$, it also has a corresponding pole (zero) of the type $e^{j\varphi}/|a|$. Consequently, the poles (and zeros) of $P_h(z)$ come in pairs of the type $e^{j\varphi}|a|$, $e^{j\varphi}/|a|$.

From the last relation in Table 1.6, we obtain the relation between the PSDs of input and output signals, i.e.

$$\mathcal{P}_{y}(f) = \mathcal{P}_{x}(f) \left| H(e^{j2\pi f T_{c}}) \right|^{2}$$
(1.186)

In the case of white noise input

$$P_{y}(z) = \sigma_{x}^{2} H(z) H^{*}\left(\frac{1}{z^{*}}\right)$$
(1.187)

and

$$\mathcal{P}_{y}(f) = T_{c}\sigma_{x}^{2} \left| H(e^{j2\pi fT_{c}}) \right|^{2}$$
(1.188)

In other words, $\mathcal{P}_{v}(f)$ has the same shape as the filter frequency response.

In the case of real filters

$$H^*\left(\frac{1}{z^*}\right) = H(z^{-1})$$
(1.189)

Among the various applications of (1.188), it is worth mentioning the *process synthesis*, which deals with the generation of a random process having a pre-assigned PSD. Two methods are shown in Section 4.1.9.

Minimum-phase spectral factorization

In the previous section, we introduced the relation between an impulse response $\{h(k)\}$ and its ACS $\{r_h(n)\}$ in terms of the z-transform. In many practical applications, it is interesting to determine the minimum-phase impulse response for a given autocorrelation function: with this intent we state the following theorem [5].

Theorem 1.4 (Spectral factorization for discrete-time processes)

Consider the process y with ACS $\{r_y(n)\}$ having z-transform $P_y(z)$, which satisfies the Paley–Wiener condition for discrete-time systems, i.e.

$$\int_{1/T_c} \left| \ln P_y(e^{j2\pi fT_c}) \right| df < \infty$$
(1.190)

where the integration is over an arbitrarily chosen interval $1/T_c$. Then the function $P_y(z)$ can be factorized as follows:

$$P_{y}(z) = f_{0}^{2} \tilde{F}(z) \tilde{F}^{*}\left(\frac{1}{z^{*}}\right)$$
(1.191)

where

$$\tilde{F}(z) = 1 + \tilde{f}_1 z^{-1} + \tilde{f}_2 z^{-2} + \cdots$$
(1.192)

is monic, minimum phase, and associated with a causal sequence $\{1, \tilde{f}_1, \tilde{f}_2, ...\}$. The factor f_0 in (1.191) is the geometric mean of $P_y(e^{j2\pi f T_c})$:

$$\ln f_0^2 = T_c \int_{1/T_c} \ln P_y(e^{i2\pi f T_c}) df$$
(1.193)

The logarithms in (1.190) and (1.193) may have any common base.

The Paley–Wiener criterion implies that $P_y(z)$ may have only a discrete set of zeros on the unit circle, and that the spectral factorization (1.191) (with the constraint that $\tilde{F}(z)$ is causal, monic and minimum

phase) is unique. For rational $P_y(z)$, the function $f_0\tilde{F}(z)$ is obtained by extracting the poles and zeros of $P_y(z)$ that lie inside the unit circle (see (1.453) and the considerations relative to (1.185)). Moreover, in (1.191) $f_0\tilde{F}^*(1/z^*)$ is the z-transform of an anticausal sequence $f_0\{\ldots, \tilde{f}_2^*, \tilde{f}_1^*, 1\}$, associated with poles and zeros of $P_y(z)$ that lie outside the unit circle.

1.5.4 PSD of passband processes

Definition 1.13

A WSS random process x is said to be PB (BB) if its PSD is of PB (BB) type.

PSD of in-phase and quadrature components

Let x be a real PB WSS process. Our aim is to derive the PSD of the in-phase and quadrature components of the process. We assume that x does not have direct current (DC) components, i.e. a frequency component at f = 0, hence, its mean is zero and consequently also $x^{(a)}$ and $x^{(bb)}$ have zero mean.

We introduce the two (ideal) filters with frequency response

$$\mathcal{H}^{(+)}(f) = 1(f) \quad \text{and} \quad \mathcal{H}^{(-)}(f) = 1(-f)$$
 (1.194)

Note that they have non-overlapping passbands. For the same input *x*, the output of the two filters is, respectively, $x^{(+)}$ and $x^{(-)}$. We find that

$$x(t) = x^{(+)}(t) + x^{(-)}(t)$$
(1.195)

with $x^{(-)}(t) = x^{(+)*}(t)$. The following relations hold

$$\mathcal{P}_{r^{(+)}}(f) = |\mathcal{H}^{(+)}(f)|^2 \mathcal{P}_r(f) = \mathcal{P}_r(f) \mathbf{1}(f)$$
(1.196)

$$\mathcal{P}_{\chi^{(-)}}(f) = |\mathcal{H}^{(-)}(f)|^2 \mathcal{P}_{\chi}(f) = \mathcal{P}_{\chi}(f) 1(-f)$$
(1.197)

and

$$\mathcal{P}_{\chi^{(+)}\chi^{(-)}}(f) = 0 \tag{1.198}$$

as $x^{(+)}$ and $x^{(-)}$ have non-overlapping passbands. Then $x^{(+)} \perp x^{(-)}$, and (1.195) yields

$$\mathcal{P}_{x}(f) = \mathcal{P}_{x^{(+)}}(f) + \mathcal{P}_{x^{(-)}}(f)$$
(1.199)

where $\mathcal{P}_{x^{(-)}}(f) = \mathcal{P}_{x^{(+)*}}(f) = \mathcal{P}_{x^{(+)}}(-f)$, using Property 5 of the PSD. The analytic signal $x^{(a)}$ is equal to $2x^{(+)}$, hence,

$$\mathbf{r}_{\chi^{(a)}}(\tau) = 4\mathbf{r}_{\chi^{(+)}}(\tau) \tag{1.200}$$

and

$$\mathcal{P}_{x^{(a)}}(f) = 4\mathcal{P}_{x^{(+)}}(f) \tag{1.201}$$

Moreover, being $x^{(a)*} = 2x^{(-)}$, it follows that $x^{(a)} \perp x^{(a)*}$ and

$$\mathbf{r}_{\mathbf{x}^{(a)}\mathbf{x}^{(a)*}}(\tau) = 0 \tag{1.202}$$

The complex envelope $x^{(bb)}$ is related to $x^{(a)}$ by (1.86) and

$$\mathbf{r}_{x^{(bb)}}(\tau) = \mathbf{r}_{x^{(a)}}(\tau)e^{-j2\pi f_0 \tau}$$
(1.203)

Hence,

$$\mathcal{P}_{x^{(bb)}}(f) = \mathcal{P}_{x^{(a)}}(f+f_0) = 4\mathcal{P}_{x^{(+)}}(f+f_0) \tag{1.204}$$

Moreover, from (1.202), it follows that $x^{(bb)} \perp x^{(bb)*}$.

Using (1.204), (1.199) can be written as

$$\mathcal{P}_{x}(f) = \frac{1}{4} [\mathcal{P}_{x^{(bb)}}(f - f_0) + \mathcal{P}_{x^{(bb)}}(-f - f_0)]$$
(1.205)

Finally, from

$$x_{I}^{(bb)}(t) = Re\left[x^{(bb)}(t)\right] = \frac{x^{(bb)}(t) + x^{(bb)*}(t)}{2}$$
(1.206)

and

$$x_{Q}^{(bb)}(t) = Im[x^{(bb)}(t)] = \frac{x^{(bb)}(t) - x^{(bb)*}(t)}{2j}$$
(1.207)

we obtain the following relations:

$$\mathbf{r}_{\chi_{I}^{(bb)}}(\tau) = \frac{1}{2} Re\left[\mathbf{r}_{\chi^{(bb)}}(\tau)\right]$$
(1.208)

$$\mathcal{P}_{x_{I}^{(bb)}}(f) = \frac{1}{4} [\mathcal{P}_{x^{(bb)}}(f) + \mathcal{P}_{x^{(bb)}}(-f)]$$
(1.209)

$$\mathbf{r}_{x_{Q}^{(bb)}}(\tau) = \mathbf{r}_{x_{I}^{(bb)}}(\tau) \tag{1.210}$$

$$\mathbf{r}_{x_{Q}^{(bb)}x_{I}^{(bb)}}(\tau) = \frac{1}{2}Im[\mathbf{r}_{x^{(bb)}}(\tau)]$$
(1.211)

$$\mathcal{P}_{x_{Q}^{(bb)}x_{I}^{(bb)}}(f) = \frac{1}{4j} [\mathcal{P}_{x^{(bb)}}(f) - \mathcal{P}_{x^{(bb)}}(-f)]$$
(1.212)

$$\mathbf{r}_{x_{l}^{(bb)}x_{Q}^{(bb)}}(\tau) = -\mathbf{r}_{x_{Q}^{(bb)}x_{l}^{(bb)}}(\tau) = -\mathbf{r}_{x_{l}^{(bb)}x_{Q}^{(bb)}}(-\tau)$$
(1.213)

The second equality in (1.213) follows from Property 4 of ACS.

From (1.213), we note that $r_{x_{l}^{(bb)}x_{Q}^{(bb)}}(\tau)$ is an odd function. Moreover, from (1.212), we obtain $x_{l}^{(bb)} \perp x_{Q}^{(bb)}$ only if $\mathcal{P}_{x^{(bb)}}$ is an even function; in any case, the random variables $x_{l}^{(bb)}(t)$ and $x_{Q}^{(bb)}(t)$ are always orthogonal since $r_{x_{l}^{(bb)}x_{Q}^{(bb)}}(0) = 0$. Referring to the block diagram in Figure 1.18b, as

$$\mathcal{P}_{\chi^{(h)}}(f) = \mathcal{P}_{\chi}(f) \quad \text{and} \quad \mathcal{P}_{\chi^{(h)}\chi}(f) = -j \operatorname{sgn}(f) \ \mathcal{P}_{\chi}(f)$$
(1.214)

we obtain

$$\mathbf{r}_{x^{(h)}}(\tau) = \mathbf{r}_{x}(\tau) \quad \text{and} \quad \mathbf{r}_{x^{(h)}x}(\tau) = \mathbf{r}_{x}^{(h)}(\tau)$$
 (1.215)

Then

$$\mathbf{r}_{x_{l}^{(bb)}}(\tau) = \mathbf{r}_{x_{Q}^{(bb)}}(\tau) = \mathbf{r}_{x}(\tau)\cos(2\pi f_{0}\tau) + \mathbf{r}_{x}^{(h)}(\tau)\sin(2\pi f_{0}\tau)$$
(1.216)

and

$$\mathbf{r}_{x_{l}^{(bb)}x_{Q}^{(bb)}}(\tau) = -\mathbf{r}_{x}^{(h)}(\tau)\cos(2\pi f_{0}\tau) + \mathbf{r}_{x}(\tau)\sin(2\pi f_{0}\tau)$$
(1.217)

In terms of statistical power, the following relations hold:

$$\mathbf{r}_{x^{(+)}}(0) = \mathbf{r}_{x^{(-)}}(0) = \frac{1}{2}\mathbf{r}_{x}(0)$$
(1.218)

$$\mathbf{r}_{\mathbf{x}^{(bb)}}(0) = \mathbf{r}_{\mathbf{x}^{(a)}}(0) = 4\mathbf{r}_{\mathbf{x}^{(+)}}(0) = 2\mathbf{r}_{\mathbf{x}}(0)$$
(1.219)

$$\mathbf{r}_{x_{l}^{(bb)}}(0) = \mathbf{r}_{x_{Q}^{(bb)}}(0) = \mathbf{r}_{x}(0)$$
(1.220)

$$\mathbf{r}_{x^{(h)}}(0) = \mathbf{r}_{x}(0) \tag{1.221}$$

Example 1.5.2 Let *x* be a WSS process with PSD

$$\mathcal{P}_{x}(f) = \frac{N_{0}}{2} \left[\operatorname{rect}\left(\frac{f - f_{0}}{B}\right) + \operatorname{rect}\left(\frac{f + f_{0}}{B}\right) \right]$$
(1.222)



Figure 1.23 Spectral representation of a PB process and its BB components.

depicted in Figure 1.23. It is immediate to get

$$\mathcal{P}_{\chi^{(a)}}(f) = 2N_0 \operatorname{rect}\left(\frac{f-f_0}{B}\right)$$
(1.223)

and

$$\mathcal{P}_{\chi^{(bb)}}(f) = 2N_0 \operatorname{rect}\left(\frac{f}{B}\right)$$
(1.224)

Then

$$\mathcal{P}_{x_{l}^{(bb)}}(f) = \mathcal{P}_{x_{Q}^{(bb)}}(f) = \frac{1}{2}\mathcal{P}_{x^{(bb)}}(f) = N_{0}\operatorname{rect}\left(\frac{f}{B}\right)$$
(1.225)

Moreover, being $\mathcal{P}_{x_{l}^{(bb)}x_{Q}^{(bb)}}(f) = 0$, we have that $x_{l}^{(bb)} \perp x_{Q}^{(bb)}$.

Cyclostationary processes

We have seen that, if x is a real passband WSS process, then its complex envelope is WSS, and $x^{(bb)} \perp x^{(bb)*}$. The converse is also true: if $x^{(bb)}$ is a WSS process and $x^{(bb)} \perp x^{(bb)*}$, then

$$x(t) = Re \left[x^{(bb)}(t) \ e^{j2\pi f_0 t} \right]$$
(1.226)

is WSS with PSD given by (1.205). If $x^{(bb)}$ is WSS, however, with

$$r_{x^{(bb)}x^{(bb)*}}(\tau) \neq 0$$
 (1.227)

observing (1.226) we find that the autocorrelation of x is a periodic function in t of period $1/f_0$:

$$\mathbf{r}_{x}(t,t-\tau) = \frac{1}{4} \left[\mathbf{r}_{x^{(bb)}}(\tau) e^{j2\pi f_{0}\tau} + \mathbf{r}_{x^{(bb)}}^{*}(\tau) e^{-j2\pi f_{0}\tau} + \mathbf{r}_{x^{(bb)}x^{(bb)*}}(\tau) e^{-j2\pi f_{0}\tau} e^{j4\pi f_{0}t} + \mathbf{r}_{x^{(bb)}x^{(bb)*}}^{*}(\tau) e^{j2\pi f_{0}\tau} e^{-j4\pi f_{0}t} \right]$$
(1.228)

In other words, x is a cyclostationary process of period $T_0 = 1/f_0$.¹³

In this case, it is convenient to introduce the average correlation

$$\overline{\mathbf{r}}_{x}(\tau) = \frac{1}{T_{0}} \int_{0}^{T_{0}} \mathbf{r}_{x}(t, t - \tau) dt$$
(1.229)

whose Fourier transform is the average PSD

$$\overline{\mathcal{P}}_{x}(f) = \mathcal{F}[\overline{\mathbf{r}}_{x}(\tau)] = \frac{1}{T_{0}} \int_{0}^{T_{0}} \mathcal{P}_{x}(f, t) dt$$
(1.230)

where

$$\mathcal{P}_{x}(f,t) = \mathcal{F}_{\tau}[\mathbf{r}_{x}(t,t-\tau)]$$
(1.231)

In (1.231), \mathcal{F}_{τ} denotes the Fourier transform with respect to the variable τ . In our case, it is

$$\overline{\mathcal{P}}_{x}(f) = \frac{1}{4} \left[\mathcal{P}_{x^{(bb)}}(f - f_0) + \mathcal{P}_{x^{(bb)}}(-f - f_0) \right]$$
(1.232)

as in the stationary case (1.205).

Example 1.5.3

Let x be a modulated DSB signal (see (1.123)), i.e.

$$x(t) = a(t)\cos(2\pi f_0 t + \varphi_0)$$
(1.233)

with *a* real random BB WSS process with bandwidth $B_a < f_0$ and autocorrelation $r_a(\tau)$. From (1.126) it results $x^{(bb)}(t) = a(t) e^{i\varphi_0}$. Hence, we have

$$\mathbf{r}_{x^{(bb)}}(\tau) = \mathbf{r}_{a}(\tau), \quad \mathbf{r}_{x^{(bb)}x^{(bb)*}}(\tau) = \mathbf{r}_{a}(\tau) \ e^{j2\varphi_{0}}$$
 (1.234)

Because $r_a(\tau)$ is not identically zero, observing (1.227) we find that x is cyclostationary with period $1/f_0$. From (1.232), the average PSD of x is given by

$$\overline{\mathcal{P}}_{x}(f) = \frac{1}{4} [\mathcal{P}_{a}(f - f_{0}) + \mathcal{P}_{a}(f + f_{0})]$$
(1.235)

Therefore, x has a bandwidth equal to $2B_a$ and an average statistical power

$$\overline{\mathbb{M}}_{x} = \frac{1}{2} \mathbb{M}_{a} \tag{1.236}$$

We note that one finds the same result (1.235) assuming that φ_0 is a uniform r.v. in $[0, 2\pi)$; in this case x turns out to be WSS.

Example 1.5.4

Let x be a modulated single sideband (SSB) with an upper sideband, i.e.

$$\begin{aligned} x(t) &= Re\left[\frac{1}{2}\left(a(t) + ja^{(h)}(t)\right)e^{j(2\pi f_0 t + \varphi_0)}\right] \\ &= \frac{1}{2}a(t)\cos(2\pi f_0 t + \varphi_0) - \frac{1}{2}a^{(h)}(t)\sin(2\pi f_0 t + \varphi_0) \end{aligned}$$
(1.237)

¹³ To be precise, x is cyclostationary in mean value with period $T_0 = 1/f_0$, while it is cyclostationary in correlation with period $T_0/2$.



Figure 1.24 Coherent DSB demodulator and baseband-equivalent scheme. (a) Coherent DSB demodulator and (b) baseband-equivalent scheme.

where $a^{(h)}$ is the Hilbert transform of *a*, a real WSS random process with autocorrelation $r_a(\tau)$ and bandwidth B_a .

We note that the modulating signal $(a(t) + ja^{(h)}(t))$ coincides with the analytic signal $a^{(a)}$ and its spectral support contains only positive frequencies.

Being

$$x^{(bb)}(t) = \frac{1}{2}(a(t) + ja^{(h)}(t))e^{j\varphi_0}$$

it results that $x^{(bb)}$ and $x^{(bb)*}$ have non-overlapping passbands and

$$\mathbf{r}_{x^{(bb)}x^{(bb)*}}(\tau) = 0 \tag{1.238}$$

The process (1.237) is then stationary with

$$\mathcal{P}_{x}(f) = \frac{1}{4} \left[\mathcal{P}_{a^{(+)}}(f - f_0) + \mathcal{P}_{a^{(+)}}(-f - f_0) \right]$$
(1.239)

where $a^{(+)}$ is defined in (1.195). In this case, x has bandwidth equal to B_a and statistical power given by

$$\mathbb{M}_x = \frac{1}{4} \mathbb{M}_a \tag{1.240}$$

Example 1.5.5 (DSB and SSB demodulators)

Let the signal r be the sum of a desired part x and additive white noise w with PSD equal to $\mathcal{P}_w(f) = N_0/2$,

$$r(t) = x(t) + w(t)$$
(1.241)

where the signal x is modulated DSB (1.233). To obtain the signal a from r, one can use the coherent demodulation scheme illustrated in Figure 1.24 (see Figure 1.21b), where h is an ideal lowpass filter, having a frequency response

$$\mathcal{H}(f) = \mathcal{H}_0 \operatorname{rect}\left(\frac{f}{2B_a}\right) \tag{1.242}$$

Let r_o be the output signal of the demodulator, given by the sum of the desired part x_o and noise w_o :

$$r_{o}(t) = x_{o}(t) + w_{o}(t) \tag{1.243}$$

We evaluate now the ratio between the powers of the signals in (1.243),

$$\Lambda_o = \frac{\mathsf{M}_{x_o}}{\mathsf{M}_{w_o}} \tag{1.244}$$

in terms of the reference signal-to-noise ratio

$$\Gamma = \frac{M_x}{(N_0/2) \ 2B_a} \tag{1.245}$$

Using the equivalent block scheme of Figure 1.24 and (1.126), we have

$$r^{(bb)}(t) = a(t) e^{j\varphi_0} + w^{(bb)}(t)$$
(1.246)

with $P_{w^{(bb)}}(f) = 2N_0 \ 1(f + f_0)$. Being

$$h * a(t) = \mathcal{H}_0 a(t) \tag{1.247}$$

it results

$$\begin{aligned} x_o(t) &= Re\left[h * \frac{1}{2} e^{-j\varphi_1} a e^{j\varphi_0}\right](t) \\ &= \frac{\mathcal{H}_0}{2} a(t) \cos(\varphi_0 - \varphi_1) \end{aligned} \tag{1.248}$$

Hence, we get

$$M_{x_o} = \frac{H_0^2}{4} M_a \cos^2(\varphi_0 - \varphi_1)$$
(1.249)

In the same baseband equivalent scheme, we consider the noise w_{eq} at the output of filter h; we find

$$\mathcal{P}_{w_{eq}}(f) = \frac{1}{4} |\mathcal{H}(f)|^2 2N_0 1(f+f_0)$$

= $\frac{\mathcal{H}_0^2}{2} N_0 \operatorname{rect}\left(\frac{f}{2B_a}\right)$ (1.250)

Being now w WSS, $w^{(bb)}$ is uncorrelated with $w^{(bb)*}$ and thus w_{eq} with w^*_{eq} . Then, from

$$w_o(t) = w_{eq,I}(t) \tag{1.251}$$

and using (1.209) it follows

$$\mathcal{P}_{w_0}(f) = \frac{\mathcal{H}_0^2}{4} N_0 \operatorname{rect}\left(\frac{f}{2B_a}\right)$$
(1.252)

and

$$\mathbb{M}_{w_0} = \frac{\mathcal{H}_0^2}{4} N_0 \ 2B_a \tag{1.253}$$

In conclusion, using (1.236), we have

$$\Lambda_o = \frac{(\mathcal{H}_0^2/4) \, M_a \cos^2(\varphi_0 - \varphi_1)}{(\mathcal{H}_0^2/4) \, N_0 \, 2B_a} = \Gamma \cos^2(\varphi_0 - \varphi_1) \tag{1.254}$$

For $\varphi_1 = \varphi_0$ (1.254) becomes

$$\Lambda_o = \Gamma \tag{1.255}$$

It is interesting to observe that, at the demodulator input, the ratio between the power of the desired signal and the power of the noise in the *passband of x* is given by

$$\Lambda_i = \frac{M_x}{(N_0/2) \ 4B_a} = \frac{\Gamma}{2}$$
(1.256)

For $\varphi_1 = \varphi_0$ then

$$\Lambda_o = 2\Lambda_i \tag{1.257}$$

We will now analyse the case of a SSB signal x (see (1.237)), coherently demodulated, following the scheme of Figure 1.25, where h_{PB} is a filter used to eliminate the noise that otherwise, after the *mixer*, would have fallen within the passband of the desired signal. The ideal frequency response of h_{PB} is given by

$$\mathcal{H}_{PB}(f) = \operatorname{rect}\left(\frac{f - f_0 - B_a/2}{B_a}\right) + \operatorname{rect}\left(\frac{-f - f_0 - B_a/2}{B_a}\right)$$
(1.258)



Figure 1.25 (a) Coherent SSB demodulator and (b) baseband-equivalent scheme.

Note that in this scheme, we have assumed the phase of the receiver carrier equal to that of the transmitter, to avoid distortion of the desired signal.

Being

$$\mathcal{H}_{PB}^{(bb)}(f) = 2 \operatorname{rect}\left(\frac{f - B_a/2}{B_a}\right)$$
(1.259)

the filter of the baseband-equivalent scheme is given by

$$h_{eq}(t) = \frac{1}{2} h_{PB}^{(bb)} * h(t)$$
(1.260)

with frequency response

$$\mathcal{H}_{eq}(f) = \mathcal{H}_0 \operatorname{rect}\left(\frac{f - B_a/2}{B_a}\right)$$
(1.261)

We now evaluate the desired component x_o . Using the fact $x^{(bb)} * h_{eq}(t) = \mathcal{H}_0 x^{(bb)}(t)$, it results

$$\begin{aligned} x_o(t) &= Re\left[h_{eq} * \frac{1}{2} \ e^{-j\varphi_0} \ \frac{1}{2}(a+j \ a^{(h)}) \ e^{j\varphi_0}\right](t) \\ &= \frac{\mathcal{H}_0}{4} \ Re\left[a(t)+j \ a^{(h)}(t)\right] = \frac{\mathcal{H}_0}{4} \ a(t) \end{aligned}$$
(1.262)

In the baseband-equivalent scheme, the noise w_{eq} at the output of h_{eq} has a PSD given by

$$\mathcal{P}_{w_{eq}}(f) = \frac{1}{4} \left| \mathcal{H}_{eq}(f) \right|^2 2N_0 \ 1(f+f_0) = \frac{N_0}{2} \ \mathcal{H}_0^2 \ \text{rect}\left(\frac{f-B_a/2}{B_a}\right)$$
(1.263)

From the relation $w_o = w_{eq,I}$ and using (1.209), which is valid because $w_{eq} \perp w_{eq}^*$, we have

$$\mathcal{P}_{w_o}(f) = \frac{1}{4} [\mathcal{P}_{w_{eq}}(f) + \mathcal{P}_{w_{eq}}(-f)] = \frac{\mathcal{H}_0^2}{8} N_0 \operatorname{rect}\left(\frac{f}{2B_a}\right)$$
(1.264)

and

$$M_{w_o} = \frac{\mathcal{H}_0^2}{8} N_0 \ 2B_a \tag{1.265}$$

Then we obtain

$$\Lambda_o = \frac{(\mathcal{H}_0^2/16) \,\mathbb{M}_a}{(\mathcal{H}_0^2/8) \,N_0 \,2B_a} \tag{1.266}$$

which using (1.240) and (1.245) can be written as

$$\Lambda_o = \Gamma \tag{1.267}$$

We note that the SSB system yields the same performance (for $\varphi_1 = \varphi_0$) of a DSB system, even though half of the bandwidth is required. Finally, it results

$$\Lambda_i = \frac{M_x}{(N_0/2) \ 2B_a} = \Lambda_o \tag{1.268}$$
Observation 1.4

We note that also for the simple examples considered in this section, the desired signal is analysed via the various transformations, whereas the noise is analysed via the PSD. As a matter of fact, we are typically interested only in the statistical power of the noise at the system output. The demodulated signal x_o , on the other hand, must be expressed as the sum of a desired component proportional to a and an orthogonal component that represents the distortion, which is, typically, small and has the same effects as noise.

In the previous example, the considered systems do not introduce any distortion since x_o is proportional to a.

1.6 The autocorrelation matrix

Definition 1.14

Given the discrete-time wide-sense stationary random process $\{x(k)\}$, we introduce the random vector with *N* components

$$\mathbf{x}^{T}(k) = [x(k), x(k-1), \dots, x(k-N+1)]$$
(1.269)

The $N \times N$ autocorrelation matrix of $x^*(k)$ is given by

$$\boldsymbol{R} = E[\boldsymbol{x}^{*}(k)\boldsymbol{x}^{T}(k)] = \begin{bmatrix} \mathbf{r}_{x}(0) & \mathbf{r}_{x}(-1) & \dots & \mathbf{r}_{x}(-N+1) \\ \mathbf{r}_{x}(1) & \mathbf{r}_{x}(0) & \dots & \mathbf{r}_{x}(-N+2) \\ \vdots & \vdots & \ddots & \dots \\ \mathbf{r}_{x}(N-1) & \mathbf{r}_{x}(N-2) & \dots & \mathbf{r}_{x}(0) \end{bmatrix}.$$
(1.270)

Properties

- 1. **R** is Hermitian: $\mathbf{R}^{H} = \mathbf{R}$. For real random processes **R** is symmetric: $\mathbf{R}^{T} = \mathbf{R}$.
- 2. **R** is a *Toeplitz matrix*, i.e. all elements along any diagonal are equal.
- 3. **R** is *positive semi-definite* and almost always *positive definite*. Indeed, taking an arbitrary vector $\boldsymbol{v}^T = [v_0, \dots, v_{N-1}]$, and letting $y_k = \boldsymbol{x}^T(k)\boldsymbol{v}$, we have

$$E[|y_k|^2] = E[\boldsymbol{v}^H \boldsymbol{x}^*(k) \boldsymbol{x}^T(k) \boldsymbol{v}] = \boldsymbol{v}^H \boldsymbol{R} \boldsymbol{v} = \sum_{i=0}^{N-1} \sum_{j=0}^{N-1} v_i^* \boldsymbol{r}_x(i-j) v_j \ge 0$$
(1.271)

If $v^H R v > 0$, $\forall v \neq 0$, then *R* is said to be *positive definite* and all its *principal minor* determinants are positive; in particular *R* is *non-singular*.

Eigenvalues

We indicate with det **R** the determinant of a matrix **R**. The eigenvalues of **R** are the solutions λ_i , i = 1, ..., N, of the *characteristic equation* of order N

$$\det[\boldsymbol{R} - \lambda \boldsymbol{I}] = 0 \tag{1.272}$$

and the corresponding column eigenvectors u_i , i = 1, ..., N, satisfy the equation

$$\boldsymbol{R}\boldsymbol{u}_i = \lambda_i \boldsymbol{u}_i \tag{1.273}$$

Table 1.7: Correspondence between eigenvalues and eigenvectors of four matrices.

	R	\mathbf{R}^{m}	\mathbf{R}^{-1}	$I - \mu R$
Eigenvalue	λ_i	λ_i^m	λ_i^{-1}	$(1 - \mu \lambda_i)$
Eigenvector	u _i	u _i	\boldsymbol{u}_i	u _i

Example 1.6.1 Let $\{w(k)\}$ be a white noise process. Its autocorrelation matrix **R** assumes the form

$$\boldsymbol{R} = \begin{bmatrix} \sigma_{w}^{2} & 0 & \dots & 0 \\ 0 & \sigma_{w}^{2} & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & \sigma_{w}^{2} \end{bmatrix}$$
(1.274)

from which it follows that

$$\lambda_1 = \lambda_2 = \dots = \lambda_N = \sigma_w^2 \tag{1.275}$$

and

$$u_i$$
 can be any arbitrary vector $1 \le i \le N$ (1.276)

Example 1.6.2

We define a complex-valued sinusoid as

$$x(k) = e^{j(\omega k + \varphi)}, \qquad \omega = 2\pi f T_c \tag{1.277}$$

with φ a uniform r.v. in $[0, 2\pi)$. The autocorrelation matrix **R** is given by

$$\boldsymbol{R} = \begin{bmatrix} 1 & e^{-j\omega} & \dots & e^{-j(N-1)\omega} \\ e^{j\omega} & 1 & \dots & e^{-j(N-2)\omega} \\ \vdots & \vdots & \ddots & \vdots \\ e^{j(N-1)\omega} & e^{j(N-2)\omega} & \dots & 1 \end{bmatrix}$$
(1.278)

One can see that the rank of R is 1 and it will therefore have only one eigenvalue. The solution is given by

$$\lambda_1 = N \tag{1.279}$$

and the relative eigenvector is

$$\boldsymbol{u}_{1}^{T} = [1, e^{j\omega}, \dots, e^{j(N-1)\omega}]$$
(1.280)

Other properties

- 1. From $\mathbf{R}^m \mathbf{u} = \lambda^m \mathbf{u}$, we obtain the relations of Table 1.7.
- 2. If the eigenvalues are distinct, then the eigenvectors are linearly independent:

$$\sum_{i=1}^{N} c_i \boldsymbol{u}_i \neq \boldsymbol{0} \tag{1.281}$$

for all combinations of $\{c_i\}$, i = 1, 2, ..., N, not all equal to zero. Therefore, in this case, the eigenvectors form a basis in \mathbb{R}^N .

3. The *trace* of a matrix **R** is defined as the sum of the elements of the main diagonal, and we indicate it with tr **R**. It holds

$$\operatorname{tr} \boldsymbol{R} = \sum_{i=1}^{N} \lambda_i \tag{1.282}$$

Eigenvalue analysis for Hermitian matrices

As previously seen, the autocorrelation matrix \boldsymbol{R} is Hermitian, thus enjoys the following properties:

1. The eigenvalues of a Hermitian matrix are real. By left multiplying both sides of (1.273) by u_i^H , it follows

$$\boldsymbol{u}_i^H \boldsymbol{R} \boldsymbol{u}_i = \lambda_i \boldsymbol{u}_i^H \boldsymbol{u}_i \tag{1.283}$$

from which, by the definition of norm, we obtain

$$\lambda_i = \frac{\boldsymbol{u}_i^H \boldsymbol{R} \boldsymbol{u}_i}{\boldsymbol{u}_i^H \boldsymbol{u}_i} = \frac{\boldsymbol{u}_i^H \boldsymbol{R} \boldsymbol{u}_i}{\|\boldsymbol{u}_i\|^2}$$
(1.284)

The ratio (1.284) is defined as *Rayleigh quotient*. As **R** is positive semi-definite, $\boldsymbol{u}_i^H \boldsymbol{R} \boldsymbol{u}_i \ge 0$, from which $\lambda_i \ge 0$.

2. If the eigenvalues of R are distinct, then the eigenvectors are orthogonal. In fact, from (1.273), we obtain:

$$\boldsymbol{u}_i^H \boldsymbol{R} \boldsymbol{u}_j = \lambda_j \boldsymbol{u}_i^H \boldsymbol{u}_j \tag{1.285}$$

$$\boldsymbol{u}_i^H \boldsymbol{R} \boldsymbol{u}_j = \lambda_i \boldsymbol{u}_i^H \boldsymbol{u}_j \tag{1.286}$$

Subtracting the second equation from the first:

$$0 = (\lambda_j - \lambda_i) \boldsymbol{u}_i^H \boldsymbol{u}_j \tag{1.287}$$

and since $\lambda_i - \lambda_i \neq 0$ by hypothesis, it follows $\boldsymbol{u}_i^H \boldsymbol{u}_i = 0$.

3. If the eigenvalues of R are distinct and their corresponding eigenvectors are normalized, i.e.

$$\|\boldsymbol{u}_{i}\|^{2} = \boldsymbol{u}_{i}^{H}\boldsymbol{u}_{i} = \begin{cases} 1 & i = j \\ 0 & i \neq j \end{cases}$$
(1.288)

then the matrix $U = [u_1, u_2, ..., u_N]$, whose columns are the eigenvectors of R, is a unitary matrix, that is

$$\boldsymbol{U}^{-1} = \boldsymbol{U}^H \tag{1.289}$$

This property is an immediate consequence of the orthogonality of the eigenvectors $\{u_i\}$. Moreover, if we define the matrix

$$\mathbf{\Lambda} = \begin{bmatrix} \lambda_1 & 0 & \dots & 0 \\ 0 & \lambda_2 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & \lambda_N \end{bmatrix}$$
(1.290)

we get

$$\boldsymbol{U}^{H}\boldsymbol{R}\boldsymbol{U} = \boldsymbol{\Lambda} \tag{1.291}$$

From (1.291), we obtain the following important relations:

$$\boldsymbol{R} = \boldsymbol{U}\boldsymbol{\Lambda}\boldsymbol{U}^{H} = \sum_{i=1}^{N} \lambda_{i}\boldsymbol{u}_{i}\boldsymbol{u}_{i}^{H}$$
(1.292)

and

$$\boldsymbol{I} - \boldsymbol{\mu}\boldsymbol{R} = \boldsymbol{U}(\boldsymbol{I} - \boldsymbol{\mu}\boldsymbol{\Lambda})\boldsymbol{U}^{H} = \sum_{i=1}^{N} (1 - \boldsymbol{\mu}\lambda_{i})\boldsymbol{u}_{i}\boldsymbol{u}_{i}^{H}$$
(1.293)

4. The eigenvalues of a positive semi-definite autocorrelation matrix \mathbf{R} and the PSD of x are related by the inequalities,

$$\min_{f} \{\mathcal{P}_{x}(f)\} \le \lambda_{i} \le \max_{f} \{\mathcal{P}_{x}(f)\}, \quad i = 1, \dots, N$$
(1.294)

In fact, let $U_i(f)$ be the Fourier transform of the sequence represented by the elements of u_i , i.e.

$$U_i(f) = \sum_{n=1}^{N} u_{i,n} e^{-j2\pi f n T_c}$$
(1.295)

where $u_{i,n}$ is the *n*-th element of the eigenvector u_i . Observing that

$$\boldsymbol{u}_{i}^{H}\boldsymbol{R}\boldsymbol{u}_{i} = \sum_{n=1}^{N} \sum_{m=1}^{N} u_{i,n}^{*} \boldsymbol{r}_{x}(n-m) u_{i,m}$$
(1.296)

and using (1.172) and (1.295), we have

$$\boldsymbol{u}_{i}^{H}\boldsymbol{R}\boldsymbol{u}_{i} = \int_{-\frac{1}{2T_{c}}}^{\frac{1}{2T_{c}}} \mathcal{P}_{x}(f) \sum_{n=1}^{N} u_{i,n}^{*} e^{j2\pi f nT_{c}} \sum_{m=1}^{N} u_{i,m} e^{-j2\pi f mT_{c}} df$$
$$= \int_{-\frac{1}{2T_{c}}}^{\frac{1}{2T_{c}}} \mathcal{P}_{x}(f) |U_{i}(f)|^{2} df \qquad (1.297)$$

Substituting the latter result in (1.284) one finds

$$\lambda_{i} = \frac{\int_{-\frac{1}{2T_{c}}}^{\frac{1}{2T_{c}}} \mathcal{P}_{x}(f) |U_{i}(f)|^{2} df}{\int_{-\frac{1}{2T_{c}}}^{\frac{1}{2T_{c}}} |U_{i}(f)|^{2} df}$$
(1.298)

from which (1.294) follows.

If we indicate with λ_{min} and λ_{max} , respectively, the minimum and maximum eigenvalue of **R**, in view of the latter point, we can define the *eigenvalue spread* as:

$$\chi(\mathbf{R}) = \frac{\lambda_{max}}{\lambda_{min}} \le \frac{\max_{f} \{\mathcal{P}_{x}(f)\}}{\min_{f} \{\mathcal{P}_{x}(f)\}}$$
(1.299)

From (1.299), we observe that $\chi(\mathbf{R})$ may assume large values in the case $\mathcal{P}_{\chi}(f)$ exhibits large variations. Moreover, $\chi(\mathbf{R})$ assumes the minimum value of 1 for a white process.

1.7 Examples of random processes

Before reviewing some important random processes, we recall the definition of Gaussian complex-valued random vector.

Example 1.7.1

A complex r.v. with a Gaussian distribution can be generated from two r.v.s. with uniform distribution (see Appendix 1.B for an illustration of the method).

Example 1.7.2

Let $\mathbf{x}^T = [x_1, \dots, x_N]$ be a real Gaussian random vector, each component has mean \mathbf{m}_{x_i} and variance $\sigma_{x_i}^2$, denoted as $x_i \sim \mathcal{N}(\mathbf{m}_{x_i}, \sigma_{x_i}^2)$. The joint probability density function (pdf) is

$$p_{\mathbf{x}}(\boldsymbol{\xi}) = [(2\pi)^{N} \det \boldsymbol{C}_{N}]^{-\frac{1}{2}} e^{-\frac{1}{2}(\boldsymbol{\xi} - \boldsymbol{m}_{\mathbf{x}})^{T} \boldsymbol{C}_{N}^{-1}(\boldsymbol{\xi} - \boldsymbol{m}_{\mathbf{x}})}$$
(1.300)

where $\xi^T = [\xi_1, \dots, \xi_N]$, $m_x = E[x]$ is the vector of its components' mean values and $C_N = E[(x - m_x)(x - m_x)^T]$ is its covariance matrix.

Example 1.7.3

Let $\mathbf{x}^T = [x_{1,I} + jx_{1,Q}, \dots, x_{N,I} + jx_{N,Q}]$ be a complex-valued Gaussian random vector. If the in-phase component $x_{i,I}$ and the quadrature component $x_{i,Q}$ are uncorrelated,

$$E[(x_{i,I} - \mathbf{m}_{x_{i,I}})(x_{i,Q} - \mathbf{m}_{x_{i,Q}})] = 0, \quad i = 1, 2, \dots, N$$
(1.301)

Moreover, we have

$$\sigma_{x_{i,l}}^2 = \sigma_{x_{i,Q}}^2 = \frac{1}{2}\sigma_{x_i}^2 \tag{1.302}$$

then the joint pdf is

$$p_{\mathbf{x}}(\boldsymbol{\xi}) = [\pi^{N} \det \boldsymbol{C}_{N}]^{-1} e^{-(\boldsymbol{\xi} - \boldsymbol{m}_{x})^{H} \boldsymbol{C}_{N}^{-1}(\boldsymbol{\xi} - \boldsymbol{m}_{x})}$$
(1.303)

with the vector of mean values and the covariance matrix given by

$$\boldsymbol{m}_{\mathbf{x}} = \boldsymbol{E}[\boldsymbol{x}] = \boldsymbol{E}[\boldsymbol{x}_{I}] + j\boldsymbol{E}[\boldsymbol{x}_{O}] \tag{1.304}$$

$$C_N = E[(x - m_x)(x - m_x)^H]$$
(1.305)

Vector \mathbf{x} is called *circularly symmetric Gaussian* random vector. For the generic component, we write $x_i \sim C \mathcal{N}(\mathbb{m}_{x_i}, \sigma_{x_i}^2)$ and

$$p_{x_i}(\xi_i) = \frac{1}{\sqrt{2\pi\sigma_{x_{i,l}}^2}} e^{-\frac{|\xi_{i,l} - \mathbb{E}_{x_{i,l}}|^2}{2\sigma_{x_{i,l}}^2}} \frac{1}{\sqrt{2\pi\sigma_{x_{i,Q}}^2}} e^{-\frac{|\xi_{i,Q} - \mathbb{E}_{x_{i,Q}}|^2}{2\sigma_{x_{i,Q}}^2}}$$
(1.306)

$$=\frac{1}{\pi\sigma_{x_i}^2}e^{-\frac{|\xi_i-\mathbf{n}_{x_i}|^2}{\sigma_{x_i}^2}}$$
(1.307)

with $\xi_i = \xi_{i,I} + j\xi_{i,Q}$ complex valued.

Example 1.7.4

Let $\mathbf{x}^T = [x_1, \dots, x_N] = [x_1(t_1), \dots, x_N(t_N)]$ be a complex-valued Gaussian (vector) process, with each element $x_i(t_i)$ having real and imaginary components that are uncorrelated Gaussian r.v.s. whose pdf is with zero mean and equal variance for all values of t_i . The vector \mathbf{x} is called *circularly symmetric Gaussian random process*. The joint pdf in this case results

$$p_{\mathbf{x}}(\boldsymbol{\xi}) = [\pi^{N} \det \boldsymbol{C}]^{-1} e^{-\boldsymbol{\xi}^{H} \boldsymbol{C}^{-1} \boldsymbol{\xi}}$$
(1.308)

where **C** is the covariance matrix of $[x_1(t_1), x_2(t_2), \dots, x_N(t_N)]$.

Example 1.7.5

Let $x(t) = A \sin(2\pi ft + \varphi)$ be a real-valued sinusoidal signal with φ r.v. uniform in $[0, 2\pi)$, for which we will use the notation $\varphi \sim \mathcal{U}[0, 2\pi)$. The mean of x is

$$m_{x}(t) = E[x(t)]$$

$$= \int_{0}^{2\pi} \frac{1}{2\pi} A \sin(2\pi f t + a) da \qquad (1.309)$$

$$= 0$$

and the autocorrelation function is given by

$$r_{x}(\tau) = \int_{0}^{2\pi} \frac{1}{2\pi} A \sin(2\pi f t + a) A \sin[2\pi f (t - \tau) + a] da$$
$$= \frac{A^{2}}{2} \cos(2\pi f \tau)$$
(1.310)

Example 1.7.6 Consider the sum of *N* real-valued sinusoidal signals, i.e.

$$x(t) = \sum_{i=1}^{N} A_i \sin(2\pi f_i t + \varphi_i)$$
(1.311)

with $\varphi_i \sim \mathcal{U}[0, 2\pi)$ statistically independent, from Example 1.7.5 it is immediate to obtain the mean

$$\mathbf{m}_{x}(t) = \sum_{i=1}^{N} \mathbf{m}_{x_{i}}(t) = 0$$
(1.312)

and the autocorrelation function

$$\mathbf{r}_{x}(\tau) = \sum_{i=1}^{N} \frac{A_{i}^{2}}{2} \cos(2\pi f_{i}\tau)$$
(1.313)

We note that, according to the Definition 1.9, page 31, the process (1.311) is not asymptotically uncorrelated.

Example 1.7.7

Consider the sum of N complex-valued sinusoidal signals, i.e.

$$x(t) = \sum_{i=1}^{N} A_i \ e^{i(2\pi f_i t + \varphi_i)}$$
(1.314)

with $\varphi_i \sim \mathcal{U}[0, 2\pi)$ statistically independent. Following a similar procedure to that used in Examples 1.7.5 and 1.7.6, we find

$$\mathbf{r}_{x}(\tau) = \sum_{i=1}^{N} |A_{i}|^{2} e^{j2\pi f_{i}\tau}$$
(1.315)

We note that the process (1.315) is not asymptotically uncorrelated.

$$x(k)$$
 h_{Tx} $y(t)$

Figure 1.26 Modulator of a PAM system as interpolator filter.

Example 1.7.8

Let the discrete-time random process y(k) = x(k) + w(k) be given by the sum of the random process x of Example 1.7.7 and white noise w with variance σ_w^2 . Moreover, we assume x and w uncorrelated. In this case,

$$\mathbf{r}_{y}(n) = \sum_{i=1}^{N} |A_{i}|^{2} e^{j2\pi f_{i}nT_{c}} + \sigma_{w}^{2}\delta_{n}$$
(1.316)

Example 1.7.9

We consider a signal obtained by pulse-amplitude modulation (PAM), expressed as

$$y(t) = \sum_{k=-\infty}^{+\infty} x(k) h_{Tx}(t - kT)$$
(1.317)

The signal y is the output of the system shown in Figure 1.26, where h_{Tx} is a finite-energy pulse and $\{x(k)\}$ is a discrete-time (with *T*-spaced samples) WSS sequence, having PSD $\mathcal{P}_x(f)$. We note that $\mathcal{P}_x(f)$ is a periodic function of period 1/T.

Let the deterministic autocorrelation of the signal h_{Tx} be

$$\mathbf{r}_{h_{T_x}}(\tau) = \int_{-\infty}^{+\infty} h_{T_x}(t) h_{T_x}^*(t-\tau) dt = [h_{T_x}(t) * h_{T_x}^*(-t)](\tau)$$
(1.318)

with Fourier transform $|\mathcal{H}_{T_x}(f)|^2$. In general, y is a cyclostationary process of period T. In fact, we have

1. Mean

$$m_{y}(t) = m_{x} \sum_{k=-\infty}^{+\infty} h_{Tx}(t - kT)$$
(1.319)

2. Correlation

$$\mathbf{r}_{y}(t,t-\tau) = \sum_{i=-\infty}^{+\infty} \mathbf{r}_{x}(i) \sum_{m=-\infty}^{+\infty} h_{Tx}(t-(i+m)\ T)h_{Tx}^{*}(t-\tau-mT)$$
(1.320)

If we introduce the average spectral analysis

$$\bar{\mathbf{m}}_{y} = \frac{1}{T} \int_{0}^{T} \mathbf{m}_{y}(t) \, dt = \mathbf{m}_{x} \mathcal{H}_{Tx}(0) \tag{1.321}$$

$$\overline{\mathbf{r}}_{y}(\tau) = \frac{1}{T} \int_{0}^{T} \mathbf{r}_{y}(t, t - \tau) dt = \frac{1}{T} \sum_{i=-\infty}^{+\infty} \mathbf{r}_{x}(i) \mathbf{r}_{h_{Tx}}(\tau - iT)$$
(1.322)

and

$$\overline{\mathcal{P}}_{y}(f) = \mathcal{F}[\overline{\mathbf{r}}_{y}(\tau)] = \left|\frac{1}{T}\mathcal{H}_{Tx}(f)\right|^{2}\mathcal{P}_{x}(f)$$
(1.323)

we observe that the modulator of a PAM system may be regarded as an interpolator filter with frequency response \mathcal{H}_{Tx}/T .

3. Average power for a white noise input For a white noise input with power M_x , from (1.322), the average statistical power of the output signal is given by

$$\overline{\mathbb{M}}_{y} = \mathbb{M}_{x} \, \frac{E_{h}}{T} \tag{1.324}$$

where $E_h = \int_{-\infty}^{+\infty} |h_{Tx}(t)|^2 dt$ is the energy of h_{Tx} . 4. Moments of y for a circularly symmetric i.i.d. input

Let $\{x(k)\}$ be a complex-valued random circularly symmetric sequence with zero mean (see (1.301) and (1.302)), i.e. letting

$$x_I(k) = Re[x(k)], \quad x_O(k) = Im[x(k)]$$
 (1.325)

we have

$$E[x_I^2(k)] = E[x_Q^2(k)] = \frac{E[|x(k)|^2]}{2}$$
(1.326)

and

$$E[x_l(k) \ x_O(k)] = 0 \tag{1.327}$$

These two relations can be merged into the single expression

$$E[x^{2}(k)] = E[x_{I}^{2}(k)] - E[x_{Q}^{2}(k)] + 2j \ E[x_{I}(k) \ x_{Q}(k)] = 0$$
(1.328)

Filtering the i.i.d. input signal $\{x(k)\}$ by using the system depicted in Figure 1.26, and from the relation

$$\mathbf{r}_{yy^*}(t,t-\tau) = \sum_{i=-\infty}^{+\infty} \mathbf{r}_{xx^*}(i) \sum_{m=-\infty}^{+\infty} h_{Tx}(t-(i+m)T)h_{Tx}(t-\tau-mT)$$
(1.329)

we have

$$\mathbf{r}_{xx^*}(i) = E[x^2(k)]\delta(i) = 0 \tag{1.330}$$

and

$$\mathbf{r}_{vv^*}(t, t - \tau) = 0 \tag{1.331}$$

that is $y \perp y^*$. In particular, we have that y is circularly symmetric, i.e.

$$E[y^2(t)] = 0 \tag{1.332}$$

We note that the condition (1.331) can be obtained assuming the less stringent condition that $x \perp x^*$; on the other hand, this requires that the following two conditions are verified

$$\mathbf{r}_{x_i}(i) = \mathbf{r}_{x_0}(i) \tag{1.333}$$

and

$$\mathbf{r}_{x_{l}x_{o}}(i) = -\mathbf{r}_{x_{l}x_{o}}(-i) \tag{1.334}$$

Observation 1.5

It can be shown that if the filter h_{Tx} has a bandwidth smaller than 1/(2T) and $\{x(k)\}$ is a WSS sequence, then $\{y(k)\}$ is WSS with PSD (1.323).

Example 1.7.10

Let us consider a PAM signal sampled with period $T_Q = T/Q_0$, where Q_0 is a positive integer number. Let

$$y_q = y(q T_Q), \quad h_p = h_{Tx}(p T_Q)$$
 (1.335)

from (1.317) it follows

$$y_q = \sum_{k=-\infty}^{+\infty} x(k) \ h_{q-kQ_0}$$
(1.336)

If $Q_0 \neq 1$, (1.336) describes the input–output relation of an interpolator filter (see (1.536)). We recall the statistical analysis given in Table 1.6, page 34. We denote with $\mathcal{H}(f)$ the Fourier transform (see (1.17)) and with $r_h(n)$ the deterministic autocorrelation (see (1.184)) of the sequence $\{h_p\}$. We also assume that $\{x(k)\}$ is a WSS random sequence with mean m_x and autocorrelation $r_x(n)$. In general, $\{y_q\}$ is a cyclostationary random sequence of period Q_0 with

1. Mean

$$m_{y}(q) = m_{x} \sum_{k=-\infty}^{+\infty} h_{q-kQ_{0}}$$
(1.337)

2. Correlation

$$r_{y}(q,q-n) = \sum_{i=-\infty}^{+\infty} r_{x}(i) \sum_{m=-\infty}^{+\infty} h_{q-(i+m)Q_{0}} h_{q-n-m Q_{0}}^{*}$$
(1.338)

By the average spectral analysis, we obtain

$$\overline{m}_{y} = \frac{1}{Q_{0}} \sum_{q=0}^{Q_{0}-1} m_{y}(q) = m_{x} \frac{\mathcal{H}(0)}{Q_{0}}$$
(1.339)

where

$$\mathcal{H}(0) = \sum_{p=-\infty}^{+\infty} h_p \tag{1.340}$$

and

$$\overline{\mathbf{r}}_{y}(n) = \frac{1}{Q_{0}} \sum_{q=0}^{Q_{0}-1} \mathbf{r}_{y}(q, q-n) = \frac{1}{Q_{0}} \sum_{i=-\infty}^{+\infty} \mathbf{r}_{x}(i) \mathbf{r}_{h}(n-iQ_{0})$$
(1.341)

Consequently, the average PSD is given by

$$\overline{\mathcal{P}}_{y}(f) = T_{Q} \mathcal{F}[\overline{r}_{y}(n)] = \left| \frac{1}{Q_{0}} \mathcal{H}(f) \right|^{2} \mathcal{P}_{x}(f)$$
(1.342)

If $\{x(k)\}$ is white noise with power M_x , from (1.341) it results

$$\overline{\mathbf{r}}_{y}(n) = \mathbb{M}_{x} \frac{\mathbf{r}_{h}(n)}{Q_{0}}$$
(1.343)

In particular, the average power of the filter output signal is given by

$$\overline{\mathbb{M}}_{y} = \mathbb{M}_{x} \frac{E_{h}}{Q_{0}}$$
(1.344)

where $E_h = \sum_{p=-\infty}^{+\infty} |h_p|^2$ is the energy of $\{h_p\}$. We point out that the condition $\overline{\mathbb{M}}_y = \mathbb{M}_x$ is satisfied if the energy of the filter impulse response is equal to the interpolation factor Q_0 .

$$\begin{aligned} x(t) &= g(t) + w(t) \\ & \searrow \\ g_M \end{aligned} \qquad y(t) \qquad \swarrow \\ y(t_0) &= g_u(t_0) + w_u(t_0) \\ & \searrow \\ \mathcal{G}_M(f) &= K \frac{\mathcal{G}^*(f)}{\mathcal{P}_w(f)} e^{-j2\pi f t_0} \end{aligned}$$

Figure 1.27 Reference scheme for the matched filter.

1.8 Matched filter

Referring to Figure 1.27, we consider a finite-energy signal pulse g in the presence of additive noise w having zero mean and PSD \mathcal{P}_w . The signal

$$x(t) = g(t) + w(t)$$
(1.345)

is filtered with a filter having impulse response g_M . We indicate with g_u and w_u , respectively, the desired signal and the noise component at the output:

$$g_u(t) = g_M * g(t)$$
 (1.346)

$$w_u(t) = g_M * w(t)$$
 (1.347)

The output is

$$y(t) = g_u(t) + w_u(t)$$
(1.348)

We now suppose that y is observed at a given instant t_0 . The problem is to determine g_M so that the ratio between the square amplitude of $g_u(t_0)$ and the power of the noise component $w_u(t_0)$ is maximum, i.e.

$$g_M: \max_{g_M} \frac{|g_u(t_0)|^2}{E[|w_u(t_0)|^2]}$$
(1.349)

The optimum filter has frequency response

$$\mathcal{G}_{M}(f) = K \frac{\mathcal{G}^{*}(f)}{\mathcal{P}_{w}(f)} e^{-j2\pi f t_{0}}$$
(1.350)

where K is a constant. In other words, the best filter selects the frequency components of the desired input signal and weights them with weights that are inversely proportional to the noise level.

Proof. $g_u(t_0)$ coincides with the inverse Fourier transform of $\mathcal{G}_M(f)\mathcal{G}(f)$ evaluated in $t = t_0$, while the power of $w_u(t_0)$ is equal to

$$\mathbf{r}_{w_{u}}(0) = \int_{-\infty}^{+\infty} \mathcal{P}_{w}(f) |\mathcal{G}_{M}(f)|^{2} df$$
(1.351)

Then we have

$$\frac{|g_{u}(t_{0})|^{2}}{\mathbf{r}_{w_{u}}(0)} = \frac{\left|\int_{-\infty}^{+\infty} \mathcal{G}_{M}(f)\mathcal{G}(f)e^{j2\pi ft_{0}}df\right|^{2}}{\int_{-\infty}^{+\infty} \mathcal{P}_{w}(f)|\mathcal{G}_{M}(f)|^{2}df}$$
$$= \frac{\left|\int_{-\infty}^{+\infty} \mathcal{G}_{M}(f)\sqrt{\mathcal{P}_{w}(f)}\frac{\mathcal{G}(f)}{\sqrt{\mathcal{P}_{w}(f)}}e^{j2\pi ft_{0}}df\right|^{2}}{\int_{-\infty}^{+\infty} \mathcal{P}_{w}(f)|\mathcal{G}_{M}(f)|^{2}df}$$
(1.352)

where the integrand at the numerator was divided and multiplied by $\sqrt{\mathcal{P}_w(f)}$. Implicitly, it is assumed that $\mathcal{P}_w(f) \neq 0$. Applying the Schwarz inequality¹⁴ to the functions

$$\mathcal{G}_M(f)\sqrt{\mathcal{P}_w(f)} \tag{1.355}$$

and

$$\frac{\mathcal{G}^*(f)}{\sqrt{\mathcal{P}_w(f)}} \ e^{-j2\pi f t_0}$$
(1.356)

it turns out

$$\frac{|g_{u}(t_{0})|^{2}}{r_{w_{u}}(0)} \leq \int_{-\infty}^{+\infty} \left| \frac{\mathcal{G}(f)}{\sqrt{\mathcal{P}_{w}(f)}} e^{j2\pi f t_{0}} \right|^{2} df = \int_{-\infty}^{+\infty} \left| \frac{\mathcal{G}(f)}{\sqrt{\mathcal{P}_{w}(f)}} \right|^{2} df$$
(1.357)

Therefore, the maximum value is equal to the right-hand side of (1.357) and is achieved for

$$\mathcal{G}_{M}(f)\sqrt{\mathcal{P}_{w}(f)} = K \frac{\mathcal{G}^{*}(f)}{\sqrt{\mathcal{P}_{w}(f)}} e^{-j2\pi f t_{0}}$$
(1.358)

where K is a constant. From (1.358), the solution (1.350) follows immediately.

White noise case

If w is white, then $\mathcal{P}_w(f) = \mathcal{P}_w$ is a constant and the optimum solution (1.350) becomes

$$\mathcal{G}_{M}(f) = K\mathcal{G}^{*}(f)e^{-j2\pi ft_{0}}$$
(1.359)

Correspondingly, the filter has impulse response

$$g_M(t) = Kg^*(t_0 - t) \tag{1.360}$$

from which the name of *matched filter* (MF), i.e. matched to the input signal pulse. The desired signal pulse at the filter output has the frequency response

$$\mathcal{G}_{\mu}(f) = K |\mathcal{G}(f)|^2 e^{-j2\pi f t_0}$$
(1.361)

From the definition of the autocorrelation function of pulse *g*,

$$r_{g}(\tau) = \int_{-\infty}^{+\infty} g(a)g^{*}(a-\tau)da$$
 (1.362)

then, as depicted in Figure 1.28,

$$g_u(t) = Kr_g(t - t_0)$$
(1.363)

i.e. the pulse at the filter output coincides with the autocorrelation function of the pulse g. If E_g is the energy of g, using the relation $E_g = r_g(0)$ the maximum of the functional (1.349) becomes

$$\frac{|g_u(t_0)|^2}{\mathsf{r}_{w_u}(0)} = \frac{|K|^2 \mathsf{r}_g^2(0)}{\mathcal{P}_w[K]^2 \mathsf{r}_g(0)} = \frac{E_g}{\mathcal{P}_w}$$
(1.364)

¹⁴ Given two signals x and y it holds

$$\left| \int_{-\infty}^{\infty} x(t) y^{*}(t) dt \right|^{2} \leq \int_{-\infty}^{\infty} |x(t)|^{2} dt \int_{-\infty}^{\infty} |y(t)|^{2} dt$$
(1.353)

where equality holds if and only if

$$y(t) = Kx(t) \tag{1.354}$$

with K a complex constant.

$$x(t) = g(t) + w(t) \qquad y(t) = Kr_g(t - t_0) + w_u(t) \qquad \downarrow^{t_0} \qquad y(t_0)$$

$$g_M(t) = Kg^*(t_0 - t)$$

Figure 1.28 Matched filter for an input pulse in the presence of white noise.



Figure 1.29 Various pulse shapes related to a matched filter.

In Figure 1.29, the different pulse shapes are illustrated for a signal pulse g with limited duration t_g . Note that in this case, the matched filter has also limited duration, and it is causal if $t_0 \ge t_g$.

Example 1.8.1 (MF for a rectangular pulse) Let

$$g(t) = \mathbf{w}_T(t) = \operatorname{rect}\left(\frac{t - T/2}{T}\right)$$
(1.365)

with

$$\mathbf{r}_{g}(\tau) = T\left(1 - \frac{|\tau|}{T}\right) \operatorname{rect}\left(\frac{\tau}{2T}\right)$$
(1.366)

For $t_0 = T$, the matched filter is proportional to g

$$g_M(t) = K \mathbf{w}_T(t) \tag{1.367}$$

and the output pulse in the absence of noise is equal to

$$g_u(t) = KT\left(1 - \left|\frac{t-T}{T}\right|\right) \operatorname{rect}\left(\frac{t-T}{2T}\right)$$
(1.368)

1.9 Ergodic random processes

The functions that have been introduced in the previous sections for the analysis of random processes give a valid statistical description of an ensemble of realizations of a random process. We investigate now the possibility of moving from ensemble averaging to time averaging, that is we consider the problem of estimating a statistical descriptor of a random process from the observation of a single realization. Let *x* be a discrete-time WSS random process having mean m_x . If in the limit it holds¹⁵

$$\lim_{K \to \infty} \frac{1}{K} \sum_{k=0}^{K-1} x(k) = E[x(k)] = m_x$$
(1.369)

then x is said to be *ergodic in the mean*. In other words, for when the above limit holds, the time-average of samples tends to the statistical mean as the number of samples increases. We note that the existence of the limit (1.369) implies the condition

$$\lim_{K \to \infty} E\left[\left| \frac{1}{K} \sum_{k=0}^{K-1} x(k) - \mathbf{m}_x \right|^2 \right] = 0$$
(1.370)

or equivalently

$$\lim_{K \to \infty} \frac{1}{K} \sum_{n=-(K-1)}^{K-1} \left[1 - \frac{|n|}{K} \right] c_x(n) = 0$$
(1.371)

From (1.371), we see that for a random process to be ergodic in the mean, some conditions on the second-order statistics must be verified. Analogously to definition (1.369), we say that *x* is *ergodic in correlation* if in the limits it holds:

$$\lim_{K \to \infty} \frac{1}{K} \sum_{k=0}^{K-1} x(k) x^*(k-n) = E[x(k)x^*(k-n)] = r_x(n)$$
(1.372)

Also for processes that are ergodic in correlation, one could get a condition of ergodicity similar to that expressed by the limit (1.371). Let $y(k) = x(k)x^*(k - n)$. Observing (1.372) and (1.369), we find that the ergodicity in correlation of the process *x* is equivalent to the ergodicity in the mean of the process *y*. Therefore, it is easy to deduce that the condition (1.371) for *y* translates into a condition on the statistical moments of the fourth order for *x*.

In practice, we will assume all stationary processes to be ergodic; ergodicity is however difficult to prove for non-Gaussian random processes. We will not consider particular processes that are not ergodic such as x(k) = A, where A is a random variable, or x(k) equal to the sum of sinusoidal signals (see (1.311)).

¹⁵ The *limit* is meant in the mean square sense, that is the variance of the r.v. $\left(\frac{1}{K}\sum_{k=0}^{K-1} x(k) - \mathbb{m}_x\right)$ vanishes for $K \to \infty$.



Figure 1.30 Relation between ergodic processes and their statistical description.

The property of ergodicity assumes a fundamental importance if we observe that from a single Realization, it is possible to obtain an estimate of the autocorrelation function, and from this, the PSD. Alternatively, one could prove that under the hypothesis¹⁶

n

$$\sum_{x=-\infty}^{+\infty} |n| \mathbf{r}_x(n) < \infty \tag{1.373}$$

the following limit holds:

$$\lim_{K \to \infty} E\left[\frac{1}{KT_c} \left| T_c \sum_{k=0}^{K-1} x(k) \ e^{-j2\pi f k T_c} \right|^2 \right] = \mathcal{P}_x(f)$$
(1.374)

Then, exploiting the ergodicity of a WSS random process, one obtains the relations among the process itself, its autocorrelation function, and PSD shown in Figure 1.30. We note how the direct computation of the PSD, given by (1.374), makes use of a statistical ensemble of the Fourier transform of the process x, while the indirect method via ACS makes use of a single realization.

If we let

$$\tilde{\mathcal{X}}_{KT}(f) = T_c \mathcal{F}[x(k) w_K(k)]$$
(1.375)

where w_K is the rectangular window of length K (see (1.401)) and $T_d = KT_c$, (1.374) becomes

$$\mathcal{P}_{x}(f) = \lim_{T_{d} \to \infty} \frac{E[|\mathcal{X}_{T_{d}}(f)|^{2}]}{T_{d}}$$
(1.376)

The relation (1.376) holds also for continuous-time ergodic random processes, where $\tilde{X}_{T_d}(f)$ denotes the Fourier transform of the windowed realization of the process, with a rectangular window of duration T_d .

¹⁶ We note that for random processes with non-zero mean and/or sinusoidal components this property is not verified. Therefore, it is usually recommended that the deterministic components of the process be removed before the spectral estimation is performed.

1.9.1 Mean value estimators

Given the random process $\{x(k)\}$, we wish to estimate the mean value of a related process $\{y(k)\}$: for example to estimate the statistical power of x we set $y(k) = |x(k)|^2$, while for the estimation of the correlation of x with lag n, we set $y(k) = x(k)x^*(k - n)$. Based on a realization of $\{y(k)\}$, from (1.369) an estimate of the mean value of y is given by the expression

$$\hat{\mathbf{m}}_{y} = \frac{1}{K} \sum_{k=0}^{K-1} y(k) \tag{1.377}$$

In fact, (1.377) attempts to determine the average component of the signal $\{y(k)\}$. As illustrated in Figure 1.31a, in general, we can think of extracting the average component of $\{y(k)\}$ using an LPF filter *h* having unit gain, i.e. $\mathcal{H}(0) = 1$, and suitable bandwidth *B*. Let *K* be the length of the impulse response with support from k = 0 to k = K - 1. Note that for a unit step input signal, the transient part of the output signal will last K - 1 time instants. Therefore, we assume

$$\hat{\mathbf{m}}_{\mathbf{y}} = z(k) = h * y(k) \quad \text{for } k \ge K - 1$$
 (1.378)

We now compute the mean and variance of the estimate. From (1.378), the mean value is given by

$$E[\hat{\mathbf{m}}_{y}] = \mathbf{m}_{y} \mathcal{H}(0) = \mathbf{m}_{y} \tag{1.379}$$

as $\mathcal{H}(0) = 1$. Using the expression in Table 1.6 of the correlation of a filter output signal given the input, the variance of the estimate is given by

$$\operatorname{var}[\hat{\mathbf{m}}_{y}] = \sigma_{y}^{2} = \sum_{n=-\infty}^{+\infty} \mathbf{r}_{h}(-n)\mathbf{c}_{y}(n)$$
 (1.380)

Assuming

$$S = \sum_{n=-\infty}^{+\infty} |c_y(n)| = \sigma_y^2 \sum_{n=-\infty}^{+\infty} \frac{|c_y(n)|}{\sigma_y^2} < \infty$$
(1.381)

and being $|\mathbf{r}_h(n)| \leq \mathbf{r}_h(0)$, the variance in (1.380) is *bounded* by

$$\operatorname{var}[\hat{\mathbf{m}}_{v}] \le E_{h}S \tag{1.382}$$

where $E_h = r_h(0)$.

For an ideal lowpass filter,

$$\mathcal{H}(f) = \operatorname{rect}\left(\frac{f}{2B}\right), \quad |f| < \frac{1}{2T_c}$$
(1.383)

assuming as filter length K that of the principal lobe of $\{h(k)\}$, and neglecting a delay factor, it results as $E_h = 2B$ and $K \simeq 1/B$. Introducing the criterion that for a good estimate, it must be

$$\operatorname{var}\left[\hat{\mathbf{m}}_{\nu}\right] \le \varepsilon \tag{1.384}$$

with $\varepsilon \ll |\mathbf{m}_v|^2$, from (1.382) it follows

$$B \le \frac{\varepsilon}{2S} \tag{1.385}$$

and

$$K \ge \frac{2S}{\epsilon} \tag{1.386}$$

In other words, from (1.381) and (1.386), for a fixed ε , the length *K* of the filter impulse response must be larger, or equivalently the bandwidth *B* must be smaller, to obtain estimates for those processes $\{y(k)\}$ that exhibit larger variance and/or larger correlation among samples. Because of their simple implementation, two commonly used filters are the rectangular window and the exponential filter, whose impulse responses are shown in Figure 1.31.



Figure 1.31 (a) Time average as output of a narrow band lowpass filter. (b) Typical impulse responses: exponential filter with parameter $a = 1 - 2^{-5}$ and rectangular window with K = 33. (c) Corresponding frequency responses.

Rectangular window

For a rectangular window,

$$h(k) = \begin{cases} \frac{1}{K} & k = 0, 1, \dots, K - 1\\ 0 & \text{elsewhere} \end{cases}$$
(1.387)

the frequency response is given by (see (1.24))

$$\mathcal{H}(f) = e^{-j2\pi f \left(\frac{K-1}{2}\right)T_c} \operatorname{sinc}_K(fKT_c)$$
(1.388)

We have $E_h = 1/K$ and, adopting as bandwidth the frequency of the first zero of $|\mathcal{H}(f)|$, $B = 1/(KT_c)$. The filter output is given by

$$z(k) = \sum_{n=0}^{K-1} \frac{1}{K} y(k-n)$$
(1.389)

that can be expressed as

$$z(k) = z(k-1) + \frac{y(k) - y(k-K)}{K}$$
(1.390)

Exponential filter

For an exponential filter

$$h(k) = \begin{cases} (1-a)a^k & k \ge 0\\ 0 & \text{elsewhere} \end{cases}$$
(1.391)

with |a| < 1, the frequency response is given by

$$\mathcal{H}(f) = \frac{1 - a}{1 - ae^{-j2\pi f T_c}}$$
(1.392)

Moreover, $E_h = (1 - a)/(1 + a)$ and, adopting as length of *h* the time constant of the filter, i.e. the interval it takes for the amplitude of the impulse response to decrease of a factor *e*,

$$K - 1 = \frac{1}{\ln 1/a} \simeq \frac{1}{1 - a} \tag{1.393}$$

where the approximation holds for $a \simeq 1$. The 3 dB filter bandwidth is equal to

$$B = \frac{1-a}{2\pi} \frac{1}{T_c} \quad \text{for } a > 0.9 \tag{1.394}$$

The filter output has a simple expression given by the recursive equation

$$z(k) = az(k-1) + (1-a) y(k)$$
(1.395)

We note that choosing a as

$$a = 1 - 2^{-l} \tag{1.396}$$

then (1.395) becomes

$$z(k) = z(k-1) + 2^{-l}(y(k) - z(k-1))$$
(1.397)

whose computation requires only two additions and one *shift* of l bits. Moreover, from (1.393), the filter time constant is given by

$$K - 1 = 2^l \tag{1.398}$$

General window

In addition to the two filters described above, a general window can be defined as

$$h(k) = Aw(k) \tag{1.399}$$

with $\{w(k)\}\$ window¹⁷ of length *K*. Factor *A* in (1.399) is introduced to normalize the area of *h* to 1. We note that, for random processes with slowly time-varying statistics, (1.390) and (1.397) give an expression to update the estimates.

¹⁷ We define the *continuous-time rectangular window* with duration T_d as

$$\mathbb{W}_{T_d}(t) = \operatorname{rect}\left(\frac{t - T_d/2}{T_d}\right) = \begin{cases} 1 & 0 < t < T_d \\ 0 & \text{elsewhere} \end{cases}$$
(1.400)

Commonly used discrete-time windows are:

1.9.2 Correlation estimators

Let $\{x(k)\}, k = 0, 1, ..., K - 1$, be a realization of a random process with K samples. We examine two estimates.

Unbiased estimate

The unbiased estimate

$$\hat{\mathbf{r}}_{x}(n) = \frac{1}{K-n} \sum_{k=n}^{K-1} x(k) x^{*}(k-n) \qquad n = 0, 1, \dots, K-1$$
(1.405)

has mean

$$E[\hat{\mathbf{r}}_{x}(n)] = \frac{1}{K-n} \sum_{k=n}^{K-1} E[x(k)x^{*}(k-n)] = \mathbf{r}_{x}(n)$$
(1.406)

If the process is Gaussian, one can show that the variance of the estimate is approximately given by

$$\operatorname{var}[\hat{r}_{x}(n)] \simeq \frac{K}{(K-n)^{2}} \sum_{m=-\infty}^{+\infty} [r_{x}^{2}(m) + r_{x}(m+n)r_{x}(m-n)]$$
(1.407)

from which it follows

$$\operatorname{var}\left[\hat{r}_{x}(n)\right] \xrightarrow[K \to \infty]{} 0 \tag{1.408}$$

The above limit holds for $n \ll K$. Note that the variance of the estimate increases with the correlation lag *n*.

Biased estimate

The biased estimate

$$\check{\mathbf{r}}_{x}(n) = \frac{1}{K} \sum_{k=n}^{K-1} x(k) x^{*}(k-n) = \left(1 - \frac{|n|}{K}\right) \hat{\mathbf{r}}_{x}(n)$$
(1.409)

1. Rectangular window

$$w(k) = w_D(k) = \begin{cases} 1 \ k = 0, 1, \dots, D-1 \\ 0 \ \text{elsewhere} \end{cases}$$
 (1.401)

where D denotes the length of the rectangular window expressed in number of samples.

2. Raised cosine or Hamming window

$$w(k) = \begin{cases} 0.54 + 0.46 \cos\left(2\pi \frac{k - \frac{D-1}{2}}{D-1}\right) & k = 0, 1, \dots, D-1\\ 0 & \text{elsewhere} \end{cases}$$
(1.402)

3. Hann window

$$w(k) = \begin{cases} 0.50 + 0.50 \cos\left(2\pi \frac{k - \frac{D-1}{2}}{D-1}\right) & k = 0, 1, \dots, D-1\\ 0 & \text{elsewhere} \end{cases}$$
(1.403)

4. Triangular or Bartlett window

$$w(k) = \begin{cases} 1 - 2 \left| \frac{k - \frac{D-1}{2}}{D - 1} \right| & k = 0, 1, \dots, D - 1 \\ 0 & \text{elsewhere} \end{cases}$$
(1.404)

has mean satisfying the following relations:

$$E[\check{\mathbf{r}}_{x}(n)] = \left(1 - \frac{|n|}{K}\right) \mathbf{r}_{x}(n) \xrightarrow[K \to \infty]{} \mathbf{r}_{x}(n) \tag{1.410}$$

Unlike the unbiased estimate, the mean of the biased estimate is not equal to the autocorrelation function, but approaches it as *K* increases. Note that the biased estimate differs from the autocorrelation function by one additive constant, denoted as *bias*:

$$\mu_{bias} = E[\check{\mathbf{r}}_x(n)] - \mathbf{r}_x(n) \tag{1.411}$$

For a Gaussian process, the variance of the biased estimate is

$$\operatorname{var}[\check{\mathbf{r}}_{x}(n)] = \left(\frac{K - |n|}{K}\right)^{2} \operatorname{var}[\hat{\mathbf{r}}_{x}(n)] \simeq \frac{1}{K} \sum_{m = -\infty}^{+\infty} [\mathbf{r}_{x}^{2}(m) + \mathbf{r}_{x}(m + n)\mathbf{r}_{x}(m - n)]$$
(1.412)

In general, the biased estimate of the ACS exhibits a mean-square error¹⁸ larger than the unbiased, especially for large values of n. It should also be noted that the estimate does not necessarily yield sequences that satisfy the properties of autocorrelation functions: for example the following property may not be verified:

$$\hat{\mathbf{r}}_{\mathbf{x}}(0) \ge |\hat{\mathbf{r}}_{\mathbf{x}}(n)|, \quad n \ne 0$$
(1.414)

1.9.3 Power spectral density estimators

After examining ACS estimators, we review some spectral density estimation methods.

Periodogram or instantaneous spectrum

Let $\tilde{\mathcal{X}}(f) = T_c \mathcal{X}(f)$, where $\mathcal{X}(f)$ is the Fourier transform of $\{x(k)\}, k = 0, ..., K - 1$; an estimate of the statistical power of $\{x(k)\}$ is given by

$$\hat{\mathbb{M}}_{x} = \frac{1}{K} \sum_{k=0}^{K-1} |x(k)|^{2} = \frac{1}{KT_{c}} \int_{-\frac{1}{2T_{c}}}^{\frac{1}{2T_{c}}} |\tilde{\mathcal{X}}(f)|^{2} df$$
(1.415)

using the properties of the Fourier transform (Parseval theorem). Based on (1.415), a PSD estimator called *periodogram* is given by

$$\mathcal{P}_{PER}(f) = \frac{1}{KT_c} |\tilde{\mathcal{X}}(f)|^2 \tag{1.416}$$

We can write (1.416) as

$$\mathcal{P}_{PER}(f) = T_c \sum_{n=-(K-1)}^{K-1} \check{\mathbf{r}}_x(n) \ e^{-j2\pi f n T_c}$$
(1.417)

$$E\left[|\hat{\mathbf{r}}_{x}(n) - \mathbf{r}_{x}(n)|^{2}\right] = \operatorname{var}[\hat{\mathbf{r}}_{x}(n)] + |\mu_{bias}|^{2}$$
(1.413)

 $^{^{18}}$ For example, for the estimator (1.405) the mean-square error is defined as

and, consequently,

$$E[\mathcal{P}_{PER}(f)] = T_c \sum_{n=-(K-1)}^{K-1} E[\check{\mathbf{r}}_x(n)] e^{-j2\pi f n T_c}$$

= $T_c \sum_{n=-(K-1)}^{K-1} \left(1 - \frac{|n|}{K}\right) \mathbf{r}_x(n) e^{-j2\pi f n T_c}$
= $T_c \mathcal{W}_B * \mathcal{P}_x(f)$ (1.418)

where $\mathcal{W}_{B}(f)$ is the Fourier transform of the symmetric Bartlett window

$$w_B(n) = \begin{cases} 1 - \frac{|n|}{K} & |n| \le K - 1\\ 0 & |n| > K - 1 \end{cases}$$
(1.419)

and

$$\mathcal{W}_B(f) = K \left[\operatorname{sinc}_K(fkT_c) \right]^2 \tag{1.420}$$

We note the periodogram estimate is affected by *bias* for finite *K*. Moreover, it also exhibits a large variance, as $\mathcal{P}_{PER}(f)$ is computed using the samples of $\check{r}_x(n)$ even for lags up to K - 1, whose variance is very large.

Welch periodogram

This method is based on applying (1.374) for finite *K*. Given a sequence of *K* samples, different subsequences of consecutive *D* samples are extracted. Subsequences may partially overlap. Let $x^{(s)}$ be the *s*-th subsequence, characterized by *S* samples in common with the preceding subsequence $x^{(s-1)}$ and with the following one $x^{(s+1)}$. In general, $0 \le S \le D/2$, with the choice S = 0 yielding subsequences with no overlap and therefore with less correlation. The number of subsequences N_s is¹⁹

$$N_s = \left\lfloor \frac{K - D}{D - S} + 1 \right\rfloor \tag{1.421}$$

Let w be a window (see footnote 17 on page 59) of D samples: then

$$x^{(s)}(k) = w(k) x(k + s(D - S)), \qquad k = 0, 1, \dots, D - 1 \ s = 0, 1, \dots, N_s - 1$$
(1.422)

For each s, compute the Fourier transform

$$\tilde{\mathcal{X}}^{(s)}(f) = T_c \sum_{k=0}^{D-1} x^{(s)}(k) e^{-j2\pi f k T_c}$$
(1.423)

and obtain

$$\mathcal{P}_{PER}^{(s)}(f) = \frac{1}{DT_c M_w} |\tilde{\mathcal{X}}^{(s)}(f)|^2$$
(1.424)

where

$$M_{\rm w} = \frac{1}{D} \sum_{k=0}^{D-1} {\rm w}^2(k) \tag{1.425}$$

is the normalized energy of the window. As a last step, for each frequency, average the periodograms:

$$\mathcal{P}_{WE}(f) = \frac{1}{N_s} \sum_{s=0}^{N_s - 1} \mathcal{P}_{PER}^{(s)}(f)$$
(1.426)

¹⁹ The symbol $\lfloor a \rfloor$ denotes the function *floor*, that is the largest integer smaller than or equal to *a*. The symbol $\lceil a \rceil$ denotes the function *ceiling*, that is the smallest integer larger than or equal to *a*.

The mean of the estimate is given by

$$E[\mathcal{P}_{WE}(f)] = T_c[|\mathcal{W}|^2 * \mathcal{P}_x](f) \tag{1.427}$$

where

$$\mathcal{W}(f) = \sum_{k=0}^{D-1} w(k) e^{-j2\pi f k T_c}$$
(1.428)

Assuming the process Gaussian and the different subsequences statistically independent, we get²⁰

$$\operatorname{var}[\mathcal{P}_{WE}(f)] \propto \frac{1}{N_s} \mathcal{P}_x^2(f) \tag{1.429}$$

Note that the partial overlap introduces correlation between subsequences. From (1.429), we see that the variance of the estimate is reduced by increasing the number of subsequences. In general, D must be large enough so that the *generic subsequence represents the process*²¹ and also N_s must be large to obtain a reliable estimate (see (1.429)); therefore, the application of the Welch method requires many samples.

Blackman and Tukey correlogram

For an *unbiased* estimate of the ACS, $\{\hat{\mathbf{r}}_{x}(n)\}, n = -L, \dots, L$, consider the Fourier transform

$$\mathcal{P}_{BT}(f) = T_c \sum_{n=-L}^{L} w(n) \hat{\mathbf{r}}_x(n) \ e^{-j2\pi f n T_c}$$
(1.430)

where w is a window²² of length 2L + 1, with w(0) = 1. If K is the number of samples of the realization sequence, we require that $L \le K/5$ to reduce the variance of the estimate. Then if the Bartlett window (1.420) is chosen, one finds that $\mathcal{P}_{BT}(f) \ge 0$.

In terms of the mean value of the estimate, we find

$$E[\mathcal{P}_{BT}(f)] = T_c(\mathcal{W} * \mathcal{P}_x)(f) \tag{1.431}$$

For a Gaussian process, if the Bartlett window is chosen, the variance of the estimate is given by

$$\operatorname{var}[\mathcal{P}_{BT}(f)] = \frac{1}{K} \mathcal{P}_{x}^{2}(f) E_{w} = \frac{2}{3} \frac{L}{K} \mathcal{P}_{x}^{2}(f)$$
(1.432)

Windowing and window closing

The windowing operation of time sequence in the periodogram, and of the ACS in the correlogram, has a strong effect on the performance of the estimate. In fact, any truncation of a sequence is equivalent to a windowing operation, carried out via the rect function. The choice of the window type in the frequency domain depends on the compromise between a narrow central lobe (to reduce *smearing*) and a fast decay of secondary lobes (to reduce *leakage*). *Smearing* yields a lower spectral resolution, that is the capability to distinguish two spectral lines that are close. On the other hand, *leakage* can mask spectral components that are further apart and have different amplitudes.

The choice of the window length is based on the compromise between spectral resolution and the variance of the estimate. An example has already been seen in the correlogram, where the condition $L \le K/5$ must be satisfied. Another example is the Welch periodogram. For a given observation of K samples, it is

²⁰ Notation $a \propto b$ means that *a* is proportional to *b*.

²¹ For example, if x is a sinusoidal process, DT_c must at least be greater than 5 or 10 periods of x.

²² The windows used in (1.430) are the same introduced in footnote 17: the only difference is that they are now centered around zero instead of (D-1)/2. To simplify the notation, we will use the same symbol in both cases.

initially better to choose a small number of samples over which to perform the DFT, and therefore a large number of windows (subsequences) over which to average the estimate. The estimate is then repeated by increasing the number of samples per window, thus decreasing the number of windows. In this way, we get estimates with not only a higher resolution but also characterized by an increasing variance. The procedure is terminated once it is found that the increase in variance is no longer compensated by an increase in the spectral resolution. The aforementioned method is called *window closing*.

Example 1.9.1

Consider a realization of $K = 10\ 000$ samples of the signal:

$$y(kT_c) = \frac{1}{A_h} \sum_{n=-16}^{16} h(nT_c) w((k-n)T_c) + A_1 \cos(2\pi f_1 kT_c + \varphi_1) + A_2 \cos(2\pi f_2 kT_c + \varphi_2)$$
(1.433)

where $\varphi_1, \varphi_2 \sim U[0, 2\pi), w(nT_c)$ is a white random process with zero mean and variance $\sigma_w^2 = 5, T_c = 0.2, A_1 = 1/20, f_1 = 1.5, A_2 = 1/40, f_2 = 1.75$, and

$$A_h = \sum_{-16}^{16} h(kT_c) \tag{1.434}$$

Moreover

$$h(kT_c) = \frac{\sin\left(\pi(1-\rho)\frac{kT_c}{T}\right) + 4\rho\frac{kT_c}{T}\cos\left(\pi(1+\rho)\frac{kT_c}{T}\right)}{\pi\left[1 - \left(4\rho\frac{kT_c}{T}\right)^2\right]\frac{kT_c}{T}}\operatorname{rect}\left(\frac{kT_c}{8T+T_c}\right)$$
(1.435)

with $T = 4T_c$ and $\rho = 0.32$.

Actually y is the sum of two sinusoidal signals and filtered white noise through h. Consequently, observing (1.188) and (1.313),

$$\mathcal{P}_{y}(f) = \sigma_{w}^{2} T_{c} \frac{|\mathcal{H}(f)|^{2}}{A_{h}^{2}} + \frac{A_{1}^{2}}{4} (\delta(f - f_{1}) + \delta(f + f_{1})) + \frac{A_{2}^{2}}{4} (\delta(f - f_{2}) + \delta(f + f_{2}))$$
(1.436)

where $\mathcal{H}(f)$ is the Fourier transform of $\{h(kT_c)\}$.

The shape of the PSD in (1.436) is shown in Figures 1.32–1.34 as a solid line. A Dirac impulse is represented by an isosceles triangle having a base equal to twice the desired frequency resolution F_q . Consequently, a Dirac impulse, for example of area $A_1^2/4$ will have a height equal to $A_1^2/(4F_q)$, thus, maintaining the equivalence in statistical power between different representations.

We now compare several spectral estimates, obtained using the previously described methods; in particular, we will emphasize the effect on the resolution of the type of window used and the number of samples for each window.

We state beforehand the following important result. Windowing a complex sinusoidal signal $\{e^{j2\pi f_1 kT_c}\}$ with $\{w(k)\}$ produces a signal having Fourier transform equal to $\mathcal{W}(f - f_1)$, where $\mathcal{W}(f)$ is the Fourier transform of w. Therefore, in the frequency domain, the spectral line of a sinusoidal signal becomes a signal with shape $\mathcal{W}(f)$ centred around f_1 .

In general, from (1.424), the periodogram of a real sinusoidal signal with amplitude A_1 and frequency f_1 is

$$\mathcal{P}_{PER}(f) = \frac{T_c}{DM_w} \left(\frac{A_1}{2}\right)^2 |\mathcal{W}(f - f_1) + \mathcal{W}(f + f_1)|^2$$
(1.437)



Figure 1.32 Comparison between spectral estimates obtained with Welch periodogram method, using the Hamming or the rectangular window, and the analytical PSD given by (1.436).

Figure 1.32 shows, in addition to the analytical PSD (1.436), the estimate obtained by the Welch periodogram method using the Hamming or the rectangular windows. Parameters used in (1.423) and (1.426) are: D = 1000, $N_s = 19$, and 50% overlap between windows. We observe that the use of the Hamming window yields an improvement of the estimate due to less *leakage*. Likewise Figure 1.33 shows how the Hamming window also improves the estimate carried out with the correlogram; in particular, the estimates of Figure 1.33 were obtained using in (1.430) L = 500. Lastly, Figure 1.34 shows how the resolution and the variance of the estimate obtained by the Welch periodogram vary with the parameters D and N_s , using the Hamming window. Note that by increasing D, and hence decreasing N_s , both resolution and variance of the estimate increase.

1.10 Parametric models of random processes

ARMA

Let us consider the realization of a random process x according to the *auto-regressive moving average* (ARMA) model illustrated in Figure 1.35. In other words, the process x, also called observed sequence, is the output of an IIR filter having as input white noise with variance σ_w^2 , and is given by the recursive



Figure 1.33 Comparison between spectral estimates obtained with the correlogram using the Hamming or the rectangular window, and the analytical PSD given by (1.436).

equation23

$$x(k) = -\sum_{n=1}^{p} a_n x(k-n) + \sum_{n=0}^{q} b_n w(k-n)$$
(1.438)

and the model is denoted as ARMA(p, q).

Rewriting (1.438) in terms of the filter impulse response h_{ARMA} , we find in general

$$x(k) = \sum_{n=0}^{+\infty} h_{ARMA}(n)w(k-n)$$
(1.439)

which indicates that the filter used to realize the ARMA model is causal. From (1.63), one finds that the filter transfer function is given by

$$H_{ARMA}(z) = \frac{B(z)}{A(z)} \text{ where } \begin{cases} B(z) = \sum_{n=0}^{q} b_n z^{-n} \\ A(z) = \sum_{n=0}^{p} a_n z^{-n} \text{ assuming } a_0 = 1 \end{cases}$$
(1.440)

Using (1.188), the PSD of the process x is given by

$$\mathcal{P}_{x}(f) = T_{c}\sigma_{w}^{2} \left| \frac{\mathcal{B}(f)}{\mathcal{A}(f)} \right|^{2} \quad \text{where} \quad \begin{cases} \mathcal{B}(f) = \mathcal{B}(e^{j2\pi fT_{c}}) \\ \mathcal{A}(f) = \mathcal{A}(e^{j2\pi fT_{c}}) \end{cases}$$
(1.441)

²³ In a simulation of the process, the first samples x(k) generated by (1.438) should be neglected because they depend on the initial conditions. Specifically, if N_{ARMA} is the *length* of the filter impulse response h_{ARMA} , the minimum number of samples to be ignored is $N_{ARMA} - 1$, equal to the filter transient.